5 Discovering Domain Entities

Section 2 briefly characterised, informally and also a bit more formally, what we mean by a domain. Section 3 informally and systematically characterised the four categories of entities: parts, actions, events and behaviours. Section 4 more-or-less “repeated” Sect. 3’s material but by now giving more terse narratives (that is, informal descriptions) and, for the first time, also formalisations. Section 4 did not hint at how one discovers domain parts (i.e., their types), actions, events and behaviours. In this section we try unravel a set of techniques and tools — so-called ‘discoverers’ and ‘analysers’ — using which the domain describer (scientist and/or engineer) can more-or-less systematically discover, analyse and describe a domain, informally and formally.

5.1 Preliminaries

Before we present the discoverers and analysers we need establish some concepts.

5.1.1 Part Signatures

Let us consider a part \( p : P \). Let \( p : P \), by definition, be the principal part of a domain. Now we need to identify (i) the type, \( P \), of that part; (ii) the types, \( S_1, \ldots, S_m \), of its proper sub-parts (if \( p \) is composite); (iii) the type, \( PI \), of its unique identifier; (iv) the possible types, \( MI_1, \ldots, MI_n \), of its mereology; and (v) the types, \( A_1, \ldots, A_o \), of its attributes. We shall name that cluster of type identifications the part signature. We refer to \( P \) as identifying the part signature. Each of the \( S_i \) (for \( i \in \{1..m\} \)) identifies sub-parts and hence sub-part, i.e., part signatures.

Example 41 (Net Domain and Sub-domain Part Signatures) The part signature of the hubs and the links are here chosen to be those of(i) the (root) net type, \( N \), (ii) the (sub-domain) set of hubs type \( Hs \), (ii) the (sub-domain) set of links type \( Ls \) and (v) the type of net attributes \( \text{Net\_name}, \text{Net\_owner} \), etc. The part signatures \( Hs \) and \( Ls \) are (ii) \( Hs = H\text{-set}, H \) (iii,iv) \( Hl, Ll\text{-set} \) (v) \( \text{Hub\_Nm}, \text{Location}, H\Sigma, H\Omega, \ldots \) (ii) \( Ls = L\text{-set}, L \) (iii,iv) \( Ll, Hl\text{-set} \) (v) \( \text{Link\_Nm}, L\Sigma, L\Omega, \text{LEN} \), etc.

5.1.2 Domain Indices

By a domain index we mean a list of part type names that identify a sequence of part signatures. More specifically The domain \( \Delta \) has index \( \langle \Delta \rangle \). The sub-domains of \( \Delta \), with part types \( A, B, \ldots, C \), has indices \( \langle \Delta, A \rangle, \langle \Delta, B \rangle, \ldots, \langle \Delta, C \rangle \). The sub-domains of sub-domain with index \( \ell \) and with part types \( A, B, \ldots, C \) has indices \( \ell^{-}(A), \ell^{-}(B), \ldots, \ell^{-}(C) \).

Example 42 (Indices of a Road Pricing Domain) We refer to the the Road-pricing Transport Domain, cf. Example 36 on page 56.
The sub-domain indices of the road-pricing transport domain, $\Delta$, are: $\langle \Delta \rangle$, $\langle \Delta, N \rangle$, $\langle \Delta, F \rangle$, $\langle \Delta, M \rangle$, $\langle \Delta, N, Vs \rangle$, $\langle \Delta, N, H \rangle$, $\langle \Delta, N, L \rangle$ and $\langle \Delta, F, V \rangle$.

5.1.3 Inherited Domain Signatures

Let $\langle \Delta, A, B, C, D \rangle$ be some domain index. Then $\langle \Delta, A, B, C \rangle \langle \Delta, A, B \rangle \langle \Delta, A \rangle \langle \Delta \rangle$ are the inherited domain indices of $\langle \Delta, A, B, C, D \rangle$.

5.1.4 Domain and Sub-domain Categories

By the domain category of the domain indexed by $\ell\langle D \rangle$ we shall mean the domain signature of $D$, and the action, event and behaviour definitions whose signatures involves just the types given in the domain signature of $D$ or in inherited domain signatures.

Example 43 (The Road-pricing Domain Category) The road-pricing domain category consist of the types $N$, $F$ and $M$, the create Net create Fleet and create M actions, and corresponding Net, Fleet and M behaviours.

By a sub-domain category, of index $\ell$, we shall mean the sub-domain types of the sub-domain designated by index $\ell$, and the actions, events and behaviours whose signatures involves just the types of the $\ell$ indexed sub-domain or of any prefix of $\ell$ indexed sub-domain or of the root domain.

Example 44 (A Hub Category of a Road-pricing Transport Domain) The ancestor sub-domain types of the hub sub-domain are: $HS$, $N$ and $\Delta$. The hub category thus includes the part (etc.) types $H$, HI, ..., the insert Hub and the delete Hub actions, perhaps some saturated_hub (and/or other) event(s), but probably no hub behaviour as it would involve at least the type LI which is not in an ancestor sub-domain of the Hub sub-domain.

5.1.5 Simple and Compound Indexes

By a simple index we mean a domain or a sub-domain index. By a compound index we mean a set of two or more distinct indices of a domain $\Delta$. Compound indices, $c_{idx} : \{\ell_1, \ell_2, \ldots, \ell_k\}$, designate parts, actions, events and behaviours each of whose types and signatures involve types defined by all of the simple indexes of $c_{idx}$.

