

# NICE TO SEE YOU BACK

# **Begin of Lecture 2: Last Session** — **Discrete Endurant Entities**

#### Parts

#### FM 2012 Tutorial, Dines Bjørner, Paris, 28 August 2012

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

# **Tutorial Schedule**

• Lectures 1–2 9:00-9:40 + 9:50-10:301 Introduction Slides 1-35 $\sqrt{2}$  Endurant Entities: Slides 36–110 Parts • Lectures 3–5 11:00-11:15 + 11:20-11:45 + 11:50-12:30**3 Endurant Entities: Materials, States** Slides 111–142 **4** Perdurant Entities: Actions and Events Slides 143–174 Slides 175–285 5 Perdurant Entities: Behaviours 12:30 - 14:00Lunch • Lectures 6–7 14:00-14:40 + 14:50-15:30Slides 286–339 6 A Calculus: Analysers, Parts and Materials 7 A Calculus: Function Signatures and Laws Slides 340–377 • Lecture 8–9 16:00-16:40 + 16:50-17:30**8 Domain and Interface Requirements** Slides 378–424 9 **Conclusion: Comparison to Other Work** Slides 428–460 **Conclusion:** What Have We Achieved Slides 425-427 + 461-472

# 2. Domain Entities

- In this tutorial we shall divide the phenomena we can observe and whose properties we can ascertain into two kinds:
  - $\circledast$  the endurant entities and
  - $\otimes$  the perdurant entities.
- Another "division" is of the phenomena and their properties into

   • the discrete entities and
   • the continuous entities.
- You can have it, i.e., the the analysis and the presentation, either way.

- By a **domain** we shall understand a suitably delineated set of **observable entities** and **abstraction**s of these, that is, of
  - $\otimes$  discrete parts and
  - $\otimes$  continuous materials and,
  - $\otimes$  discrete actions

(operation applications causing state changes),

 $\otimes$  discrete events

("spurious" state changes not [intentionally] caused by actions),

#### 

(seen as sets of sequences of actions, events and behaviours) and

#### **« continuous behaviour**s

(abstracted as continuous functions in space and/or time).

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

# **2.1. From Observations to Abstractions**

- When we observe a domain we observe **instances** of **entities**;
- but when we describe those instances
  - (which we shall call **value**)
  - $\otimes$  we describe, not the values,
  - $\otimes$  but their type and properties.

    - actions, events and behaviours, all, have types and values, namely as expressed by their signatures; and
    - actions, events and behaviours have properties, namely as expressed by their function definitions.

# 2.2. Algebras

 ${\tt $\infty$}$  a set of  ${\tt parts}$  and

 ${\tt $\infty$}$  a set of materials

and

« a set of **perdurant**s: operations on entities.

These operations yield parts or materials.

• With that in mind we shall try view a domain as an algebra, of some kind, of

**∞ part**s and

 $\circledast$  actions, events and behaviours.

# 2.3. Phenomena

- **Phenomena:** By a **phenomenon** we shall understand
  - something that can be observed by the human sensesor by equipment based on laws of physics and chemistry.
- Those phenomena that can be observed by
  - the human eye or
    touched, for example, by human hands
  - « we call parts and materials.
- Those phenomena that can be observed of parts and materials
  & can usually be measured
  & and we call them properties of these parts and those materials.

# **2.4.** Entities

- Ontologically we distinguish between two kinds of domain entities:

   **endurant entities** and
  - « perdurant entities.
- We shall characterise these two terms:
  - $\circledast \mbox{endurants}$  on Slide 49 and
  - **∞ perdurant**s on Slide 144.
- This distinction is supported by current literature on **ontology** [BarrySmith1993].
- In this section of this lecture we shall not enter a discourse on

$\ll$ "things",	$\otimes objects,$
$\otimes$ entities,	∞ etcetera.

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

#### **2.4.1. A Description Bias**

- One of several "twists"
  - ☆ that make the TripTych form of domain engineering
  - $\circledast$  distinct from that of ontological engineering
  - $\otimes$  is that we use a model-oriented formal specification approach<sup>12</sup>
  - where usual ontology formalisation languages are variants of Lisp's [Lisp1] S-expressions.
  - & KIF: Knowledge Interchange Format, http://www.ksl.stanford.edu/knowledge-sharing/kif/ is a leading example.

<sup>&</sup>lt;sup>12</sup>RAISE [RaiseMethod]. Our remarks in this section apply equally well had we instead chosen either of the Alloy [alloy], Event B [JRAbrial:TheBBooks], VDM [e:db:Bj78bwo,e:db:Bj82b,jf-pgl-97] or Z [m:z:jd+jcppw96] formal specification languages.

