

# **Trust-based Accounting in Grid Systems**

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# Summary

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Distributed systems which offer access to computing resources via the use of the Grid paradigm, it is important to be able to find a suitable "price" which is to be paid for the use of these resources.

In traditional centralised accounting systems in computers, the price is simply a function of the number of resources used, the length of time for which they are occupied and possibly the urgency of the computing task must be completed. In turn this means that the price will depend on the load on the system and the Quality of Service desired by the user.

In Grid systems, where previously unknown users may start to use the system, the situation is slightly more complicated, because it may not be possible to trust all the users to fulfill their part of the contract, i.e. to keep within the limits of resource usage which they have asked to be allocated. Similarly, a user may not necessarily trust a provider to provide correct results or to supply the agreed resources. A just price for resource usage may therefore depend on the risk which the user and the resource provider are willing to run, and to the degree the provider and user trust one another.

In this M.Sc. project, the aim was to develop a system for trust management which can be used to evaluate a reasonable price for use of computing resources, with due consideration for the degree to the trust relationship between the site and user and vice versa.

The fields studied was the Grid architecture for Computational Economy, Beta trust reputation and revenue management. It lead to a model formulation for price calculation based on a trust value and lead to a heuristic for how a site and user should negotiate with each other. An simulation program was implemented

to investigate the price calculation and to find the optimum resource allocation from an economical point of view.

**Keywords:** Grid systems, Economical Grid, reputational trust, revenue management, optimum capacity allocation, price calculation.

# Sammenfatning

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I et distribueret miljø som tilbyder adgang til computer ressourcer vha. Grid paradigmet, er det vigtigt at kunne finde en passende pris, som skal betales for at benytte disse ressourcer.

I traditionelle centraliseret computer accounting systemer, er prisen en simpel funktion af det antal ressource man har brugt, det tidsrum det har været brugt i og måske med hvilken prioritet som opgaven skal gennemføres efter.

I et Grid system, hvor tidligere ukendte brugere kan benytte systemet, er situationerne mere kompleks, da det ikke er sikkert at man kan stole på at alle brugere opfylder del af en kontrakt. I denne kontrakt er f.eks. sat en grænse for hvor meget en ressource må bruges. Tilsvarende kan en bruger ikke altid stole på at en udbyder giver det rigtige resultat eller allokere den mængde ressource som aftalt. En korrekt pris for en ressource kan derfor måske afhænge på den risiko som en bruger og en udbyder er villige til at tage samt til hvilken grad brugeren og udbyderen stoler på hinanden.

I dette projekt har målet været at udvikle et system for at håndtere tillid, som kunne bruges to at lave en fornuftig pris for en computer ressource med hensynstagen til de tillidsforhold som udbyder og bruger har.

Følgende felter er blevet undersøgt: Grid arkitektur for computer økonomi, Beta tillidsrygte og indtjeningsstyring. Dette til en modelformulering for prisberegning baseret på tillidsværdier og førte til heuristikker om hvordan bruger og udbyder skal kommunikere med hinanden. Et simulerings program var lavet for at undersøge pris beregning og for at finde optimalet ressource allokering ud fra et økonomisk synspunkt.

**Nøgleord:** Grid systemer, Økonomisk Grid, Rygte tillid, Indtjeningsstyring, Optimal allokering af kapacitet, pris beregning.

# Preface

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This thesis was carried out in the period of September 1st 2005 to February 28th 2006 at Informatics Mathematical Modelling, the Technical University of Denmark in partial fulfillment of the requirements for acquiring the Master degree in engineering. The thesis was supervised by Robin Sharp.

The thesis deals with different aspects of mathematical modeling of a computational grid using data and partial knowledge about the structure of computational grids. The main focus is on extensions on the economic aspects of computational grids.