Example 45 (Compound Indices of the Road-pricing System) We show just one compound index: $\{\langle \Delta, N, HS, H \rangle, \langle \Delta, N, LS, L \rangle\}$. 
5.1.6 Simple and Compound Domain Categories

By a simple domain category we shall mean any $\ell$-indexed [sub-]domain category. By the compound domain category of compound index $c_{idx} : \{\ell_i, \ell_j, \ldots, \ell_k\}$, we shall mean the set of types, actions, events and behaviours as induced by compound index $c_{idx}$, that is, parts, actions, events and behaviours each of whose types and signatures involve types defined by all of the simple indexes of $c_{idx}$.

**Example 46 (The Compound Domain Category of Hubs and Links)** The compound domain category designated by \{\langle $\Delta$, N, HS, H \rangle, \langle $\Delta$, N, LS, L \rangle\} includes:

\[
\begin{align*}
\text{type} & \\
HIs &= HI\text{-set} & \text{mereo}_L: L \rightarrow HIs \\
\text{axiom } & \forall \ his:HIs\cdot \text{card } his=2 & \text{mereo}_H: H \rightarrow LIs \\
LIs &= LI\text{-set} & \text{attr}_H\Sigma: H \rightarrow H\Sigma \\
H\Sigma &= (LI \times LI)\text{-set} & \text{attr}_L\Sigma: L \rightarrow L\Sigma \\
L\Sigma &= (HI \times HI)\text{-set} & \text{attr}_L\Omega: L \rightarrow L\Omega \\
H\Omega &= H\Sigma\text{-set} & \text{attr}_H\Omega: H \rightarrow H\Omega \\
L\Omega &= L\Sigma\text{-set} & \forall h:H\cdot \text{attr}_H\Sigma(h) \subseteq \text{attr}_H\Omega(h) \\
\text{value} & \\
\end{align*}
\]

5.1.7 Examples

We repeat some examples, but now “formalised”.

**Example 47 (The Root Domain Category)** We start at the root, $\Delta$, of the Road Pricing Domain. See Fig. 13.

\[
\Delta
\]

\[
<\Delta>
\]

Figure 13: The $\langle \Delta \rangle$ Root

At the root we ‘discover’ the net, fleet and road pricing monitor. See Fig. 14 on the facing page.

When observing the very essence of the road pricing domain “at the $\langle \Delta \rangle$ level” one observes:
New indices: \{<\Delta,N>,<\Delta,F>,<\Delta,M>\}

Figure 14: Exploring the root index \(\langle \Delta \rangle\) Index

type
N, F, M
value
\begin{align*}
\text{obs}_N &: \Delta \rightarrow N \\
\text{obs}_F &: \Delta \rightarrow F \\
\text{obs}_M &: \Delta \rightarrow M \\
\text{attr}_{\ldots} &: \Delta \rightarrow \ldots
\end{align*}

where \(\ldots\) stands for types of road pricing domain attributes.

Example 48 (The Net Domain Category) We then proceed to explore the domain at index \(\langle \Delta,N \rangle\). See Fig. 15.

New indices: \{<\Delta,N,Hs>,<\Delta,N,Ls>\}

Figure 15: Exploring the \(\langle \Delta,N \rangle\) Index

When observing the very essence of the Net domain, “at the \(\langle \Delta,N \rangle\) level” one observes:

type
\begin{align*}
\text{Hs} &= \text{H-set} \\
\text{Ls} &= \text{L-set} \\
\text{H} &= \ldots \\
\text{L} &= \ldots
\end{align*}

value
\begin{align*}
\text{obs}_{\text{Hs}} &: N \rightarrow \text{Hs} \\
\text{obs}_{\text{Ls}} &: N \rightarrow \text{Ls} \\
\text{attr}_{\text{Hs}} &: \text{Hs} \rightarrow \ldots \\
\text{attr}_{\text{Ls}} &: \text{Ls} \rightarrow \ldots
\end{align*}
where ... stand for attributes of the Hs and the Ls parts of N.

Example 49 (The Fleet Domain Category) We then proceed to explore the domain at index \( \langle \Delta, F \rangle \). See Fig. 16.

![Figure 16: Exploring the \( \langle \Delta, F \rangle \) Index](image)

When observing the very essence of the Fleet domain, “at the \( \langle \Delta, F \rangle \) level” one observes:

```
type
  Vs = V-set
  V ...
value
  obs\_Vs: F \to Vs
  attr\_... Vs \to ...
```

where ... stand for attributes that we may wish to associate with Fleets of vehicles.

Example 50 (The Hub Domain Category) We now switch “back” to explore the domain at index \( \langle \Delta, N, Hs \rangle \). See Fig. 17 on the next page.

When observing the very essence of the Fleet domain, “at the \( \langle \Delta, N, Hs, H \rangle \) level” one observes:

```
type
  HI ...
value
  uid\_HI: H \to HI
  attr\_...: H \to ...
```

where ... stand for LOCation, etc.
Example 51 (The Link Domain Category) Next we explore the link domain. See Fig. 18.

When observing the very essence of the Fleet domain, “at the $\langle \Delta, N, Hs, H \rangle$ level” one observes:

```
<table>
<thead>
<tr>
<th>type</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>LI</td>
<td>uid_{LI} : L \rightarrow LI</td>
</tr>
<tr>
<td></td>
<td>attr_{...} : L \rightarrow ...</td>
</tr>
</tbody>
</table>
```

where ... stand for LOCation, LENgth, etc.