- The bias is now this:
  - The model-oriented languages mentioned in this section all share the following:
    - ∞ (a) a **type concept** and facilities for defining types, that is: endurants (parts), and
    - (b) a function concept and facilities for defining functions (notably including predicates), that is: perdurants (actions and events).
    - $\infty$  (c) RSL further has constructs for defining processes, which we shall use to model behaviours.

# 2.4.2. An 'Upper Ontology'

- $\bullet$  By an upper ontology we shall understand
  - $\otimes$  a relatively small, ground set of ontology expressions
  - ∞ which form a basis for a usually very much larger set of ontology expressions.

- The need for introducing the notion of an **upper ontology** arose, in the late 1980s to early 1990s as follows:
  - $\otimes$  usually an ontology was (is) expressed in some very basic language, viz., Lisp-like S-expressions^{13}.
  - $\otimes$  This was necessitated by the desire to be able to share ontologies between many computing applications worldwide.
  - Then it was found that several ontologies shared initial bases in terms of which the rest of their ontologies were formulated.

<sup>&</sup>lt;sup>13</sup>Ontology languages: KIF http://www.ksl.stanford.edu/knowledge-sharing/kif/-#manual, OWL [Ontology Web Language] [OWL:2009], ISO Common Logic [ISO:CL:2007]

<sup>©</sup> Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

- We therefore consider the following model-oriented specification language constructs as forming an upper ontology:
  - $\circledast$  types, ground types, type expressions and type definitions;
  - $\circledast$  functions, function signatures and function definitions;
  - $\circledast$  processes, process signatures and process definitions,
  - as constituting an **upper level ontology** for **TripTych domain description**s.
- That is, every domain description is structured with respect to: \* parts and materials using types,

  - « events using predicates,
  - **« discrete** behaviours using **process**es and
  - « continuous behaviours using partial differential equations.

# 3. Endurants

- There is sort of a dichotomy buried in our treating endurants before perdurants. The dichotomy is this:
  - one could claim that the perdurants, i.e., the actions, events and behaviours is "what it, the domain, is all about";
- To describe these, however, we need refer to endurants!

48

# 3.1. General

Wikipedia:

• By an endurant (also known as a continuant or a substance) we shall understand an entity

*∞* that can be observed, i.e., perceived or conceived,

 $\otimes$  as a complete concept,

 $\otimes$  at no matter which given snapshot of time.

• Were we to freeze time

 $\otimes$  we would still be able to observe the entire endurant.

#### **3.2. Discrete and Continuous Endurants**

• We distinguish between

**« discrete endurant**s, which we shall call **part**s, and **« continuous endurant**s, which we shall call **material**s.

We motivate and characterise this distinction.

- By a **discrete endurant**, that is, a **part**, we shall understand something which is
  - $\otimes$  separate or distinct in form or concept,
  - $\otimes$  consisting of distinct or separate parts.
- By a **continuous endurant**, that is, a **material**, we shall understand something which is
  - $\otimes$  prolonged without interruption,
  - $\otimes$  in an unbroken series or pattern.
- $\bullet$  We shall
  - ☆ first treat the idea of discrete endurant, that is, a part (Slides 51−110),

# 4. Discrete Endurants: Parts 4.1. Atomic and Composite Parts

- Parts may be analysed into disjoint sets of
- atomic parts and

- composite parts.
- Atomic parts are those which,
  - $\otimes$  in a given context,
- Composite parts are those which,
  - $\otimes$  in a given context,
- A sub-part is a part.

**Example: 4 Atomic and/or Composite Parts.** To one person a part may be atomic; to another person the same part may be composite.

- It is the domain describer who decides the outcome of this aspect of domain analysis.
  - - $\infty$  For the domain of ferrying cars with passengers
    - ∞ persons are considered parts.
  - - $\infty$  For the domain of medical surgery
    - $\infty$  persons may be considered composite parts.

52

### **Example:** 5 **Container Lines.**

- We shall presently consider **container**s (as used in container line shipping) to be atomic parts.
- And we shall consider a **container vessel** to be a composite part consisting of
  - $\circledast$  an indexed set of container bays
  - $\otimes$  where each container bay consists of indexed set of container rows
  - w where each container row consists of indexed set of container
    stacks
  - $\otimes$  where each container stack consists of a linearly indexed sequence of containers.
- Thus container vessels, container bays, container rows and container stacks are composite parts.

