

Reduced order models for parameterized problems: Lecture One

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SPRINGER BRIEFS IN MATHEMATIC

Certified Reduced

Basis Methods

Equations

(bcar

for Parametrized

Partial Differentia

Lecture I: Introduction, motivation, basics

Lecture 2: Certified reduced methods

Lecture 3: The 'non's 'etc

Hesthaven, Rozza, Stamm Certified Reduced Basis Methods for Parametrized Partial Differential Equations Springer Briefs in Mathematics, 2015

Free: https://infoscience.epfl.ch/record/213266?ln=en





Understand Reduced models





What we seek



What we need is an accurate way to evaluate the solution at new parameter values at reduced complexity.



What we need is an **accurate** way to evaluate the solution at new parameter values at reduced complexity.

Fast input-output procedure:

input: parameter value $\mu \in \mathcal{D}$





We do not consider reduced physics -

<u>High-frequency vs low-frequency EM</u>

 $\nabla \times \nabla \times \mathbf{E} + \omega^2 \mathbf{E} = \mathbf{f} \qquad \mathbf{V}\mathbf{S} \qquad -\nabla^2 \mathbf{E} = \mathbf{f}$

Viscous vs inviscid fluid flows

$$\begin{aligned} \frac{\partial \mathbf{u}}{\partial t} + \mathbf{u} \cdot \nabla \mathbf{u} &= -\nabla \mathbf{p} + \nu \nabla^2 \mathbf{u} & \frac{\partial \mathbf{u}}{\partial t} + \mathbf{u} \cdot \nabla \mathbf{u} &= -\nabla \mathbf{p} \\ \nabla \cdot \mathbf{u} &= \mathbf{0} & \text{vs} & \nabla \cdot \mathbf{u} &= \mathbf{0} \end{aligned}$$



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$$\begin{split} \frac{\partial \mathbf{u}}{\partial t} + \mathbf{u} \cdot \nabla \mathbf{u} &= -\nabla \mathbf{p} + \nu \nabla^2 \mathbf{u} & \frac{\partial \mathbf{u}}{\partial t} + \mathbf{u} \cdot \nabla \mathbf{u} &= -\nabla \mathbf{p} \\ \nabla \cdot \mathbf{u} &= \mathbf{0} & \mathbf{vs} & \nabla \cdot \mathbf{u} &= \mathbf{0} \end{split}$$

.. but reduced representations of the full problem





Assume we are interested in

$$-\nabla^2 u(\mathbf{x},\mu) = \mathbf{f}(\mathbf{x},\mu) \qquad \mathbf{x} \in \mathbf{\Omega}$$

and wish to solve it accurately for many values of 'some' parameter μ





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and wish to solve it accurately for many values of 'some' parameter μ

We can use our favorite numerical method

$$A_h \mathbf{u}_h(\mathbf{x}, \mu) = \mathbf{f}_h(\mathbf{x}, \mu) \qquad \dim(\mathbf{u}_h) = \mathcal{N} \gg 1$$

For many parameter values, this is expensive - and slow !



Assume we (somehow) know

$$\mathbf{u}_{\mathbf{h}}(\mathbf{x}, \mu) \simeq \mathbf{u}_{\mathbf{RB}}(\mathbf{x}, \mu) = \mathbf{V}\mathbf{a}(\mu) \qquad \mathbf{V}^T \mathbf{V} = \mathbf{I}$$
$$\dim(\mathbf{a}) = N \qquad \dim(V) = \mathcal{N} \times N$$



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Then we can recover a solution for a new parameter as little cost

$$(\mathbf{V}^T \mathbf{A}_h \mathbf{V}) \mathbf{V}^T \mathbf{u}_h(\mu) = \mathbf{V}^T \mathbf{f}_h(\mu)$$



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So IF

-)...we know the orthonormal basis $\,\,\,\,\,V$
-)... and it allows an accurate representation $u_{RB}(\mu)$
- ...and we can evaluate RHS 'fast'- $\mathcal{O}(N)$
- we can evaluate new solutions at cost $\mathcal{O}(N)$



So IF

-)...we know the orthonormal basis $\,\,\,\,\,V$
-)... and it allows an accurate representation $u_{RB}(\mu)$
- ... and we can evaluate RHS 'fast'-
- we can evaluate new solutions at cost -

So WHY ? - a promise to do more with less





Examples in many application domains

- Optimization/inversion/control problems
- Simulation based data bases
- Uncertainty quantification
- Sub-scale models in multi-scale modeling
- In-situ/deployed modeling

When is that relevant ?