Example 52 (The Compound Hub and Link Domain Category) We next explore a compound domain. See Fig. 19 on the next page.

When observing the very essence of the Fleet domain, at the $\{\langle \Delta, N, Hs, H \rangle, \langle \Delta, N, Ls, L \rangle\}$ level one observes:
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**Figure 19:** Exploring composite index \(\{\langle \Delta, N, Hs, H \rangle, \langle \Delta, N, Ls, L \rangle\}\)

**type**

\[H\Sigma = (LI \times LI)\text{-set}, \ H\Omega = H\Sigma\text{-set},\]
\[L\Sigma = (HI \times HI)\text{-set}, \ L\Omega = L\Sigma\text{-set}\]

**value**

\[\text{attr}_{H\Sigma}: H \to H\Sigma, \ \text{attr}_{H\Omega}: H \to H\Omega\]
\[\text{attr}_{L\Sigma}: L \to L\Sigma, \ \text{attr}_{L\Omega}: L \to L\Omega\]
\[\text{mero}_L: L \to HI\text{-set axiom } \forall l:L:\text{card mero}_L(l) = 2\]
\[\text{mero}_H: H \to LI\text{-set } (= LI\text{-set})\]
\[\text{remove}_H: HI \to N \not\sim N\]
\[\text{insert}_L: L \to N \sim N\]
\[\text{remove}_L: LI \to N \not\sim N\]

\[\ldots\]

**axiom**

\[\forall h:\sigma:H\Sigma \ldots, \ \forall h:\omega:H\Omega \ldots\]
\[\forall l:\sigma:H\Sigma \ldots, \ \forall l:\omega:H\Omega \ldots\]

**5.1.8 Discussion**

The previous examples (47–52), especially the last one (52), illustrates the complexity of a domain category; from just observing sub-part types and attributes (as in Examples 47–49), Example 52 observations grow to intricate mereologies etcetera. The ‘discoverers’ that we shall propose aim at structuring the discovery process by focusing, in turn, on part sorts, concrete part types, unique identifier types of parts, part mereology, part attributes, action signatures, event signatures, etc.
5.2 Proposed Type and Signature ‘Discoverers’

By a ‘domain discoverer’ we shall understand a tool and a set of principles and techniques for using this tool in the discovery of the entities of a domain.

In this section we shall put forward a set of type and signature discoverers. Each discoverer is indexed by a simple or a compound domain index. And each discoverer is dedicated to some aspect of some entities. Together the proposed discoverers should cover the most salient aspects of domains. Our presentation of type and signature discoverer does not claim to help analyse “all” of a domain.

We need formally define what an index is.

type

\[
\text{Index} = \text{Smpl}\_\text{Idx} \mid \text{Cmpd}\_\text{Idx} \\
\text{Smpl}\_\text{Idx} = \{ \langle \Delta \rangle \hat{\text{id}}x \mid \text{id}x: \text{Type\_Name}^* \} \\
\text{Cmpd}\_\text{Idx}' = \text{Smpl}\_\text{Idx}\text{-set} \\
\text{Cmpd}\_\text{Idx} = \{ \text{sis}\_\text{Cmpd}\_\text{Idx}' \cdot \text{wf}\_\text{Cmpd}\_\text{Idx}(\text{sis}) \mid \} }
\]

value

\[
\text{wf}\_\text{Cmpd}\_\text{Idx}: \text{Cmpd}\_\text{Idx}' \rightarrow \text{Bool} \\
\text{wf}\_\text{Cmpd}\_\text{Idx}(\text{sis}) \equiv \forall \text{si},\text{si}':\text{Smpl}\_\text{Idx} \cdot \{\text{si},\text{si}'\} \subseteq \text{sis} \wedge \text{si} \neq \text{si}'
\]

\[
\text{DISCOVERER\_KIND}: \text{Index} \rightarrow \text{Text} \\
\text{DISCOVERER\_KIND}(\ell {\Uparrow} (t)) \text{ as text} \\
\text{pre}: \ell {\Uparrow} (t) \text{ is a valid index beginning with } \Delta \\
\text{post}: \text{text is some, in our case, RSL text}
\]

The idea of the \( \ell {\Uparrow} (t) \) index is that it identifies a sub-domain, \( t \), of \( \Delta \) where \( \text{DISCOVERER\_KIND} \) is any of the several different “kinds” of domain forms:

[90 (Page 76)] PART\_SORTS,
[91 (Page 77)] HAS\_A\_CONCRETE\_TYPE,
[92 (Page 78)] PART\_TYPES,
[93 (Page 79)] UNIQUE\_ID,
[96 (Page 79)] MEREOLGY,
[98 (Page 80)] ATTRIBUTES,
[100 (Page 82)] ACTION\_SIGNATURES,
[101 (Page 83)] EVENT\_SIGNATURES and
[103 (Page 84)] BEHAVIOUR\_SIGNATURES.

In a domain analysis (i.e., discovery) the domain description emerges “bit-by-bit”. Initially types are discovered and hence texts which define unique identifier types and functions, mereology types and functions, and attribute types and functions. Then the signatures of actions, events and behaviours.
You may consider these “piece-wise” texts as being “added” to a (hence) growing reservoir of (RSL) texts with this reservoir being continually inspected by the domain analyser.

5.2.1 Analysing Domain Parts

The two most important aspects of an algebra are those of its parts and its operations. Rather than identifying, that is, discovering or analysing individual parts we focus on discovering their types — initially by defining these as sorts. And rather than focusing on defining what the operations achieve we concentrate on the signature, i.e., the types of the operations.