A Precursor for Requirements Engineering

### 4.1.1. Atomic Parts

# • When we observe

what we have decided, i.e., analysed, to be an endurant,
more specifically an atomic part, of a domain,
we are observing an instance of an atomic part.

- When we describe those instances
  - ∞ we describe, not their values, i.e., the instances,
    ∞ but their
    - ${\scriptstyle \scriptsize \ensuremath{\varpi}}$  type and
    - or properties.

- In this section on **endurant entities** we shall unfold what these properties might be.
- But, for now, we focus on the **type** of the observed **atomic part**.

• What does it mean for a number of atomic parts to be of "the same kind" ?

 $\otimes$  It means

∞ that we have decided,

∞ for any pair of **part**s considered of the same kind,

 $\infty$  that the kinds of properties,

\* for such two parts,

∞ are "the same",

\* that is, of the same type, but possibly of different values,
and that a number of different, other "facets",
are not taken into consideration.

# • That is,

- $\circledast$  we abstract a collection of atomic parts
- $\otimes$  to be of the same kind,
- - $\infty$  those that are of the analysed kind, and
  - $\infty$  those that are not.

- It is now our description choice to associate with a set of **atomic parts** of *"the same kind"* 
  - $\circledast$  a part type (by suggesting a name for that type, for example,  $\mathsf{T})$  and
  - ∞ a set of **properties** (of its values):
    - $\mbox{$\ensuremath{$\odot$}$}$  unique identifier,
    - ${\scriptstyle \circledcirc}$  mereology and
    - attributes.

- Later we shall introduce **discrete perdurant**s (actions, events and behaviours) whose **signature**s involves (possibly amongst others) type **T**.
- Now we can characterise *"of the same kind"* atomic part facets<sup>14</sup>
  - $\circledast$  being of the same, named part type,
  - $\otimes$  having the same unique identifier type,
  - $\otimes$  having the same  $\ensuremath{\mathsf{mereology}}$ 
    - (but not necessarily the same mereology values), and
  - $\otimes$  having the same set of  ${\tt attribute}{\sf s}$ 
    - (but not necessarily of the same attribute values),
- The *"same kind"* criteria apply equally well to composite part facets.

<sup>&</sup>lt;sup>14</sup>as well as "of the same kind" **composite part facets**.

# **Example:** 6 Transport Nets: Atomic Parts (I).

- The types of atomic transportation net parts are:
  - $\otimes$  hubs, say of type  $\mathsf{H},$  and
  - $\otimes$  links, say of type  $\mathsf{L}.$
- The chosen mereology associates with every hub and link a
  - $\otimes$  distinct unique identifiers
  - $\otimes$  (of types HI and LI respectively), and, vice versa,
  - $\otimes$  how hubs and links are connected:
    - $\infty$  hubs to any number of links and
    - $\infty$  links to exactly two distinct hubs.

• The chosen attributes of

« hubs include

multiple hub location,
 15

 $\infty$  hub design<sup>15</sup>,

 $\otimes$  and of links include

link location,

 $\textcircled{\ } \texttt{o} \ \texttt{link} \ \texttt{length},$ 

hub traffic state<sup>16</sup>,
hub traffic state space<sup>17</sup>, etc.;

link traffic state<sup>18</sup>,
link traffic state space<sup>19</sup>, etc.

• With these mereologies and attributes we see that we can consider hubs and links as different kinds of atomic parts.

<sup>15</sup>Design: simple crossing, freeway "cloverleaf" interchange, etc.

<sup>16</sup>A hub traffic state is (for example) a set of pairs of link identifiers where each such pair designates that traffic can move from the first designated link to the second.

<sup>17</sup>A hub state space is (for example) the set of all hub traffic states that a hub may range over.

<sup>18</sup>A link traffic state is (for example) a set of zero to two distinct pairs of the hub identifiers of the link mereology.

<sup>19</sup>A link traffic state space is (for example) the set of all link traffic states that a link may range over.

### **Observers for Atomic Parts**

- Let the domain describer decide
  - $\otimes$  that a type, A (or  $\Delta$ ), is atomic,
  - « hence that it does not consists of sub-parts.
- Hence there are no **observer** to be associated with A (or  $\Delta$ ).

# 4.1.2. Composite Parts

- $\bullet$  The domain describer has chosen to consider

  - ∞ to be a composite part (i.e., a composite part type).
- Now the domain describer has to analyse the types of the sub-parts of the composite part.
  - ∞ There may be just one "kind of" sub-part of a composite part<sup>20</sup>,
    ∞ or there may be more than one "kind of"<sup>21</sup>.
- For each such **sub-part type** 
  - $\otimes$  the domain describer decides on
  - $\otimes$  an appropriate, distinct  $\ensuremath{\mathsf{type}}\xspace$  and
  - $\otimes$  a sub-part observer (i.e., a function signature).