Examples in many application domains

Optimization/inver





D. Knezevic et al, 2010



We consider projection based techniques, i.e.

$$\mathbf{u}_{\mathbf{h}}(\mathbf{x},\mu) \simeq \mathbf{u}_{\mathbf{RB}}(\mathbf{x},\mu) = \mathbf{Va}(\mu)$$

There is a substantial literature for linear systems

$$\begin{cases} \boldsymbol{C}\dot{\boldsymbol{x}}(t) + \boldsymbol{G}\boldsymbol{x}(t) = \boldsymbol{B}\boldsymbol{u}(t), \\ \boldsymbol{y}(t) = \boldsymbol{L}^{\mathrm{T}}\boldsymbol{x}(t), \end{cases} \Leftrightarrow \begin{cases} \boldsymbol{C}_{n}\dot{\boldsymbol{z}}(t) + \boldsymbol{G}_{n}\boldsymbol{z}(t) = \boldsymbol{B}_{n}\boldsymbol{u}(t), \\ \tilde{\boldsymbol{y}}(t) = \boldsymbol{L}_{n}^{\mathrm{T}}\boldsymbol{z}(t), \end{cases}$$

Typically seeks to approximate the transfer function $H(s) = L^{T}(G + sC)^{-1}B. \iff H_{n}(s) = L_{n}^{T}(G_{n} + sC_{n})^{-1}B_{n}.$

Pade approximations
Krylov subspace methods
Balanced truncation

Non-linear problems ?















Parametrized problem setting:

Let $\mu = (\mu_1, \mu_2, \dots, \mu_{15}) \in \mathbb{P} = [\mu_-, \mu_+]^{15}$. Then, for any $\mu \in \mathbb{P}$, compute $s(u(\mu))$.

Solutions $u(\mu)$ for different values of $\mu \in \mathbb{P}$:





Variational setting: Define

$$\mathbb{V} = \{ v \in H^{1}(\Omega) \mid v|_{\Gamma_{top}} = 0 \},\$$
$$a(w, v; \mu) = \sum_{i=1}^{15} \mu_{i} \int_{\mathcal{R}_{i}} \nabla w \cdot \nabla v + \int_{\mathcal{R}_{P+1}} \nabla w \cdot \nabla v,\$$
$$f(v) = \ell(v) = \int_{\Gamma_{1}} v.$$

Parametrized problem setting:

For any $\mu \in \mathbb{P}$, find $u(\mu) \in \mathbb{V}$ s.t.

$$a(u(\mu), v; \mu) = f(v), \quad \forall v \in \mathbb{V}.$$

Then, compute $s(\mu) = \ell(u(\mu)) = \int_{\Gamma_1} u(\mu)$.







Assume that Ω is a homogenous media with magnetic permeability μ_0 and electrical permittivity ε_0 . Assume that Ω is a homogenous media with magnetic permeability μ_0 and electrical permeability μ_0 and μ_0 and electrical permeability μ_0 and μ_0

Assume that Ω is a homogenous media with magnetic permeability μ_0 and elec-The control of the set of th

 $\begin{array}{c} \mbox{Then, the electric field } E(\mu) = E^i(\mu) + E^s(\mu) \in H(\operatorname{curl},\Omega) \text{ satisfies} \\ \mbox{curl curl } E(\mu) - k^2 E(\mu) = 0 & \text{in } \Omega, & \text{Maxwell} \\ \mbox{curl curl } E(\mu) = E(\mu) E(\mu) = 0 & \text{oin} \Gamma\Omega, & \text{boundary conditions} \\ \mbox{curl} E^s(x;\mu) \times \frac{x}{|x|} - ik E^t(\mu);\mu) = \mathcal{O}\left(\frac{1}{|x|}\right) & \mbox{aspin} \Gamma \rightarrow \infty. & \text{Silver-Müllerbaundars condition} \\ \mbox{curl} E^s(x;\mu) \times \frac{x}{|x|} - ik E^s(x;\mu) = \mathcal{O}\left(\frac{1}{|x|}\right) & \mbox{aspin} \Gamma \rightarrow \infty. & \text{Silver-Müller radiation cond.} \\ \mbox{Boundary condition is equivalent to } \gamma_t E(\mu) = 0 & \text{where } \gamma_t & \text{denotes the tangential} \\ \mbox{Trace operator on surface } F, \gamma_t E(\mu) = n \times (E(\mu) \times n). & \mbox{Ic} \\ \mbox{Jeries operator on surface } \Gamma, \gamma_t E(\mu) = n \times (E(\mu) \times n). & \mbox{Ic} \\ \mbox{Jeries operator on surface } \Gamma, \gamma_t E(\mu) = n \times (E(\mu) \times n). & \mbox{Jeries operator on surface } \Gamma, \gamma_t E(\mu) = n \times (E(\mu) \times n). \end{array}$