It (therefore) seems wise to start with the discovery of parts, and hence of their types. Part types are present in the signatures of all actions, events and behaviours. When observing part types we also observe a variety of part type analysers: possible unique identities of parts, the possible mereologies of composite parts, and the types of the attributes of these parts.

Domain Part Sorts and Their Observers Initially we “discover” parts — by deciding upon their types, in the form, first of sorts, subsequently and possibly in the form of concrete types.

A Domain Sort Discoverer:

90. A part type discoverer applies to a simply indexed domain, index, and yields

- a set of type names
- each paired with a part (sort) observer.

Example 53 (Some Part Sort Discoveries) We apply a concrete version of the above sort discoverer to the road-pricing transport domain $\Delta$:

\[
\text{PART\_SORTS}(\langle \Delta \rangle):
\]

- type
  - N, F, M
- value
  - obs$_N$: $\Delta \rightarrow N$
  - obs$_F$: $\Delta \rightarrow F$
obs\(_M\): \(\Delta \rightarrow M\)

\[
\text{PART\_SORTS}(⟨\Delta, N⟩):
\]

\[
\begin{align*}
\text{type} & : Hs, Ls \\
\text{value} & :
\begin{align*}
\text{obs\_Hs}: & N \rightarrow Hs \\
\text{obs\_Ls}: & N \rightarrow Ls
\end{align*}
\]

\[
\text{PART\_SORTS}(⟨\Delta, F⟩):
\]

\[
\begin{align*}
\text{type} & : Vs \\
\text{value} & :
\begin{align*}
\text{obs\_Cs}: & F \rightarrow Vs
\end{align*}
\]

**Domain Part Types and Their Observers**

**Do a Sort Have a Concrete Type?** Sometimes we find it expedient to endow a “discovered” sort with a concrete type expression, that is, “turn” a sort definition into a concrete type definition.

91. Thus we introduce the “discoverer”:

91 \(\text{HAS\_A\_CONCRETE\_TYPE}: \text{Index} \rightarrow \text{Bool}\)

91 \(\text{HAS\_A\_CONCRETE\_TYPE}(\ell\gamma(t)):\text{true}|\text{false}\)

**Example 54 (Some Type Definition Discoveries)** We exemplify two true expressions:

\[
\begin{align*}
\text{HAS\_A\_CONCRETE\_TYPE}(⟨\Delta, N, Hs⟩) \\
\text{HAS\_A\_CONCRETE\_TYPE}(⟨\Delta, N, Ls⟩) \\
\sim \text{HAS\_A\_CONCRETE\_TYPE}(⟨\Delta, N, Hs, H⟩) \\
\sim \text{HAS\_A\_CONCRETE\_TYPE}(⟨\Delta, N, Ls, L⟩)
\end{align*}
\]
A Domain Part Type Observer: The PARTTYPES invocation yields one or more sort definitions of part types together with their observer functions. The domain analyser can decide that some parts can be immediately analysed into concrete types. Thus, together with yielding a type name, the PARTTYPES can be expected to yield also a type definition, that is, a type expression (paired with the type name). Not all type expressions make sense. We suggest that only some make sense.

92. The PARTTYPES discoverer applies to a composite type, t, and yields
   a a type definition, T = TE;
   b together with the sort and/or type definitions of so far undefined type names of TE.
   c The PARTTYPES discoverer is not defined if the designated sort is judged to not warrant a concrete type definition.

Example 55 (Some Part Type Discoveries) We exemplify two discoveries:

PARTTYPES(⟨Δ,N,Hs⟩):
   type
   H
   Hs = H-set

PARTTYPES(⟨Δ,N,Ls⟩):
   type
   L
   Ls = L-set

PARTTYPES(⟨Δ,F⟩):
   type
   V
   Vs = V-set
Concrete Part Types: In Example 55 on the facing page we illustrated one kind of concrete part type: sets. Practice shows that sorts often can be analysed into sets. Other analyses of part sorts are Cartesians, list, and simple maps:

90b. \( te: tn_1 \times tn_2 \times \ldots \times tn_m \)
90b. \( te: tn^* \)
90b. \( te: \text{Token } \frac{m}{m} \ tn \)

where \( tn \)'s are part type – usually sort – names some of which may have already been defined, and where \( \text{Token} \) is some simple atomic (non-part) type.

Part Type Analysers There are three kinds of analysers: unique identity analysers, mereology analysers and general attribute analysers.

Unique Identity Analysers: We associate with every part type \( T \), a unique identity \( TI \).

93. So, for every part type \( T \) we postulate a unique identity analyser function \( \text{uid}_{TI} \).

value
93. \( \text{UNIQUE}_\text{ID}: \text{Index} \rightarrow \text{Text} \)
93. \( \text{UNIQUE}_\text{ID}(\ell^\langle T \rangle) \):
93. \( \text{type} \)
93. \( TI \)
93. \( \text{value} \)
93. \( \text{uid}_{TI}: T \rightarrow TI \)

Mereology Analysers: We remind the reader of Sects. 3.1.6 on page 32. Given a part, \( p \), of type \( T \), the mereology, \( \text{MERELOGY} \), of that part is the set of all the unique identifiers of the other parts to which part \( p \) is partship-related as “revealed” by the \( \text{mero}_{TI} \) functions applied to \( p \).

