Combining Logics

Walter Carnielli

IFCH/CLE - State University of Campinas - Brazil ICR-CSC - University of Luxembourg - Luxembourg

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Walter Carnielli (Unicamp and Uni.Lu)

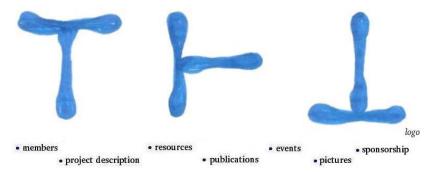
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- Do we really reason using propositional, quantified, epistemic, alethic, doxastic, temporal, many-valued, fuzzy, intuitionistic, paraconsistent · · · logics?
- Or we do combine everything, and perhaps more?
- How is really the reasoning in domains like legal reasoning, computer systems, economic reasoning, etc, expressed in terms of elementary concepts?

Project ConsRel - CLE at Campinas

ConsRel: Logical Consequence and Combinations of Logics



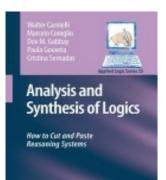
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- If we can combine reasoning, or at least combine logics, why not decompose them?
- If a logic is decomposed into "elementary" sublogics, is it possible to recover it by combining such fragments?
- What kind of properties of logics (like completeness, decidability, interpolation properties, axiomatizability, computable efficiency, etc.) can be transferred to their combinations?

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A consequence:

General methods for combining logics, lots of examples and some suggested applications.



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Walter Carnielli (Unicamp and Uni.Lu)

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"For as this ought, or ought not, expresses some new relation or affirmation, it is necessary that it should be observed and explained; and at the same time that a reason should be given, for what seems altogether inconceivable, how this new relation can be a deduction from others, which are entirely different from it. But as authors do not commonly use this precaution..."

 Hume was asking for something we may call "bridge principle", without which we seem to be unable to handle combined reasoning...

- expressed in "A Treatise of Human Nature" (Book 3, Part 1, Section 1, paragraph 27)...
- ... generated a controversy about the legitimacy of statements that bind factualities to norms
- ... and inaugurated the idea of "bridge principles" as necessary principles for mixed reasoning.

- spontaneous or hidden bridge principles pose intriguing questions to combined logics;
- bridge principles may spontaneously arise in the operation of combining logics ...
- they may have however, desirable or undesirable consequences for combined reasoning;
- moreover, we also find collapsing and anti-collapsing problems.

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- In order to perform such a jump from 'is' to 'ought', one might appeal to an explicit "bridge principle", which specifically connects 'is' and 'ought';
- $\alpha \rightarrow \bigcirc \alpha$ is a simple bridge principle representing 'is-ought';
- Are bridge principles *necessary*?

- Bridge principles may not be necessarily analytical, in the sense that they might not be true because of the meaning of their symbols alone;
- Yet, bridge principles in a broad sense may appear spontaneously when combining logics;
- How can something non-analytical appear analytically?

Definition

G. Schurz: An axiom schema *A* is a bridge principle iff *A* contains at least one schematic letter which has at least one occurrence within the scope of a deontic "obligation" operator \bigcirc , and at least one occurrence outside the scope of any \bigcirc .

- 'Ought-implies-can': $\bigcirc \alpha \rightarrow \Diamond \alpha$;
- But this can be widely extended beyond modal logics.

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Splicing (combining) versus splitting (decomposing) logics

- Most relevant methods: fusion, product of modal logics and fibring.
- Paradigmatic splicing method: algebraic fribring.
- Used in computer science and knowledge representation; used less by logicians, and very timidly by philosophers.
- Integrating several reasoning modules: temporal, epistemic, alethic. and more...
- Paradigmatic splitting method: possible-translations semantics.

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- Is there a unique "correct" logic (monism), or many (pluralism), or none (instrumentalism)?
- Does composition of logics restore the unity from "fragments", or create more specimens, expanding the "pluralism"?
- Philosophers of logic should take combined logics into account!

- Introduced by R. Thomason in 1984 (but anticipated, in examples of fusing alethic and deontic modalities, by M. Fitting in 1968);
- The fusion of L₁ with L₂ is the bimodal logic L, defined over a language with two boxes. The rest of the connectives are assumed to be classical, and so they are shared by L₁ and L₂.