<sup>20</sup>that is, only one sub-part type <sup>21</sup>that is, more than one sub-part type

# **Example: 7 Container Vessels: Composite Parts.** We bring pairs of informal, narrative description texts and formalisations.

- $\bullet$  For a container vessel, say of type V, we have
  - *⊗* Narrative:
    - $\infty$  A container vessel, v:V, consists of container bays, bs:BS.
    - ∞ A container bay, b:B, consists of container rows, rs:RS.
    - $\infty$  A container row, r:R, consists of container stacks, ss:SS.
    - A container stack, s:S, consists of a linearly indexed sequence of containers.

#### $\otimes$ Formalisation:

type V,BS, value obs\_BS:  $V \rightarrow BS$ , type B,RS, value obs\_RS:  $B \rightarrow RS$ , type R,SS, value obs\_CS:  $R \rightarrow SS$ , type SS,S, value obs\_S:  $SS \rightarrow S$ , type S = C\*.

# 4.1.3. Abstract Types, Sorts, and Concrete Types

- - ∞ but is otherwise undefined, that is,
    - ∞ is a space of undefined mathematical quantities,
      - \* where these are given properties
      - \* which we may express in terms of axioms over sort (including property) values.

• By a concrete type we shall understand a type, T,

which has been given both a nameand a defining type expression of, for example the form

 $\otimes$  where A, B, ..., C are type names or type expressions.

# **Example: 8 Container Bays.** We continue Example 7 on page 64. **type** $Bs = Bld \overrightarrow{m} B$ , **value** $obs\_Bs: BS \rightarrow Bs$ , **type** $Rs = Rld \overrightarrow{m} R$ , **value** $obs\_Rs: B \rightarrow Rs$ , **type** $Ss = Sld \overrightarrow{m} S$ ,

value obs\_Ss:  $R \rightarrow Ss$ ,

type 
$$S = C^*$$
.

# **Observers for Composite Parts I/II**

- $\bullet$  We can initially consider these types B, C, ..., D, as abstract types, or sorts, as we shall mostly call them.

#### **Observers for Composite Parts II/II**

# 4.2. Properties

- Endurants have properties.
  - $\otimes$  Properties are
    - what makes up a parts (and materials) and,
    - with property values distinguishes one part from another part and
      - one material from another material.
  - $\otimes$  We name properties.
    - **• Properties** of **parts** and **materials** can be given distinct names.
    - ${\scriptstyle \scriptsize \varpi}$  We let these names also be the property type name.
    - member Hence two parts (materials) of the same part type (material type)

have the same set of property type names.

- Properties are all that distinguishes parts (and materials).
  - The part types (material types) in themselves do not express properties.
  - ∞ They express a class of parts (respectively materials).

  - $\otimes$  have the same property types.

- For pragmatic reasons we distinguish between three kinds of properties:
  - « unique identifiers, « mereology, and » attributes.
- If you "remove" a property from a part
  - $\otimes$  it "looses" its (former) part type,
  - $\ll$  to, in a sense, attain another part type:
    - perhaps of another, existing one,
    - ∞ or a new "created" one.
- But we do not know how to model removal of a property from an endurant value!<sup>22</sup>

<sup>&</sup>lt;sup>22</sup>And we see no need for describing such type-changes. Crude oil does not "morph" into fuel oil, diesel oil, kerosene and petroleum. Crude oil is consumed and the fractions result from distillation, for example, in an oil refinery.

#### **Example:** 9 Atomic Part Property Kinds.

- We distinguish between two kinds of persons:
  - « 'living persons' and 'deceased persons';
  - - $\infty$  LP: living person, with a set of properties,
    - $\infty$  DP: deceased person, with a, most likely, different set of properties.
- All persons have been born, hence have a birth date (static attributes).
- Only deceased persons have a (well-defined) death date.

- All persons also have height and weight profiles (i.e., with dated values, i.e., dynamic attributes).
- One can always associate a **unique identifier** with each person.
- Persons are related, family-wise:
  - « have parents (living or deceased),
  - ∞ (up to four known) grandparents, etc.,
  - « may have brothers and sisters (zero or more),
  - ∞ may have children (zero or more), etc.