94. Let types \( T_1, T_2, \ldots, T_n \) be the types of all parts of a domain.
95. Let types \( TI_1, TI_2, \ldots, TI_n \), be the types of the unique identifiers of all parts of that domain.
96. The mereology analyser \( \text{MERELOGY} \) is a generic function which applies to an index and yields the set of

\[ a \text{ zero,} \]

\[ ^{29}\text{We here assume that all parts have unique identifications.} \]
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b one or
c more

mereology observers.

\textbf{type}
94. \( T = T_1 \mid T_2 \mid \ldots \mid T_n \)
95. \( T_{\text{idx}} = T_1 T_2 \mid \ldots \mid T_n \)
96. \textbf{MERELOGY: Index \(\rightarrow\) Text}
96. \textbf{MERELOGY}(\{\ell_i \sim \langle T_j \rangle, \ldots, \ell_k \sim \langle T_l \rangle\}):
96a. either: \{\}
96b. or: \text{mereo}_{T_i}: T \rightarrow (T_{\text{I}_i}[T_{\text{I}_i}-\text{set}])
96c. or: \{ \text{mereo}_{T_i}: T \rightarrow (T_{\text{I}_i}[T_{\text{I}_i}-\text{set}]),
96c. \text{mereo}_{T_j}: T \rightarrow (T_{\text{I}_j}[T_{\text{I}_j}-\text{set}]),
96c. \ldots,
96c. \text{mereo}_{T_z}: T \rightarrow (T_{\text{I}_z}[T_{\text{I}_z}-\text{set}]) \}

where none of \( T_{\text{I}_x}, T_{\text{I}_y}, \ldots, T_{\text{I}_z} \) are equal to \( T \) and each is some \( T_{\text{idx}} \).

\textbf{General Attribute Analysers:} A general attribute analyser analyses parts beyond their unique identities and possible mereologies.

97. Part attributes have names. We consider these names to also abstractly name the corresponding attribute types, that is, the names function both as attribute names and sort names. Finally we allow attributes of two or more otherwise distinct part types to be the same.

98. \textbf{ATTRIBUTES} applies to parts of any part type \( t \) and yields

99. the set of attribute observer functions \text{attr}_{at}, one for each attribute sort \( at \) of \( t \).

\textbf{type}
97. \( AT = AT_1 \mid AT_2 \mid \ldots \mid AT_n \)
\textbf{value}
98. \textbf{ATTRIBUTES: Index \(\rightarrow\) Text}
98. \textbf{ATTRIBUTES}(\ell \sim (T)):
99. \textbf{type}
99. \( AT_1, AT_2, \ldots, AT_m \)
99. \textbf{value}
99. \text{attr}_{AT_1}: T \rightarrow AT_1
99. \text{attr}_{AT_2}: T \rightarrow AT_2
99. \ldots,
99. \text{attr}_{AT_m}: T \rightarrow AT_m, m \leq n
Example 56 (Example Part Attributes) We exemplify attributes of composite and of atomic parts:

\[
\text{ATTRIBUTES}(\langle \Delta \rangle):
\]

\[
\text{type}
\]

\[
\text{Domain\_Name, ...}
\]

\[
\text{value}
\]

\[
\text{attr\_Name: } \Delta \to \text{Domain\_Name}
\]

\[
\text{...}
\]

where \text{Domain\_Name} could include \textit{State Roads} or \textit{Rail Net}, etcetera.

\[
\text{ATTRIBUTES}(\langle \Delta, N \rangle):
\]

\[
\text{type}
\]

\[
\text{Sub\_Domain\_Location, Sub\_Domain\_Owner, Kms, ...}
\]

\[
\text{value}
\]

\[
\text{attr\_Location: } N \to \text{Sub\_Domain\_Location}
\]

\[
\text{attr\_Owner: } N \to \text{Sub\_Domain\_Owner}
\]

\[
\text{attr\_Length: } N \to \text{Kms}
\]

\[
\text{...}
\]

where \text{Sub\_Domain\_Location} could include \textit{Denmark}, \text{Sub\_Domain\_Owner} could include \textit{The Danish Road Directorate}\textsuperscript{30}, respectively \textit{BaneDanmark}\textsuperscript{31}, etcetera.

\[
\text{ATTRIBUTES}(\langle \Delta, N, Hs, L \rangle):
\]

\[
\text{type}
\]

\[
\text{LOC, LEN, ...}
\]

\[
\text{value}
\]

\[
\text{attr\_LOC: } L \to \text{LOC}
\]

\[
\text{attr\_LEN: } L \to \text{LEN}
\]

\[
\text{...}
\]

\[
\text{ATTRIBUTES} (\{ \langle \Delta, N, Hs, L \rangle, \langle \Delta, N, Hs, H \rangle \}):
\]

\[
\text{type}
\]

\[
\text{L}\Sigma = \text{HI-set, } L\Omega - \text{L}\Sigma\text{-set}
\]

\[
\text{H}\Sigma = \text{LI-set, } H\Omega - \text{H}\Sigma\text{-set}
\]

\[
\text{value}
\]

\[
\text{attr\_L}\Sigma: L \to L\Sigma
\]

\textsuperscript{30}\text{http://www.vejdirektoratet.dk/roaddirectory.asp?page=dept&objno=1024}

\textsuperscript{31}\text{http://uk.bane.dk/default_eng.asp?artikelID=931}
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\[\text{attr}_{\text{L}\Omega}: \text{L} \rightarrow \text{L}\Omega\]
\[\text{attr}_{\text{H}\Sigma}: \text{H} \rightarrow \text{H}\Sigma\]
\[\text{attr}_{\text{H}\Omega}: \text{H} \rightarrow \text{H}\Omega\]

where \text{LOC} might reveal some Bezier curve\(^{32}\) representation of the possibly curved three-dimensional location of the link in question, \text{LEN} might designate length in meters, \text{L}\Sigma [\text{H}\Sigma] designates the state of the link [hub], \text{L}\Omega [\text{H}\Omega] designates the space of all allowed states of the link [hub], etcetera.