 $k=\omega\sqrt{\mu_0\varepsilon_0}$ is wave number and ω the angular frequency of the time-harmonic ansatz

$$\hat{E}(x,t;\mu) = e^{-i\omega t} E(x;\mu).$$







Parametrized problems

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The parameters can describe

- Materials
- Sources
- Geometries
- Parameterized uncertainty
- Time
- etc



Parametrized problems



The parameters can describe

- Materials
- Sources
- Geometries
- Parameterized uncertainty
- Time
- etc



Does this always work, i.e., does a reduced model always exist ?

Probably not - we need to understand when and how to check

The solution manifold



Consider the "exact" and "discrete" solution manifolds

Consider the "exact" and \mathcal{M} is crete "(solution mappinolds \mathbb{V} ,

$$\mathcal{M} = \{u(\mu); \forall \mu \in \mathbb{P}\} \subset \mathbb{V},$$
$$\mathcal{M}_{\delta} = \{u_{\delta}(\mu); \forall \mu \in \mathbb{P}\} \subset \mathbb{V}_{\delta},$$
$$\mathcal{M}_{\delta} = \{u_{\delta}(\mu); \forall \mu \in \mathbb{P}\} \subset \mathbb{V}_{\delta},$$

where, for each $\mu \in \mathbb{P}$, $u(\mu)$ and $u_{\delta}(\mu)$ denote the solution of the underlying where, for each $\mu \in \mathbb{P}$, $u(\mu)$ and $u_{\delta}(\mu)$ denote the solution of the <u>underlying</u> exact and discrete problems respectively.



The solution manifold



The key question is how well can the solution manifold \mathcal{M} be approximated by \mathcal{M}_{δ} using an N-dimensional linear space ?



Clearly, if the solution space is (locally) smooth we have a good chance.

The solution manifold



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Clearly, if the solution space is (locally) smooth we have a good chance.

Highly sensitive/chaotic systems will be problematic as they have no structure

v well can the solution manifold \mathcal{M} resp. \mathcal{M}_{δ} be approximated by ensional finear solution manifold \mathcal{M} resp. \mathcal{M}_{δ} be approximated by ensional finear space \mathcal{M}_{rb}

For any N-dimensional space we define

$$E(\mathcal{M}, \mathbb{V}_{rb}) = \sup_{u(\mu) \in \mathcal{M}} \inf_{v_{rb} \in \mathbb{V}_{rb}} \|u(\mu) - v_{rb}\|_{\mathbb{V}}.$$

e Kolmogorov N-width is then defined by

$$d_N(\mathcal{M}) = \inf_{\mathbb{V}_{rb}} E(\mathcal{M}, \mathbb{V}_{rb}) = \inf_{\mathbb{V}_{rb}} \sup_{u(\mu) \in \mathcal{M}} \inf_{v_{rb} \in \mathbb{V}_{rb}} \|u(\mu) - v_{rb}\|_{\mathbb{V}},$$

ere the first infimum is taken over all N-dimensional subspaces \mathbb{V}_{rb}

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How well can the solution manifold \mathcal{M} resp. \mathcal{M}_{δ} be approximated by an dimensional dimensional space we define

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$$- \left(d_N(\mathcal{M}) = \inf_{\mathbb{V}_{rb}} E(\mathcal{M}, \mathbb{V}_{rb}) = \inf_{\mathbb{V}_{rb}} \sup_{u(\mu) \in \mathcal{M}} \inf_{v_{rb} \in \mathbb{V}_{rb}} \|u(\mu) - v_{rb}\|_{\mathbb{V}}, \right)$$