# 4.2.1. Unique Identification

• We can assume that all **part**s

- $\circledast$  of the same part type
- « can be uniquely distinguished,
- « hence can be given unique identifications.

## **Unique Identification**

- With every part, whether atomic or composite we shall associate a unique part identifier, of just unique identifier.
- Thus we shall associate with part type T

the unique part type identifier type TI,

 $\otimes$  and a unique part identifier observer function, uid\_TI: T \rightarrow TI.

 $\bullet$  These associations (TI and  $\mathsf{uid}_{-}\mathsf{TI})$  are, however,

∞ usually expressed explicitly,

∞ whether they are ("subsequently") needed!

- The unique identifier of a part
  - $\otimes$  can not be changed;
  - $\otimes$  hence we can say that

no matter what a given part's property values may take on,
that part cannot be confused with any other part.

• Since we can talk about this concept of **unique identification**,

∞ we can **abstract**ly describe it —

and do not have to bother about any representation,that is, whether we can humanly observe unique identifiers.

# 4.2.2. Mereology

- Mereology [CasatiVarzi1999]<sup>23</sup> (from the Greek  $\mu\epsilon\rho\sigma\varsigma$  'part') is
  - « the theory of part-hood relations:
  - $\circledast$  of the relations of part to whole and
  - « the relations of part to part within a whole.

<sup>&</sup>lt;sup>23</sup>Achille Varzi: Mereology, http://plato.stanford.edu/entries/mereology/

- For pragmatic reasons we choose to model the mereology of a domain in either of two ways
  - ∞ either by defining a concrete type as a model of the composite type,
  - ∞ or by endowing the sub-parts of the composite part with structures of unique part identifiers.
  - or by suitable combinations of these.

**Example: 10 Container Bays, Etcetera: Mereology.** First we show how to model indexed set of container bays, rows and stacks for the previous example.

- Narrative:

  - (iv) A stack is a linear indexed sequence of containers, c:C.

#### • Formalisation:

```
(i) type BS, B, Bld,
                 Bs = BId \xrightarrow{m} B,
        value obs_Bs: BS \rightarrow Bs
                  (or obs_Bs: BS \rightarrow (BId \xrightarrow{m} B));
∞ (ii) type RS, R, Rld,
                  Rs = RId \xrightarrow{\pi} R,
        value obs Rs: RS \rightarrow Rs
                  (or obs_Rs: RS\rightarrow(RId \overrightarrow{m} R));
(iii) type SS, S, Sld,
                  Ss = SId \rightarrow S;
\otimes (iv) type C,
                  S = C^*
```

#### **Example:** 11 **Transport Nets: Mereology.**

- We show how to model a **mereology** 
  - $\circledast$  for a transport net of links and hubs.
- Narrative:
  - (i) Hubs and links are endowed with unique hub, respectively link identifiers.
  - (ii) Each hub is furthermore endowed with a hub mereology which lists the unique link identifiers of all the links attached to the hub.
- (iii) Each link is furthermore endowed with a link mereology which lists the set of the two unique hub identifiers of the hubs attached to the link.
- (iv) Link identifiers of hubs and hub identifiers of links must designate hubs, respectively links of the net.

- Formalisation:
  - (i) **type** H, HI, L, LI; **value**
  - (ii) uid\_HI:H $\rightarrow$ HI, uid\_LI:L $\rightarrow$ LI, mereo\_H:H $\rightarrow$ LI-set, mereo\_L:L $\rightarrow$ HI-set,

#### axiom

```
(iii) \forall l:L · card mereo_L(l) = 2

(iv) \forall n:N, l:L, h:H · l \in obs_Ls(obs_LS(n)) \land h \in obs__Hs(obs_HS(n))

\forall hi:HI · hi \in mereo_L(l) \Rightarrow

\exists h':H·h' \in obs_Hs(obs_HS(n)) \land uid_HI(h)=hi

\land \forall li:LI · li \in mereo_H(h) \Rightarrow

\exists l':L·l' \in obs_Ls(obs_LS(n)) \land uid_LI(l)=li
```

#### **Concrete Models of Mereology**

The concrete mereology example models above illustrated maps and sequences as such models.

- In general we can model mereologies in terms of
  - (i) sets: A-set, (ii) lists: A\*, and
  - $\otimes$  (ii) Cartesians:  $A_1 \times A_2 \times \ldots \times A_m$ ,  $\otimes$  (iv) maps:  $A_{\overline{m}} B$ ,

where A, A<sub>1</sub>, A<sub>2</sub>,...,A<sub>m</sub> and B are types [we assume that they are type names] and where the A<sub>1</sub>, A<sub>2</sub>,...,A<sub>m</sub> type names need not be distinct.