Lyngby, Februar 2006

Peter Birkestrøm Due





# Acknowledgments

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# Introduction

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Since the concept of computational grids has emerged, the development has been going fast. The scientific community saw the potential for distributed computing very fast and is leading today in the usage of a Grid system. The object of a computer grid is like a power grid, you plug into a socket and then you have power. In this case you just have computer power ready to be used. In Europe CERN is pushing in an attempt to utilize this idea to the maximum extent in order for them to have en computational power ready when their Large Hadron Collider is done, which is scheduled to be in 2007. They have started the LHC Computing Project(LKC), with the purpose of building and maintaining a data storage and analysis infrastructure for everyone that uses the LHC[Knobloch and L.Robertson, 2005]. Since the projected data use amounts to 15 Petabytes annually and a CPU capacity of approx 25000 cpus [Eck, 2005], the terms trust and accounting is becoming more and more prevalent.

Trust since there will be more and more transactions where the users are previously unknown. This is problem, since there is no guarantee that the user will keep within the limits allocated by the system. When the limit is breached, there has to be a consequence one way or another. One way could be to let the trust be a parameter in the price the user has to pay for the resource usage.

Another issue regarding accounting is the demand for many cpus. The question

is how much of this should governmental funded and how much of this could be handled by private hosting. In order for this to happen there has to be a computational economy in place.

Large corporation like Intel, IBM, Silicon Graphics, Sun and Microsoft all invest large sums into their research in GRID computing. Intel can use the grid to simulate the communications on embedded chips, IBM can see a potential in consolidating cpu power in large centers and Microsoft believe that after a computational grid, there will come a data grid, where their SQL servers will be ready.

All of this requires an infrastructure, that can handle pricing and trust. This thesis will focus on how to derive a "just" price for the computational power in regards to trust.

## 1.1 Motivation

To find a price is not just a question of looking at the costs. There are many other parameters to be considered. First of all where is the trade taking place. From the point of view of thesis, this is done on a computational grid where accounting mechanisms exists but is not used. A study of relevant literature shows that there is only one, who has looked into the economy of resource usage and the chapter detailing this shows that there are still much to be done in this area.

With regards to trust it is all within the realm of authentication and authorization. What is unique in a Grid system is that users and sites initially don't have an established trust relationship and hence it is difficult to determine if the user or the site is trustworthy.

One way of solving this is to use reputational trust and combined this with a direct trust. This is used in this thesis as an input for the price.

In regards to economy there are many ways to calculate a price and in order for this to make sense there has to be a definition of value.

This thesis investigates how to use revenue management as well as bond theory,

since it could be applied on a Grid system. The usage of revenue management is could be used to at make the sites more attractable with the use of discount prices. The revenue management is could also be used to find the optimally allocation of resources on the resource sites. This is investigated if this can be applied to a Grid system.

## 1.2 Structure of the Report

The structure of the report is a follows:

**Chapter 1** is the introduction.

**Chapter 2** introduces the concept of Grid Architecture for Computational Economy (GRACE) and other proposed architectures for the grid.

**Chapter 3** defines the trust notation and analyzes a trust simulator.

**Chapter 4** investigates financial theory with regards to costs and capacity planning. Furthermore revenue management and bond theory is analyzed.

**Chapter 5** presents a pricing model and show process flows for how the users and sites interact. A pricing function for resource usage is also derived.

**Chapter 6** analyzes the simulation requirements and presents the different components in the simulation.

**Chapter 7** presents simulation experiments and introduces the different test cases. The results of the experiments are also presented.

**Chapter 8** presents a discussion of the simulation results and possible future work.

**Chapter 9** concludes on the work presented in this thesis.



# Background: The Grid

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## 2.1 Overview

The Grid is a service for sharing computer power and data storage over the Internet. The aim of is to turn a global network of computers into one vast computational resource and the idea is shown in figure 2.1 on the following page. Since jobs on the Grid normally involves large-scale resource intensive applications, the need for resource management is ever present. But job management is very complex as resources are geographically distributed, heterogeneous, and owned by different organizations, firms etc. who have diverse policies etc. Furthermore different sites all have different access and cost models.

In this chapter two architectures are investigated with the aim of establishing an economic model to be employed on a Grid system.