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How well can the solution manifold \mathcal{M} resp. \mathcal{M}_{δ} be approximated by an dimensional dimensional space we define

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ere the first infimum is taken over all N-dimensional subspaces V_{rb} of **If this decays rapidly with N, we are in good shape**

Identifying the optimal linear space by the Kolmogorov N-width is not practical — cost exponential in N


The behavior of the Kolmogorov N-width is non-trivial

$$\mathcal{M} = \Big\{ u(x,\mu) = |x-\mu|^{0.5} \ \Big| \ \mu, x \in (0,1) \Big\}.$$

Singularity at varying location.





sider on heebehaviom of the Kolmogorov N-width is non-trivial

$$\mathcal{M} = \left\{ \mathcal{M}(\underline{x}, \underline{k})_{u}(\overline{x}, \underline{k}) = \mu | x - \mu | x -$$





$$\mathcal{M} = \Big\{ u(x,\mu) = |x - 0.5|^{\mu + 0.5} \ \Big| \ \mu, x \in (0,1) \Big\}.$$

Singularity at fixed location with varying width.



Consider on the other hand



$$\mathcal{M} = \mathcal{M}\left\{ u \in x \left\{ u \in x, \mu \right\} = \left[x 5 \downarrow^{\mu} 0.5 \right]^{\mu} = \left[x 4 \downarrow^{\mu} 0.5 \right]^{\mu} = \left[x$$



Morale: We need to check if a reduced space exists before going ahead



We can get a good sense by a feasibility study

o Define a point-set $\mathbb{P}_h = \{\mu_1, \ldots, \mu_M\} \subset \mathbb{P}$.

• Compute for each μ_i the truth solution $u(\mu_i)$ using a simplified model.

• Store the degrees of freedom row-wise in a matrix A.



This samples the solution manifold



3D EM scattering with the angle varying 0-360 deg. RCS is computed every 2 deg.

Computing the SVD of the 180 solutions shows that less than 60 samples would suffice -- and likely much less for applications



Computation by CERFACS



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Computation by CERFACS







Answer: Depends really on the problem. For uniformly elliptic problems,

Compteting the swip of the solution hat is gives ault exists but exponential convergence is in practise observed and for other problems it measures of the Kolmogorov N-width

Parameters n(k,d) seriet 25 at [On π] of 19 fixed \times [0, π], compute the BEM-

For solution is contracted parameter of the scattering problem: two dimensional parameterization with polar and contracted the singular values of the scattering problem: two dimensional parameterization with polar and contracted the singular values of the singular values of



Rigorous results are sparse for this - but there are some of the nature $d_N(\mathcal{M}) \leq C e^{-cN}$

Basic setting



We consider physical systems of the form

$$\mathcal{L}(\mathbf{x}, \mu) u(\mathbf{x}, \mu) = f(\mathbf{x}, \mu) \qquad \mathbf{x} \in \Omega$$
$$u(\mathbf{x}, \mu) = g(\mathbf{x}, \mu) \qquad \mathbf{x} \in \partial \Omega$$

where the solutions are implicitly parameterized by

$$\mu \in \mathcal{D} \in \mathcal{R}^N$$



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where the solutions are implicitly parameterized by

$$\mu \in \mathcal{D} \in \mathcal{R}^N$$

- How do we find the basis.
- How do we ensure accuracy under parameter variation ?
- What about speed ?

NSolutions and their behavior ation space. We have the solution of the parameter value $\mu \in \mathbb{P}$, and $a(\mu) \in \mathbb{V}$ such that Exact solution: For some parameter value $\mu \in \mathbb{P}$, and $a(\mu) \in \mathbb{V}$ such that $a(u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in \mathbb{V}$. Then, compute the value of the output functional $s(\mu) = \ell(u(\mu); \mu)$. Then, compute the value of the output functional $s(\mu) = \ell(u(\mu); \mu)$. In practise, the exact PDE cannot be solved. A popular discretisation technique Is phactGaletrk in a appBotachmoR eplaotyche A" comptibutolis constrained to the value of the solved of the solved

is in then **Goaler kirb**s **ppproach** CR that the "continuous" space \mathbb{V} by the finite $\frac{1}{1}$ it is dimensional subspace \mathbb{W} and that

$$\lim_{\delta \to 0} \inf_{v_{\delta} \in \mathbb{V}_{\delta}} \|v - v_{\delta}\|_{\mathbb{V}} = 0, \qquad \forall v \in \mathbb{V}.$$