---

**Attribute Sort Exploration:** Once the attribute sorts of a part type have been determined there remains to be “discovered” the concrete types of these sorts. We omit treatment of this point in the present version of these research notes.

### 5.2.2 Discovering Action Signatures

**General** We really should discover actions, but actually analyse function definitions. And we focus, in these research notes, on just “discovering” the function signatures of these actions. By a function signature, to repeat, we understand a functions name, say \text{fct}, and a function type expression (\text{te}), say \text{dte} \rightarrow \text{rte} where \text{dte} defines the type of the function’s definition set and \text{rte} defines the type of the function’s image, or range set.

**Function Signatures Usually Depend on Compound Domains** We use the term ‘functions’ to cover actions, events and behaviours.

We shall in general find that the signatures of actions, events and behaviours depend on types of more than one domain. Hence the schematic index set \{\ell_1\langle T_1 \rangle, \ell_2\langle T_2 \rangle, \ldots, \ell_n\langle T_n \rangle\} is used in all actions, events and behaviours discoverers.

**The \text{ACTION\_SIGNATURES} Discoverer**

100. The \text{ACTION\_SIGNATURES} meta-function applies to an index set and yields

- a set of action signatures each consisting of an action name and a pair of definition set and range type expressions where
- the type names that occur in these type expressions are defined by in the domains indexed by the index set.

\[\text{ACTION\_SIGNATURES}: \text{Index} \rightarrow \text{Text}\]

\[\text{ACTION\_SIGNATURES} (\{\ell_1\langle T_1 \rangle, \ell_2\langle T_2 \rangle, \ldots, \ell_n\langle T_n \rangle\}):\]

100a\[\text{act}_{\text{fct}_i}: \text{te}_{i\ell} \rightarrow \text{te}_{i\text{r}},\]
100b\[\text{act}_{\text{fct}_j}: \text{te}_{j\ell} \rightarrow \text{te}_{j\text{r}},\]

---

\(^{32}\)http://en.wikipedia.org/wiki/Bézier_curve


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5.2.3 Discovering Event Signature

Events are from the point of view of signatures very much like actions.

101. The EVENT\_SIGNATURES meta-function applies to an index set and yields

\textbf{a} a set of action signatures each consisting of an action name and a pair of definition set and range type expressions where

\textbf{b} the type names that occur in these type expressions are defined either in the domains indexed by the index set or by the environment (i.e., “outside” the domain $\Delta$).

\begin{align*}
\text{EVENT\_SIGNATURES:} & \text{ Index } \overset{\sim}{\rightarrow} \text{ Text} \\
\text{EVENT\_SIGNATURES}(\{\ell_1 \overset{\sim}{\rightarrow} \langle T_1 \rangle, \ell_2 \overset{\sim}{\rightarrow} \langle T_2 \rangle, \ldots, \ell_n \overset{\sim}{\rightarrow} \langle T_n \rangle\}): & \\
\text{101a } \text{evt}_f_{\text{ct}_i}: & \text{ te}_{i_d} \overset{\sim}{\rightarrow} \text{ te}_{i_r}, \\
\text{101a } \text{evt}_f_{\text{ct}_j}: & \text{ te}_{j_d} \overset{\sim}{\rightarrow} \text{ te}_{j_r}, \\
\text{101a } \ldots, & \\
\text{101a } \text{evt}_f_{\text{ct}_k}: & \text{ te}_{k_d} \overset{\sim}{\rightarrow} \text{ te}_{k_r} \\
\text{101b } & \text{ where:} \\
\text{101b } & \text{type names of } \text{ te}_{(i|j|\ldots|[k])_d} \text{ and } \text{ te}_{(i|j|\ldots|[k])_r} \text{ are type names defined by the indices which are prefixes of } \\
\text{101b } & \ell_m \overset{\sim}{\rightarrow} \langle T_m \rangle \text{ and where } T_m \text{ is in some signature } \text{ act}_f_{\text{ct}_{(i|j|\ldots|[k])}.}
\end{align*}

5.2.4 Discovering Behaviour Signatures

We choose, in these research notes, to model behaviours in CSP\textsuperscript{33}. This means that we model (synchronisation and) communication between behaviours by means of messages $m$ of type $M$, CSP channels (channel $\text{ch}:M$) and CSP

\begin{align*}
\text{output: } & \text{ch}!e \text{ [offer to deliver value of expression } e \text{ on channel ch], and} \\
\text{input: } & \text{ch}? \text{ [offer to accept a value on channel ch].}
\end{align*}

\textsuperscript{33}Other behaviour modelling languages are Petri Nets, MSCs: Message Sequence Charts, Statechart etc. We invite the reader to suggest corresponding ‘discovery’ techniques and tools.