Definition

- Semantic fusion of \mathcal{L}_1 (with \Box_1) and \mathcal{L}_2 (with \Box_2): bimodal \mathcal{L} with \Box_1 and \Box_2 characterized by general Kripke frames $\langle W, R_1, R_2 \rangle$ with a set of worlds W and two relations R_1 and R_2 over W.
- 2 $\langle W, R_1, R_2 \rangle$ is such that $\langle W, R_1 \rangle$ and $\langle W, R_2 \rangle$ are Kripke frames for \mathcal{L}_1 and \mathcal{L}_2 .

 The Hilbert calculi of L is the merging of the axioms and rules of both logics (but in L they can be instantiated with mixed formulas).

 Introduced by K. Segerberg in 1973 and by V. Shehtman in 1978 (in two papers with the same title...).

Definition

Product of \mathcal{L}_1 (with \Box_1) and \mathcal{L}_2 (with \Box_2) is also a bimodal logic \mathcal{L} with \Box_1 and \Box_2 is \mathcal{L} , characterized by all Kripke models $\langle W_1 \times W_2, \overline{R}_1, \overline{R}_2, V_1 \times V_2 \rangle$

Definition

- $\overline{R}_i \subseteq (W_1 \times W_2) \times (W_1 \times W_2)$ is defined from R_i as:
 - $(w_1, w_2)\overline{R}_1(u_1, u_2)$ iff $w_1R_1u_1$ and $w_2 = u_2$;

• $(w_1, w_2)\overline{R}_2(u_1, u_2)$ iff $w_2R_2u_2$ and $w_1 = u_1$.

2
$$V_1 \times V_2 : \mathbb{P} \longrightarrow \wp(W_1 \times W_2)$$
 is the mapping $(V_1 \times V_2)(p) = V_1(p) \times V_2(p)$, such that $V_i : \mathbb{P} \longrightarrow \wp(W_i)$ is a valuation in $\langle W_i, R_i, V_i \rangle$

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Fibring modal logics: putting yourself in somebody else's shoes

• D. Gabbay, 1996; also generates bi-modal logics.

Definition

- Given L₁ and L₂ and their Kripke models, take transfer maps: h₁ from worlds of models M₁ of L₁ into models M₂ of L₂, and h₂ vice-versa.
- A Kripke model of L₁ evaluates □₂φ at the actual world w₁ by transferring the validity checking to checking □₂φ within the Kripke model h₁(w₁) at its actual world.
- Solution Vice-versa for $\Box_1 \varphi$ within a Kripke model for \mathcal{L}_2 .

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- Designed to overcome the limitations of fusion, product and fribing (all them for modal logics only).
- Proposed by A. Sernadas, C. Sernadas and C. Caleiro in 1999.

Definition

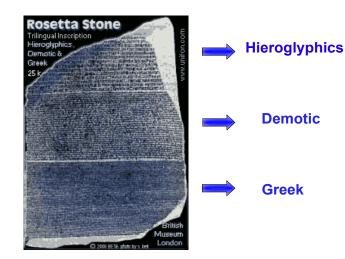
- The categorial fibring of L₁ and L₂ is the *least* logic L over the combined language which extends L₁ and L₂.
- 2 It is the *coproduct* of \mathcal{L}_1 and \mathcal{L}_2 in the category of logics and their morphisms.

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- Categorial fibring is universal in the sense of category theory, and generalize fusion and fibring;
- Metafibring, a restriction proposed by M. Coniglio in 2005, a categorial construction where morphisms preserve meta-properties of the logics.
- Metafibring permits a logic to be recovered from its fragments (that is, from logics defined over sub-languages).

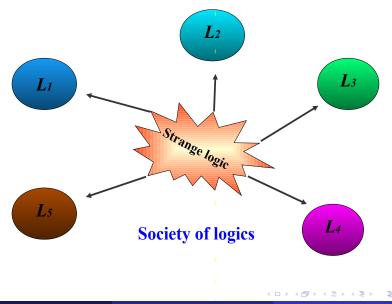
- Idealized for paraconsistent logics, specially for Logics of Formal Inconsistency (LFIs).
- However, they are applicable in several other cases.
- They constitute a most general method for decomposing (splitting) logics.