Galerkin solution: For some parameter $\mu \in \mathbb{P}$, find $u_{\delta}(\mu) \in \mathbb{V}_{\delta}$ such that

$$a(u_{\delta}(\mu), v_{\delta}; \mu) = f(v_{\delta}; \mu) \quad \forall v_{\delta} \in \mathbb{V}_{\delta}.$$

Then, compute the value of the output functional $s_{\delta}(\mu) = \ell(u_{\delta}(\mu); \mu)$. Then, comp $\mathbf{A}_{\delta}^{\mu} \mathbf{u}_{\delta}^{\mu} = \mathbf{f}_{\delta}^{\mu}$.^{output} functional $s_{\delta}(\mu) = (\mathbf{u}_{\delta}^{\mu})^{T} \mathbf{f}_{\delta}^{\mu}$.

Convergence and stability



For coercive problems: Since $\mathbb{V}_{\delta} \subset \mathbb{V}$, there holds that the discrete coercivity constant

$$\alpha_{\delta}(\mu) = \inf_{v_{\delta} \in \mathbb{V}_{\delta}} \frac{a(v_{\delta}, v_{\delta}; \mu)}{\|v_{\delta}\|_{\mathbb{W}}^2}$$

is uniformly bounded from below:

$$0 < \alpha \le \alpha_{\delta}(\mu),$$

since

$$0 < \alpha \le \alpha(\mu) = \inf_{v \in \mathbb{V}} \frac{a(v, v; \mu)}{\|v\|_{\mathbb{V}}^2} \le \inf_{v_{\delta} \in \mathbb{V}_{\delta}} \frac{a(v_{\delta}, v_{\delta}; \mu)}{\|v_{\delta}\|_{\mathbb{V}}^2}$$

Convergence and stability



For coercive problems: Since $\mathbb{V}_{\delta} \subset \mathbb{V}$, there holds that the discrete coercivity constant $u(\mu)$: Exact solution

 $u_{\delta}(\mu) : \operatorname{Truth}(\operatorname{Discretef}_{v_{\delta} \in \mathbb{V}_{\delta}} \operatorname{Galerkin}; \operatorname{H}) \text{ proximation} \\ v_{\delta} : \operatorname{Any} \operatorname{discrete}^{v_{\delta} \in \mathbb{V}_{\delta}} \operatorname{function}^{v_{\delta}}(\operatorname{V}_{\delta}^{2} \in \mathbb{V}_{\delta})$

is unif cody diving, Cfaver kin with ogonality and continuity (Cea's lemm

since

$$0 < \begin{bmatrix} \|u(\mu) - u_{\delta}(\mu)\|_{\mathbb{V}}^{2} & \stackrel{\text{coer.}}{\leq} & \frac{1}{\alpha(\mu)}a\Big(u(\mu) - u_{\delta}(\mu), u(\mu) - u_{\delta}(\mu); \mu\Big) \\ & \stackrel{\text{G.O.}}{=} & \frac{1}{\alpha(\mu)}a\Big(u(\mu) - u_{\delta}(\mu), u(\mu) - v_{\delta}; \mu\Big) \\ & \stackrel{\text{cont.}}{\leq} & \frac{\gamma(\mu)}{\alpha(\mu)}\|u(\mu) - u_{\delta}(\mu)\|_{\mathbb{V}}\|u(\mu) - v_{\delta}\|_{\mathbb{V}} \\ & \text{Then} \end{bmatrix}$$

$$\left\| u(\mu) - u_{\delta}(\mu) \|_{\mathbb{V}} \le \frac{\gamma(\mu)}{\alpha(\mu)} \| u(\mu) - v_{\delta} \|_{\mathbb{V}}, \quad \forall v_{\delta} \in \mathbb{V}_{\delta} \right)$$



Ultimative goal: a fast input-output computation: $\mu \mapsto s(\mu)$.