The analyzed models are:

- Grid Architecture for Computational Economy (GRACE)
- G-commerce



Figure 2.1: Generic view of the Grid

## 2.2 Grid Architecture for Computational Economy

In order to make an economic model, one has to realize that the producers (resource owners) and consumers (resource users) all have different goals, objectives and strategies. In his Ph.d. thesis Rajkumar Buyya proposed [Buyya, 2002] a framework for economy-based distributed resource management and resource scheduling for Grid computing. Buyya points out that there are many applications for the Grid such as Globus<sup>1</sup> and Legion<sup>2</sup>, but none address the issue of computational economy.

In order to make a real world scalable Grid system there is a need for computational economy in order have mechanisms to regulate supply and demand.

Buyaa introduces the following terms for the two major players in GRACE as shown on figure 2.2 on the next page:

**Resource Providers** These are also called Grid Service Providers (GSP), who owns and administrate resources. They seek to maximize their profit

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<sup>1</sup><http://www.globus.org/>

<sup>2</sup><http://www.cs.virginia.edu/legion/>



and consequently their resource utilization.

**Resource Consumers** These are also called Grid Resource Brokers (GRB). The GRB's act as the consumer's software agent. The consumer's strategy is to have its jobs solved in the shortest amount of time within a given time frame. The brokers are also referred to as super-schedulers.

In the Grid environment Buyaa also introduces the Grid Market Directory(GMD), which discovers resources on the grid and handles communication with the Broker. Each resource site has a Grid Trade Server (GTS), that handles current jobs and makes sure that the site is listed in the GMD. Also Deal Template is introduced as a form of contract between the GSP and GRB. This gives a typical market-oriented Grid environment as shown in figure 2.2 It is a prerequisite that there is a bank in the environment, that controls the flow of money.

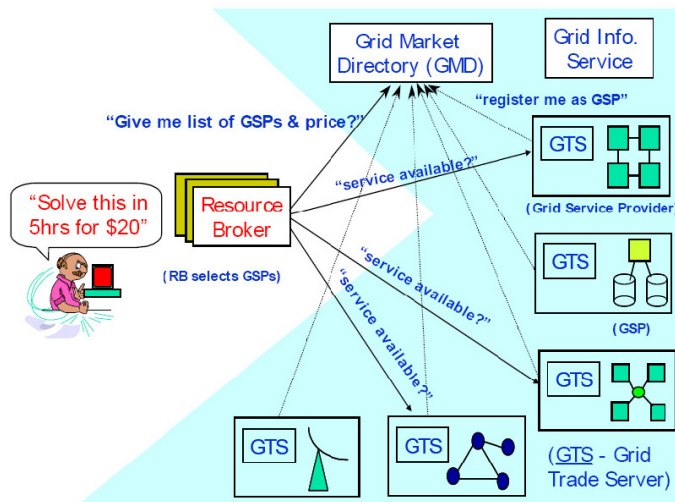


Figure 2.2: Generic view of interaction between players in an Economy grid

The GSP and the consumers are meant to communicate as follows:

1. Client (agent) initiates communication with the trade server at the GSP to obtain the price for resource usage,
2. Negotiation takes place according to the current GSP's pricing scheme,
3. Client and GSP agrees and the trade server requests a credit check in the bank.

4. Job is added to the site's job queue with the state wait.
5. Job is started when Bank returns with a credit check ok.
6. Job starts
7. Job finishes.
8. Resource usage is accounted and money is transfer to the site's account in the bank.

This way no job is started without the user having money.

The broker's job consists of the following:

1. The broker finds GSP's
2. It identifies suitable resources and locates a price (negotiation with the different trade server listed in the GMD).
3. It selects the GSP which matches the utility function and deadlines the user has specified.
4. Job is submitted and money is subtracted from the user's account in the bank.

In a given grid economy the economic model has to be defined for the market. There are many models but the ones relevant to Grid systems are:

- A Commodity Market Model
- Posted Price Model.
- An Auction Model

In a commodity market the resource providers set the price and advertise their prices. The pricing policies can be static or dynamic and the users choose sites after a cost-benefit analysis. The Posted Price Model is the same as the Commodity Market Model except the sites give the price before the scheduling. In the Auction Model the GSP invites bids and stops when there are no more new bids from users.