Rosetta stone: how translations work



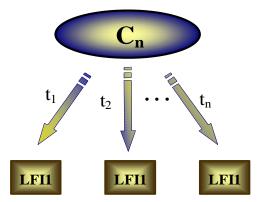
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Translations acting together



Walter Carnielli (Unicamp and Uni.Lu)

Possible-translations semantics for LFIs



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- "Bridge principles" wide sense: interactions (i.e., derivations) among distinct logic operators which are not instances of valid derivations in the logics being combined.
- So, e.g., in the logic *L* obtained by combining ∧ with ∨ via metafibring:
- *p* ∧ *r* ⊢ (*p* ∧ *r*) ∨ *q* is not a bridge principle, as it is derived by substitution from *p* ⊢ *p* ∨ *q*, valid in the logic of ∨.

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- However, p ∧ (q ∨ r) ⊢ (p ∧ q) ∨ r (distributivity of ∧ over ∨) is not obtained in the logic of ∧, nor in the logic of ∨, but appears spontaneously in the combination! (Béziau & Coniglio).
- Another case of spontaneous emergence of a bridge principle: in the metafibring of the logic of classical ¬ and the logic of classical ∨, the law of excluded-middle p ∨ ¬p emerges unavoidably in the combined logic (Coniglio).

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What is the meaning of spontaneous laws?

- Also, in the metafibring of the logics of classical ¬ and classical →, the Principle of Pseudo-Scotus p → (¬p → q) emerges unavoidably (Coniglio).
- The first case obtains the full \lor - \land fragment of **PC**.
- The second and third cases obtain full PC.
- In all cases, the bridge principles arise spontaneously due to the nature of the combination process.
- Are they expected? No, from the point of view of intuitionists or paraconsistentists!

- Self-generated bridge principles in the product of two normal modal logics L₁ and L₂ (whose languages have, respectively, □₁ and ◊₁, and □₂ and ◊₂) (Gabbay):
 - $(\Box_1 \Box_2 \alpha \leftrightarrow \Box_2 \Box_1 \alpha)$ \Box -commutativity;
 - $(\Diamond_1 \Diamond_2 \alpha \leftrightarrow \Diamond_2 \Diamond_1 \alpha)$ \Diamond -commutativity;
 - $(\Diamond_1 \Box_2 \alpha \rightarrow \Box_2 \Diamond_1 \alpha)$ (1,2)-Church-Rosser;
 - $(\Diamond_2 \Box_1 \alpha \to \Box_1 \Diamond_2 \alpha)$ (2,1)-Church-Rosser.

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- Consequently, other bridge principles will be also derivable:
 - $(\Diamond_1^k \Box_2^m \alpha \to \Box_2^m \Diamond_1^k \alpha)$
 - $(\Diamond_2^k \Box_1^m \alpha \to \Box_1^m \Diamond_2^k \alpha)$

 $(1^k, 2^m)$ -Church-Rosser property;

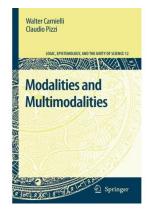
 $(2^k, 1^m)$ -Church-Rosser property.

- A form of (1,2)-Church-Rosser property for "knowledge" K:
 ◊Kα → K◊α will emerge spontaneously in the product of any normal modal logics.
- We may inform a previously ignorant person about the following fact *p*: "there exists an egg-laying mammal" (namely, the platypus).
- So, ◊*Kp* is true (as we may inform her), but *K*◊*p* is false (indeed, her ignorance excludes *a priori* the possibility of her knowing about the existence of such an animal).

- Interactions are only revealed after a careful semantic analysis, and one has no control on which bridge principles might crop up.
- A side effect: it is not possible to obtain *a priori* a complete Hilbert calculus for products of modal logics.
- Additional bridge principles might have to be explicitly added to ensure completeness.
- Within multimodalities a profusion of bridge principles naturally appears.

- The collapsing problem (D. Gabbay, and independently L. Fariñas del Cerro, A. Herzig, 1996).
- By freely combining **PC** and intuitionistic propositional logic the resulting logic collapses to classical logic: intuitionistic implication becomes classic.
- The collapsing phenomenon–a spontaneous and undesirable bridge principles: α₁ →_c α₂ ⊢ α₁ →_i α₂ and α₁ →_i α₂ ⊢ α₁ →_c α₂ (where →_c and →_i are, respectively, **PC** and **HI** implication).

Modalities and Multimodalities a newborn consequence



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