That is, for each $\mu \in \mathbb{P}$, evaluate $s(\mu) = \ell(u(\mu); \mu)$ where $u(\mu)$ is the solution of the parametrised weak problem: Find $u(\mu) \in \mathbb{V}$ s.t.

$$a(u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in \mathbb{V}.$$

In practise, the truth approximation is considered (instead of the exact solution): for $\mu \in \mathbb{P}$, evaluate $s_{\delta}(\mu) = \ell(u_{\delta}(\mu); \mu)$ where $u_{\delta}(\mu)$ is solution of: Find $u_{\delta}(\mu) \in \mathbb{V}_{\delta}$ s.t.

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Poor man's approach or **brute force approach**: For each new $\mu \in \mathbb{P}$, solve the above discrete problem to obtain $u_{\delta}(\mu)$ and compute the output $s_{\delta}(\mu)$.

One evaluation of the map $s_{\delta} : \mu \mapsto s_{\delta}(\mu)$ is quite expensive for some desired accuracy tol. Ok for a couple of evaluations but not in a many query context where many evaluations are need.



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One evaluation of the map $s_{\delta} : \mu \mapsto s_{\delta}(\mu)$ is quite expensive for some desired accuracy tol. Ok for a couple of evaluations but not in a many query context where many evaluations are need:



Ultimative goal: a fast input-output computation: $\mu \mapsto s(\mu)$. Ultimative goal: a fast input-output computation: $\mu \mapsto s(\mu)$. That is, for each $\mu \in \mathbb{P}$, evaluate $s(\mu) = \ell(u(\mu); \mu)$ where $u(\mu)$ is the solution of the parametrised weak problem Find $u(\mu) \in W$ is the solution of the parametrised weak problem μ Find $u(\mu) \in W$ is $\forall v \in W$. In practice, the truth approximation is considered (instead of the exact solution): $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. In practice, the truth approximation is considered (instead of the exact solution): $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. In practice, the truth approximation is considered (instead of the exact solution): $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$. $\lim_{t \to 0} |u(\mu), v; \mu) = f(v; \mu), \quad \forall v \in W$.

 $\begin{array}{c} a(u_{\delta}(\mu), v_{\delta}; \mu) \equiv f(v_{\delta}; \mu), \quad \forall v_{\delta} \in \mathbb{V}_{\delta}, \\ a(u_{\delta}(\mu), v_{\delta}; \mu) = f(v_{\delta}; \mu), \quad \forall v_{\delta} \in \mathbb{V}_{\delta}. \end{array}$

Poor man's approach of brute force approach. For each new $\mu \in \mathbb{P}$, solve the above discrete problem to obtain $u_{sl}(\mu)$ and compute the output $s_s(\mu)$. **Poor man's approach** or brute force approach: For each new $\mu \in \mathbb{P}$, solve (the above discrete problem to obtain $u_{sl}(\mu)$ and compute the output $s_{\delta}(\mu)$. (the above discrete problem to obtain $u_{\delta}(\mu)$ and compute the output $s_{\delta}(\mu)$. accuracy tol. Ok for a couple of evaluations but not in a many query context where many evaluations are need.



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Seeking
$$u(x,t) \simeq u_{\delta}(x,t) \simeq V u_{rb}$$

Let \mathbb{P}_h be a finite set of M points in \mathbb{P} that are sampled "finely". Introduce the error measure

$$\sqrt{\frac{1}{M}\sum_{\mu\in\mathbb{P}_h}\inf_{v_{\mathsf{rb}}\in\mathbb{V}_{\mathsf{rb}}}\|u(\mu)-v_{\mathsf{rb}}\|_{\mathbb{V}}^2}.$$

 \Rightarrow Average error of best approximation in \mathbb{V}_{rb} over \mathbb{P}_h .

The POD-space \mathbb{V}_{POD} is the *N*-dimensional sub-space of \mathbb{V}_{δ} that minimises the above error measure.

Remark: The corresponding discrete Kolmogorov space would be the one that minimises the error measure

$$\sup_{\mu \in \mathbb{P}_h} \inf_{v_{\mathsf{rb}} \in \mathbb{V}_{\mathsf{rb}}} \| u(\mu) - v_{\mathsf{rb}} \|_{\mathbb{V}}.$$

Thus, this is question of L^{∞} vs. L^2 over the parameter space \mathbb{P}_h .