Which model to choose is a matter of optimization. Section 2.3 on the next page includes there is a comparison between the Commodity and Auction Model.

A fixed pricing scheme is ill suited as it does not support various levels of Quality of Service for the consumers whereas consumer requirements will be biased in terms of runtime, price, etc.

The demand function of the consumers needs to be investigated in order to do a reasonable forecast function. The reason for this is in a population of quality-sensitive buyers all pricing strategies leads to a price equilibrium, where in a population of price sensitive buyers the strategies leads to price wars. Buyaa [Buyya et al., 2001] defines the following items to be included in accounting in order to locate a price:

- CPU usage time
- Memory
- Maximum resident set size - page size
- Amount of memory used
- Page Faults
- Storage used
- Network activity
- Signals received
- Utilized software libraries

Another parameter for the price specification could be time of the day etc. Other parameters can be:

- Consumer id
- Peak timer price
- Lunch time price
- Offpeak time price
- Discount price (not fully loaded)

- High rate (full load)
- Holidays price
- Wall clock time price (within 2 hours, over 2 hours)

Today the price is defined from production cost and profit margins. But the perceived value of resource differ between producer and consumer. The following value function is assumed:

$$\text{Resource value} = f(\text{Physical costs}, \text{Service overhead}, \text{Demand}, \text{The value from the users point of view}, \text{Preferences})$$

This is the price function, but the three last parameters are subjective and valued differently by different users. Preferences is when the job has to be executed etc.

## 2.3 G-commerce

This is another model for resource allocation in a Computational Grid. The model is called G-commerce[Wolski et al., 2001] and is built on the notation of a market economy. It focuses on 2 categories of economic models: Commodities Market and Auctions. The commodity market is used, since the grid tries to use resources as interoperable commodities. The auction market requires no global pricing information and is easy to implement.

The following criteria are used to evaluate the 2 categories:

1. Grid-wide price stability
2. Market equilibrium
3. Application efficiency
4. Resource efficiency

The price stability is required in order to maintain scheduling stability. If the price fluctuates the overall performance will be poor since the scheduling efficiency will decrease. Market equilibrium is a measure of how fair the price

is. However this assumption only holds if the users are quality-sensitive as mentioned in 2.2 on page 6. Application efficiency measures how effective the Grid is as a computational platform. Resource efficiency measures how well the grid utilizes its resources.

In the article by Wolski [Wolski et al., 2001] the following assumption is made:

*The relative worth of a resource is determined by its supply and the demand for it*

In other words it is not the production cost or profit margins alone that defines the value of a resource. Another criteria is the Market equilibrium 2.3, which is a means of price adjustments.

In the simulation section in the article [Wolski et al., 2001] the commodities are CPU and disk storage and the result turns out to be that a commodities market is a more suitable choice than an auction market.

**Market Equilibrium** The market equilibrium is the work of Debreu [Debreu, 1959]. A market economy consists of producers, consumers and supply and demand functions for each commodity, which is determined by a set of market prices for the various commodities. A unique equilibrium price is guaranteed to exist according to the theorem of Debreu [Debreu, 1959]. For all commodities their prices are represented by the price vector  $p = (p_1, p_2 \dots p_n)$ , then the excess demand vector is defined as  $z_j$  for the  $j^t$  commodity. It is assumed that all  $z$ 's are interrelated so that each  $z_j$  is a function of all prices in  $p$ . This means there is an equilibrium point  $p^*$  such that  $z(p^*) = 0$  exists. This also means that for any value of  $p$  the  $n \times n$  matrix can be derived.

$$D_z(P) = \frac{\partial z_i}{\partial p_j} \quad (2.1)$$

This means for any value of  $\lambda$ , which has the same sign as  $D_z(P)$ , the economic equilibrium is obeying the differential equation

$$D_z(P) \frac{dp}{dt} = -\lambda z(p) \quad (2.2)$$

## 2.4 Summary

In this chapter I have shown current status of a economical Grid system. The 2 most dominant ways are the GRACE and G-commerce. GRACE is an architecture for how the process flow should be for an economical Grid system. It also shows which parameters should be accounted for and how they should influence the price. However GRACE does not answer which market condition that apply for a Grid system, but it sums up how the work flows differs compared to the market condition. G-commerce believes that the market containing Grid systems will either be a commodity market or an auction market. Furthermore it is believed that on a market the governed condition is market equilibrium. With this condition it shows that commodity market is more applicable than an auction market.