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We allow for the declaration of single channels as well as of one, two, ..., \( n \) dimensional arrays of channels with indexes ranging over channel index types:

\[
\text{type} \quad \text{Idx}, \text{CIdx}, \text{RIdx} \ldots \\
\text{channel} \quad \text{ch}:\text{M}, \{ \text{ch}_v[vi]:\text{M}' | vi:\text{Idx} \}, \{ \text{ch}_m[ci,ri]:\text{M}'' | ci:\text{CIdx}, ri:\text{RIdx} \}, \ldots
\]

etcetera. We assume some familiarity with RSL/CSP.

A behaviour usually involves two or more distinct sub-domains.

**Example 57 (The Involved Subdomains of a Vehicle Behaviour)** Let us illustrate that behaviours usually involve two or more distinct sub-domains. A vehicle behaviour, for example, involves the vehicle subdomain, the hub subdomain (as vehicles pass through hubs), the link subdomain (as vehicles pass along links) and, for the road pricing system, also the monitor subdomain.

102. The \texttt{BEHAVIOUR\_SIGNATURES} is a meta function.

103. It applies to a set of indices and results in a text,

104. The text contains:

\begin{enumerate}
  \item a set of zero, one or more message types,
  \item a set of zero, one or more channel index types,
  \item a set of zero, one or more channel declarations,
  \item a set of one or more process signatures with each signature containing a behaviour name, an argument type expression, a result type expression, usually just \texttt{Unit}, and
  \item an input/output clause which refers to channels over which the signedatured behaviour may interact with its environment.
\end{enumerate}
Example 58 (A Vehicle Behaviour Signature Discovery) We refer, for example, to Examples 36 (Pages 56–61) and 40 (Pages 66–67).

```plaintext
let ih=(\Delta,N,LS,H),il=(\Delta,N,HS,L),iv=(\Delta,F,V),im=(\Delta,Monitor) in
BEHAVIOUR_SIGNATURES({iv,ih,iv,im}) as text
let n:N, hs=obs_HS(n), ls=obs_LS(n), vs=obs_F(PART_SORTS)(\Delta) in
where text:
  type
    VL_Msg, VH_Msg, VM_Msg
  channel
    62a. \{vh_ch[\text{attr\_VI}(v),\text{attr\_HI}(h)]|v:V,h:H\land v\in vd\land h\in hs\}:VH_Msg
    62b. \{vl_ch[\text{attr\_VI}(v),\text{attr\_LI}(h)]|v:V,l:L\land v\in vs\land h\in ls\}:VL_Msg
    62c. \text{m\_ch}:VM_Msg
  value
    64. \text{vehicle}: VI \rightarrow V \rightarrow VP \\
    83. \text{out},\text{in} \text{\{vl\_ch}[vi,li]|li:LI\land li\in xtr_LIs(ls)\}\text{out m\_ch,.. Unit}
end end
```

5.3 What Does Application Mean?

Now what does it actually mean “to apply” a discover function? We repeat our list of discoverers.

```
[90 (Page 76)] PART_SORTS,
[91 (Page 77)] HAS_A_CONCRETE_TYPE,
[92 (Page 78)] PART_TYPES,
[93 (Page 79)] UNIQUE_ID,
[96 (Page 79)] MEREEOLOGY,
[98 (Page 80)] ATTRIBUTES,
[100 (Page 82)] ACTION_SIGNATURES,
[101 (Page 83)] EVENT_SIGNATURES and
[103 (Page 84)] BEHAVIOUR_SIGNATURES.
```
Towards a Theory of Domain Descriptions

It is the domain engineer cum scientist\textsuperscript{34} who "issues" the "commands". The first "formal" domain inquiry is that of $\text{PART\_SORTS}((\Delta))$. We refer to Item 90 on page 76, for example as captured by the formulas, Items 90–90b (Page 86).

For the domain engineer to 'issue' one of the 'discovery commands' means that that person has (i) prepared his mind to study the domain and is open to impressions, (ii) decided which DISCOVERER KIND to focus on, and (iii) studied the "rules of engagement" of that command, that is which pre-requisite discoverers must first have been applied, with which index, that is, in which context the command invocation should be placed, and which results the invocation is generally expected to yield.

\subsection{PART\_SORTS}

Let us review the PART\_SORTS discoverer:

\begin{enumerate}
\item \text{PART\_SORTS}: Index $\leadsto$ Text
\item \text{PART\_SORTS}(\ell^\langle T \rangle):
\begin{enumerate}
\item tns:\{T_1,T_2,\ldots,T_m\}:TN-set \times
\item \{ obs,T_j\}: T \rightarrow T_j | T_j:tns
\end{enumerate}
\end{enumerate}

The domain analyser\textsuperscript{35} has decided to “position” the search at domain index $\ell^\langle T \rangle$ where $T = \Delta$ if $\ell = ()$ and where $T$ is some “previously discovered part type.

From Item 90a the domain analyser is guided (i.e., advised) to analyse the domain “at position $\ell^\langle T \rangle$; is the domain type $T$ a composite type of one or more subpart types? If so then decide which they are, that is: $T_1,T_2,\ldots,T_m$, that is, the “generation” of the text $\text{type } T_1,T_2,\ldots,T_m$, if not then $\text{tns} = \{}$ and no text is “generated”.

Item 90b, and given the domain analyser’s resolution of Item 90a, then directs the “generation” of $m$ observers $\text{obs},T_j\rightarrow T_j$ (for $j : \{1..m\}$).