- Additional concrete types, say D, can be defined by concrete type definitions, D=E, where E is either of the type expressions (i-iv) given above or (v) E<sub>i</sub>|E<sub>j</sub>, or (vi) (E<sub>i</sub>). where E<sub>k</sub> (for suitable k) are either of (i-vi).
- Finally it may be necessary to express well-formedness predicates for concretely modelled mereologies.

### **Abstract Models of Mereology**

Abstractly modelling mereology of parts, to us, means the following.

• With part types  $\mathsf{P}_1, \mathsf{P}_2, \ldots, \mathsf{P}_n$ 

 $\otimes$  is associated the unique part identifier types,  $\Pi_1, \Pi_2, \ldots, \Pi_n$ ,  $\otimes$  that is **uid**\_ $\Pi_i$ :  $\mathsf{P}_i \rightarrow \Pi_i$  for  $i \in \{1..n\}$ ,

• and with each part type,  $\mathsf{P}_i$ ,

∞ is then associated a **mereology** observer,

 $\otimes$  mereo\_P<sub>i</sub>:  $P_i \rightarrow \Pi_j$ -set  $\times \Pi_k$ -set  $\times ... \times \Pi_\ell$ -set,

 $\bullet$  such that for all  $\mathsf{p}{:}\mathsf{Pi}$  we have that

• Finally it may be necessary to express axioms for abstractly modelled mereologies.

#### A Precursor for Requirements Engineering

- How **parts** are related to other **parts** 
  - ∞ is really a modelling choice, made by the **domain describer**.
  - $\otimes$  It is not necessarily something
    - that is obvious
    - from observing the **part**s.

## **Example:** 12 **Pipelines: A Physical Mereology.**

- Let pipes of a pipe line be composed with valves, pumps, forks and joins of that pipe line.
- Pipes, valves, pumps, forks and joins (i.e., pipe line units) are given unique pipe, valve, pump, fork and join identifiers.
- A mereology for the pipe line could now endow pipes, valves and pumps with
  - $\otimes$  one input unique identifier, that of the predecessor successor unit, and
  - $\otimes$  one output unique identifier, that of the successor unit.
- Forks would then be endowed with two input unique identifiers, and
  one out put unique identifier;
- and joins "the other way around".

#### A Precursor for Requirements Engineering

### **Example:** 13 **Documents:** A Conceptual Mereology.

• The mereology of, for example, this document,

 $\otimes$  that is, of the tutorial slides,

is determined by the author.

- There unfolds, while writing the document,
  - $\otimes$  a set of unique identifiers
  - $\otimes$  for section, subsection, sub-subsection, paragraph, etc., units. and

- This occurs as the author necessarily
  - $\otimes$  inserts cross-references,
    - $\infty$  in unit texts to other units, and
    - ∞ from unit texts to other documents (i.e., 'citations');
  - $\otimes$  and while inserting "page" shifts for the slides.
- From those inserted references there emerges what we could call the document mereology.
- The "design" of mereologies improves with experience.

A Precursor for Requirements Engineering

### **Example:** 14 **Pipelines:** Mereology.

- We divert from our line of examples centered around
  - ∞ transport nets and, to some degree,
  - $\otimes$  container transport,
- to bring a second, in a series of examples
  - $\otimes$  on pipelines
  - $\otimes$  (for liquid or gaseous material flow).

- 1. A pipeline consists of connected units, u:U.
- 2. Units have unique identifiers.
- 3. And units have mereologies, ui:UI:
  - (a) pump, pu:Pu, pipe, pi:Pi, and valve, va:Va, units have one input connector and one output connector;
  - (b) fork, fo:Fo, [join, jo:Jo] units have one [two] input connector[s] and two [one] output connector[s];
  - (c) well, we:We, [sink, si:Si] units have zero [one] input connector and one [zero] output connector.
  - (d) Connectors of a unit are designated by the unit identifier of the connected unit.
  - (e) The auxiliary **sel\_Uls\_in** selector function selects the unique identifiers of pipeline units providing input to a unit;
  - (f)  $sel_Uls_out$  selects unique identifiers of output recipients.