Seeking $u(x,t) \simeq u_{\delta}(x,t) \simeq V u_{rb}$

 $\label{eq:phile} $$ I = \frac{1}{2} h = \frac{1}$

$$\sqrt{\frac{1}{M}} \sum_{\mu \notin \mathbb{P}_h^{h}} \inf_{v_{rb}^{v_{rb}} \notin \mathbb{V}_{rb}^{rb}} \| u(\mu) - v_{rb}^{v_{rb}} \|_{\mathbb{V}}^2.$$

 \Rightarrow Average error of best approximation in \mathbb{V}_{rb} over \mathbb{P}_{h}^{h} .

The POD space Web is the Mathemainal subspace of Wothat minimises the above of romaging of the space of the state of the s

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Basis by POD approach

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Proper Orthogonal Projection (POD):

1. Compute the solution $u_{\delta}(\mu)$ for all $\mu \in \mathbb{P}_h$ and define the correlation matrix

$$C_{ij} = (u(\mu_j), u(\mu_i))_{\mathbb{V}}, \qquad i, j = 1, \dots, M.$$

- 2. Find the eigen-pairs (λ_n, v_n) solution to $Cv_n = \lambda_n v_n$ for the N largest eigenvalues.
- **3.** Define basis functions as

$$\varphi_n = \sum_{i=1}^M (v_n)_i \, u(\mu_i).$$

 \Rightarrow Basis functions are linear combinations of snapshots.

4. Set $\mathbb{V}_{POD} = \operatorname{span}\{\varphi_1, \ldots, \varphi_N\}.$

Basis by POD approach



Proper Orthogonal Projection (POD):

1. Compute the solution $u_{\delta}(\mu)$ for all $\mu \in \mathbb{P}_h$ and define the correlation matrix in $u_{\delta}(\mu)$

$$C_{ij} = (u (\mu_{ij}), u(\mu_{ij}), u(\mu_{ij})) = 1, \dots, M.$$

- 2. Find the independence of the set of the
- 3. Defize Basisnenhasisn supportions as

$$\varphi_n = \sum_{i=1}^{M} (\varphi_n)_{\overline{\tau}} u \underbrace{M}_{i=1} (v_n)_i u($$

 $\begin{array}{l} \Rightarrow & \text{Basis function}\\ \mu_i \\ \mu_i \\ \text{finear combination}\\ \text{snapshots, snapshots.} \end{array}$

4. Set $\underline{\mathbb{Y}}_{\text{POD}} = \underline{\mathbb{Y}}_{\text{POD}} \{ \varphi_1 \text{ span} \{ \varphi_N \} \dots, \varphi_N \}.$

Result:

$$\frac{1}{M} \sum_{i=1}^{L} \inf_{v \in \mathbb{V}_{POD}} \|u(\mu_i) - v\|_{\mathbb{V}}^2 = \sum_{n=N+1}^{M} \lambda_n.$$

In practise: replace $u(\mu_i)$ by a truth approximation $u_{\delta}(\mu_i)$.





Find eigen-decomposition of C



The reduced model is now obtained as

$$A_h u_{\delta} = f_h \quad \Rightarrow \\ (V^T A_h V) V^T u_{\delta} = V^T f_h \qquad V^T V = I$$

or

$$A_{rb}u_{rb} = f_{rb}$$



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$$s(u) \simeq s(u_{\delta}) \simeq s(\mathbf{V}u_{rb})$$



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$$s(u) \simeq s(u_{\delta}) \simeq s(\operatorname{V} u_{rb})$$

Since $N \ll \mathcal{N}$ we have the potential for speed

POD example - Ex I



- **o** Nodal values of exact solutions used instead of FE-approximations.
- o \mathbb{P}_h : 491 equidistant points in $\mathbb{P} = [0.01, 0.5]$.





Method has several names -

Karhunen-Loeve expansions
Proper orthogonal expansions
Empirical eigenfunctions

Properties -

Simple and straightforward for linear systems
Offline cost can be high
Accuracy ? — did we sample carefully enough ?
What about online cost for nonlinear problem

 $abla^2 u(x,\mu) = f(u,\mu) \Rightarrow A_{rb}(\mu)u_{rb}(x,\mu) = V^T f(V u_{rb},\mu)$ **Depends on** \mathcal{N}



We need to develop methods that address these shortcomings

Compute what we need - nothing more
Control the error to certify results
Ensure efficiency
Deal with non-linear problems

This will be the main topics of Lecture 2

Questions ?