\subsection{HAS\_A\_CONCRETE\_TYPE}

Let us review the HAS\_A\_CONCRETE\_TYPE analyser:

\begin{enumerate}
\item \text{HAS\_A\_CONCRETE\_TYPE}: Index $\rightarrow$ Bool
\item \text{HAS\_A\_CONCRETE\_TYPE}(\ell^\langle t \rangle):true|false
\end{enumerate}

Item 91 directs the domain analyser to decide whether the domain type $T$ at “position” $\ell^\langle t \rangle$ should be given a concrete type definition. It is a decision solely at the discretion of the domain analyser whether domain type $T$ should be given a concrete type definition, and, as we shall see next, which concrete type it should then be “given”, that is, how it should be “concretely abstractly” modelled.

\subsection{PART\_TYPES}

Let us review the PART\_TYPES analyser:

\begin{enumerate}
\item \text{PART\_TYPES}: Index $\leadsto$ Text
\item \text{PART\_TYPES}(\ell^\langle t \rangle):
\begin{enumerate}
\item \text{type } T = TE,
\end{enumerate}
\end{enumerate}

\textsuperscript{34}When we write: domain engineer cum scientist we mean to say that the domain engineer really is performing a scientific inquiry.

\textsuperscript{35}We use the alternative, synonymous terms: ‘domain engineer’, ‘domain describer’, ‘domain scientist’.
92b. \( T_1 \) or \( T_1 = T_{E_1} \)
92b. \( T_2 \) or \( T_2 = T_{E_2} \)
92b. ...
92b. \( T_n \) or \( T_n = T_{E_n} \)
92c. \( \text{pre: \{HAS\_A\_CONCRETE\_TYPE}(\ell\langle t \rangle) \}

The domain analyser has decided to “position” the search at domain index \( \ell\langle T \rangle \) where \( T = \Delta \) if \( \ell = \langle \rangle \) and where \( T \) is some “previously discovered part type.

From Item 92 the domain analyser is guided (i.e., advised) to analyse the domain “at position \( \ell\langle T \rangle \): can a reasonably abstract, yet concrete type definition be given for \( T \)? If so then decide which it should be, that is, should it be an atomic type a number type, \textbf{Intg}, \textbf{Rat}, \textbf{Real}, a Boolean type, \textbf{Bool}, or a token type, f.ex. \textbf{TOKEN} A token type is a further undefined atomic type — typically used to model identifiers.; or should it be a composite type either a set type: \textbf{TE}: \textbf{T}_{\text{-set}} of \textbf{te}: \textbf{e}_{\text{-infset}}, or a Cartesian type: \textbf{TE}: \textbf{T}_1 \times \textbf{T}_2 \times \ldots \times \textbf{T}_m, or a list type: \textbf{TE}: \textbf{T}_1 \text{ or te: } \textbf{T}_1 \text{ or a map type: } \textbf{TE}: \textbf{T}_d \rightarrow \textbf{T}_r ? \) In either case the text type \( T = \textbf{TE}, T_1 \) or \( T_{E_1}, T_2 \) or \( T_{E_2}, \ldots, T_n \) or \( T_{E_n} \) is generated where \( \textbf{TE}(\textbf{TE}_a) \) is a type expression whose so far undefined type names \( T_1, T_2, \ldots, T_n \) must be defined, either as sorts, or a concrete types.

5.3.4 \textbf{UNIQUE\_ID}

Let us review the \textbf{UNIQUE\_ID} analyser:

\[
\begin{align*}
\text{value} & \quad 93. \quad \text{UNIQUE\_ID}: \text{Index} \rightarrow \text{Text} \\
\text{type} & \quad 93.a \quad \text{UNIQUE\_ID}(\ell\langle T \rangle): \\
\text{value} & \quad 93.b \quad \text{type} \\
\text{value} & \quad 93.c \quad \text{TI} \\
\text{value} & \quad 93.d \quad \text{uid\_TI}: \text{T} \rightarrow \text{TI} \\
\end{align*}
\]

Item 93.a inquires as to the Line 93.b \textbf{type} name Line 93.c of the inquired part type’s unique identifiers Line 93.d and the function signature \textbf{value} Line 93.e of the observer, \textbf{uid\_TI}, name, the definition set type (\textbf{T}, of course) and the range set type (\textbf{TI} — obviously). Thus, the only real “new” “discovery” here is the name, \textbf{TI}, of the unique identifier type.

Etcetera, etcetera.

5.4 \textbf{Discussion}

We have presented a set of discoverers:

\[
\begin{align*}
[90 \text{(Page 76)}] & \quad \text{PART\_SORTS,} \\
[91 \text{(Page 77)}] & \quad \text{HAS\_A\_CONCRETE\_TYPE,} \\
[92 \text{(Page 78)}] & \quad \text{PART\_TYPES,} \\
[93 \text{(Page 79)}] & \quad \text{UNIQUE\_ID,} \\
[96 \text{(Page 79)}] & \quad \text{MEREEOLOGY,} \\
[98 \text{(Page 80)}] & \quad \text{ATTRIBUTES,} \\
[100 \text{(Page 82)}] & \quad \text{ACTION\_SIGNATURES,} \\
[101 \text{(Page 83)}] & \quad \text{EVENT\_SIGNATURES and}
\end{align*}
\]
There is much more to be said: About a meta-state component in which is kept the “text” so far generated. A component from which one can see which indices and hence which type names have so far been “generated”, and on the basis of which one can perform tests of well-formedness of generated text, etcetera, etcetera,
6 Conclusion

6.1 General

6.2 What Have We Achieved?

6.3 Other Formal Models

6.4 Research Issues

6.5 Engineering Issues

6.6 Comparable Work

6.7 Acknowledgements