# Background: Trust Reputation Systems

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## 3.1 Overview

Trust is easy for humans to understand. When somebody behaves good, they are to be trusted whereas when somebody behaves bad the trust is low. In the business world people look at the risk of a economic loss if customers do not pay. The trust relationship is neglected initially. Over time the customer will be granted more and more trust as long as transactions are successful. If the customer ill behaves the firm will not decline to have any further business with that person.

In this chapter the trust notations that apply tot he above mentioned scenario will be investigated. Especially trust reputation systems as mentioned by Jönsang [Audun Jøsang and Faccor, 2003] will be investigated. The trust simulator by M. Brinkløv [Brinkløv and Sharp, 2006] is also investigated.

## 3.2 Trust

### 3.2.1 Notation of Trust

In order to handle trust there is two leading definitions in the literature which is something called *reliability trust* and *decision trust*, and they are defined as follows:

**Reliability Trust** Trust is the subjective probability by which an individual A, expects another individual B, to perform a given action on which B's welfare depends.[Gambetta, 1990]

**Decision Trust** Trust is the extent to which one party is willing to depend on something or somebody in a given situation with a feeling of relative security, even though negative consequences are possible.[Mcknight and Chervany, 1996]

Notice in the first definition there is dependence on the trusted party, reliability of the trusted party whereas in the second also risk attitude is accounted for. You might do an action even though negative consequences are possible. In other words there is a trade-off.

### 3.2.2 Reputational Trust

Since the Internet was invented people have used it to perform business transactions in small and large scale. There are eBay where people sell and buy goods, there is large invoice systems like the one at Novozymes<sup>1</sup> where the customers can buy enzymes on-line. Often these transactions take place where the players often have insufficient information about each other. In order to create trustworthiness you could create a reputation system. Reputation is by definition what is generally said about or believed about a person's or thing's character or standing [Jøsang and Ismail, 2002]. In an e-Commerce environment reputation can be used to generate trust among the participants. Such a system is called a reputation system and its objective is to gather, distribute and aggregate feedback about participants. Audun Jøsang proposes a reputation engine in [Jøsang and Ismail, 2002], which is based on the beta probability density function. The reputation engine calculates the user's reputation ratings from various inputs including feedback from other users.

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<sup>1</sup><http://www.novozymes.com/>



The framework for such a system can be seen in figure 3.1 as described in [Audun Jøsang and Faccar, 2003].

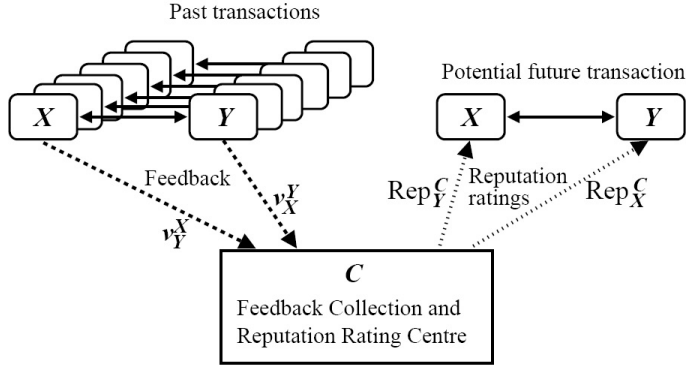


Figure 3.1: General framework for a reputation system

The reputation center collects the input from all agents and updates each agents reputation rating as a function of feedback. These updated ratings are on-line and can be used for agents to decide if the agents wants to transact with another agent. In Jøsangs beta reputation system [Jøsang and Ismail, 2002] the backbone of the system is the beta probability density function, which can be used to represent probability distributions of binary events. In the next paragraph the beta density function is elaborated.