A Precursor for Requirements Engineering

91

#### type 1. U = Pu | Pi | Va | Fo | Jo | Si | We2 UI value 2. uid U: U $\rightarrow$ UI 3. mereo\_U: $U \rightarrow UI-set \times UI-set$ 3. wf\_mereo\_U: $U \rightarrow Bool$ 3. wf\_mereo\_U(u) $\equiv$ 3(a). $is_{-}(Pu|Pi|Va)(u) \rightarrow card iusi = 1 = card ouis,$ 3(b). is\_Fo(u) $\rightarrow$ card iuis = 1 $\wedge$ card ouis = 2, 3(b). is\_Jo(u) $\rightarrow$ card iuis = 2 $\land$ card ouis = 1, 3(c). is\_We(u) $\rightarrow$ card iuis = $0 \land$ card ouis = 1, 3(d). is\_Si(u) $\rightarrow$ card iuis = 1 $\wedge$ card ouis = 0 3(e). sel\_UIs\_in 3(e). sel\_UIs\_in(u) $\equiv$ **let** (iuis,\_)=mereo\_U(u) **in** iuis **end** 3(f). sel\_out: U $\rightarrow$ UI-set

3(f). sel\_UIs\_out(u)  $\equiv$  **let** (\_\_,ouis)=mereo\_U(u) **in** ouis **end** 

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

• We omit treatment of axioms for pipeline units

« being indeed connected to existing other pipeline units.

∞ We refer to Example 22 on page 119 and 23 on page 123.

## 4.2.3. Attributes

- By an **attribute** of a part, **p**:**P**, we shall understand
  - « some observable property, some phenomenon,
  - $\circledast$  that is not a  $\mathsf{sub-part}$  of p
  - $\otimes$  but which characterises  ${\bf p}$
  - $\otimes$  such that all parts of type  $\mathsf P$  have that attribute and
  - such that "removing" that attribute from p
    (if such was possible)
    "renders" the type of p undefined.
- We ascribe types to attributes not, therefore, to be confused with types of (their) parts.

#### **Example:** 15 Attributes.

• Example attributes of links of a transport net are:

- $\otimes$  length LEN,
- $\otimes$  location LOC,
- $\otimes$  state  $\mathsf{L}\Sigma$  and
- $\otimes$  state space L $\Omega$ ,
- Example attributes of a person could be:
  - $\otimes$  name NAM,
  - $\otimes$  birth date  $\mathsf{BID},$
  - $\otimes$  gender GDR,
  - $\otimes$  weight WGT,
  - $\otimes$  height  $\mathsf{HGT}$  and
  - $\otimes$  address ADR.

- Example attributes of a transport net could be:
  - $\otimes$  name of the net,
  - $\otimes$  legal owner of the net,
  - $\otimes$  a map of the net,

 $\otimes$  etc.

- Example attributes of a container vessel could be:
  - « name of container vessel,
  - $\otimes$  vessel dimensions,
  - $\otimes$  vessel tonnage (TEU),
  - $\otimes$  vessel owner,
  - $\otimes$  current stowage plan,
  - $\otimes$  current voyage plan, etc.

## 4.2.3.1 Static and Dynamic Attributes

- By a **static attribute** we mean an attribute (of a part) whose value remains fixed.
- By a **dynamic attribute** we mean an attribute (of a part) whose value may vary.

### **Example:** 16 **Static and Dynamic Attributes.**

- The length and location attributes of links are static.
- The state and state space attributes of links and hubs are dynamic.
- The birth-date attribute of a person is considered static.
- The height and weight attributes of a person are dynamic.
- The map of a transport net may be considered dynamic.
- The current stowage and the current voyage plans of a vessel should be considered dynamic.

#### Attribute Types and Observers, I/II

- $\bullet$  Let the domain describer decide that parts of type  $\mathsf{P}$
- have attributes of types  $A_1$ ,  $A_2$ , ...,  $A_t$ .
- This means that the following two formal clauses arise:

 $\otimes$  P, A<sub>1</sub>, A<sub>2</sub>, ..., A<sub>t</sub> and  $\otimes$  attr\_A<sub>1</sub>:P $\rightarrow$ A<sub>1</sub>, attr\_A<sub>2</sub>:P $\rightarrow$ A<sub>2</sub>, ..., attr\_A<sub>t</sub>:P $\rightarrow$ A<sub>t</sub>

### Attribute Types and Observers, II/II

• We may wish to annotate the list of **attribute type names** as to whether they are static or dynamic, that is,

```
\circledast whether <code>values</code> of some attribute type
```

```
⇔ vary or
```

```
∞ remain fixed.
```

• The prefix attr\_ distinguishes attribute observers from part observers (obs\_) and mereology observers (uid\_, mereo\_).

#### 4.3. Shared Attributes and Properties

- Shared attributes and shared properties
  - « play an important rôle in understanding domains.

## 4.3.1. Attribute Naming

- We now *impose a restriction* on the naming of **part attributes**.
  - $\circledast \ If \ \text{attribute} s$ 
    - $\infty$  of two different parts
    - $\infty$  of different part types
    - $\infty$  are identically named
    - $\infty$  then attributes must be somehow related, over time!
  - « The "somehow" relationship must be described.

#### **Example:** 17 Shared Bus Time Tables.

- Let our domain include that of *bus time tables* for *busses* on a *bus transport net* as described in many examples in this tutorial.
- We can then imagine a *bus transport net* as containing the following parts:
- For the sake of argument we consider a *bus time table* to be an attribute of the *bus management system*.
- And we also consider *bus time tables* to be attributes of *busses*.

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

• We think of the bus time table of a bus

\* to be that subset of the
bus management system bus time table
\* which corresponds to the bus' line number.

• By saying that bus time tables

"corresponds" to well-defined subsets of
the bus management system bus time table

we mean the following

- $\otimes$  The value of the bus bus time table
- $\otimes$  must at every time

# 4.3.2. Attribute Sharing

• We say that two parts,

« of no matter what part type,

#### « *share* an attribute,

- $\otimes$  if the following is the case:
  - the corresponding part types (and hence the parts)
  - ${\scriptstyle \circledcirc}$  have identically named attributes.
  - We say that identically named attributes designate shared attributes.
- We do not present the corresponding invariants over parts with identically named attributes.

#### © Dines Bjørner 2012, DTU Informatics, Techn.Univ.of Denmark – July 31, 2012: 09:02

# 4.4. Shared Properties

- We say that two **parts**,
  - « of no matter what part type,
  - $\circledast$  share a property,
  - $\otimes$  if either of the following is the case:
    - (i) either the corresponding part types (and hence the parts) have shared attributes;
    - (ii) or the unique identifier type of one of the parts potentially is in the mereology type of the other part;
      (iii) or both.
  - We do not present the corresponding invariants over parts with
     shared properties.

## 4.5. Summary of Discrete Endurants

- We have introduced the **endurant** notions of **atomic part**s and **composite part**s:
  - $\circledast$  part types,
  - ∞ part observers (obs\_),
    ∞ sort observers, and
    - oncrete type observers;
  - « part properties:
    - ${\scriptstyle \textcircled{\sc o}}$  unique identifiers:
      - \* unique part identifier observers (uid\_),
      - \* unique part identifier types,

- mereology:
  - \* part mereologies,
  - \* part mereology observers
     (mereo\_);
  - and
- o attributes:
  - \* attribute observers (attr\_) and
  - \* attribute types.

The unique identifier property cannot necessarily be observed:

\* it is an abstract concept and

\* can be objectively "assigned".

That is: **unique identifier**s are not required to be manifest.

- The mereology property also cannot usually be observed: *«* it is also an abstract concept, *»* but can be deduced from careful analysis.
  That is: mereology is not required to be manifest.
- The attributes can be observed:

 $\otimes$  usually by simple physical measurements,

- $\otimes$  or by deduction from (conceptual) facts,
- That is: attributes are usually only "indirectly" manifest.

# **Discrete Endurant Modelling I/II**

Faced with a phenomenon the domain analyser has to decide

- $\bullet$  whether that  ${\sf phenomenon}$  is an  ${\sf entity}$  or not, that is, whether
  - $\circledast$  an endurant or
  - $\circledast \operatorname{\mathsf{a}}$  perdurant or
  - « neither.
- If endurant and if discrete, then whether it is
  - $\otimes$  an atomic part or
  - $\otimes$  a composite part.
- Then the **domain analyser** must decide on its type,

 $\otimes$  whether an abstract type (a sort)

∞ or a **concrete type**, and, if so, which concrete form.

# **Discrete Endurant Modelling II/II**

- Next the unique identifier and the mereology of the part type (e.g., P) must be dealt with:

  - $\otimes$  part mereology types and mereology observer name (mereo\_P).
- Finally the designer must decide on the **part type attribute**s for parts **p**:**P**:
  - $\otimes$  for each such a suitable **attribute type name**, for example,  $A_i$  for suitable *i*,
  - ⊗ a corresponding attribute observer signature, attr\_A<sub>i</sub>:P→A<sub>i</sub>, ⊗ and whether an attribute is considered static or dynamic.

### **End of Lecture 2: Last Session** — Discrete Endurant Entities

#### **Parts**

#### FM 2012 Tutorial, Dines Bjørner, Paris, 28 August 2012



# LONG BREAK