**Beta Density Function** Beta function uses the two 2 parameters  $\alpha$  and  $\beta$ , who can be expressed as a  $\Gamma$  function:

$$f(p|\alpha, \beta) = \frac{\Gamma(\alpha+\beta)}{\Gamma(\alpha)\Gamma(\beta)} p^{\alpha-1} (1-p)^{\beta-1}, \text{ where } 0 \leq p \leq 1, \alpha \succ 0 \text{ and } \beta \succ 0$$

This gives a probability value of the beta distribution:

$$E(p) = \frac{\alpha}{(\alpha + \beta)} \quad (3.1)$$

If you have a process of 2 possible outcomes  $\{x, \bar{x}\}$ , where  $r$  is the observed number of  $x$  and  $s$  is the observed number of  $\bar{x}$ , then the future possible outcome of  $x$  can be expressed as:

$$\alpha = r + 1 \text{ and } \beta = s + 1 \text{ where } r, s \geq 0$$

If you have 11 observations of  $x$  and 7 observations of  $\bar{x}$  it matches a beta function  $f(p|12,8)$  which is plotted in figure 3.2

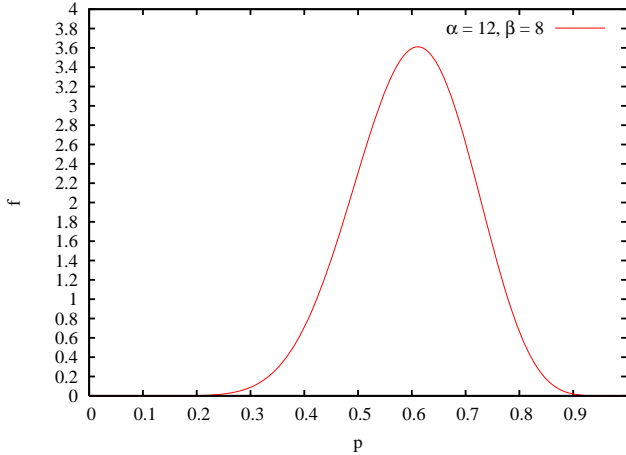


Figure 3.2: Beta function of event  $x$  after 11 observations of  $x$  and 7 observations of  $\bar{x}$

This gives an  $E(p) = 0.6$  which indicates that the outcome of  $x$  is uncertain. Had the value been 0.9 the future outcome of  $x$  would have been high. This is the mathematical background for the reputation system and reputation rating in 3.1 on the preceding page.

**Beta Reputation** Since it is build on the beta distribution, the agent will always have a reputation rating between 0 and 1. The definition of reputation rating is:

**Reputation Rating** Let  $(r(Z),s(Z))$  be  $Z$ 's reputation parameters and let  $(r_{base},s_{base})$  be the a priori parameters.  $r(Z)$  is the positive feedbacks while  $s(Z)$  is the negative. Then the reputation  $Rep(Z)$  will be:

$$Rep(Z) = \frac{r(Z) + r_{base}}{r(Z) + s(Z) + r_{base} + s_{base}} \quad (3.2)$$

This score is for one agent, but in order to be a reputation system it has to be possible to combine more than one rating. To do this the combined feedback from X's and Y's feedback on target Z is defined as follows:

$$r^{X,Y}(Z) = r^X(Z) + r^Y(Z) \tag{3.3}$$

$$s^{X,Y}(Z) = s^X(Z) + s^Y(Z) \tag{3.4}$$

Now the system can combine feedback. But some agents is more highly reputed than other and therefore should carry more weight than others. This is done by discounting the feedback as a function of the reputation of the agent who gives the feedback. To do this an opinion is needed and Jøsang proposes to use Subjective Logic fundamentals [Audun Jøsang and Kinaeter, 2005]. Moreover this can be bijective mapped to the beta distribution. The opinion  $\omega_x^A = (b, d, u, a)$  where  $b=belief, d=disbelief, u=uncertainty$  and  $a=base\ rate$ . At all times  $b + d + u = 1$  and  $a \in [0, 1]$ , which is also used to calculate the expected opinion probability density value determined by  $E(\omega_x^A) = b + au$ .

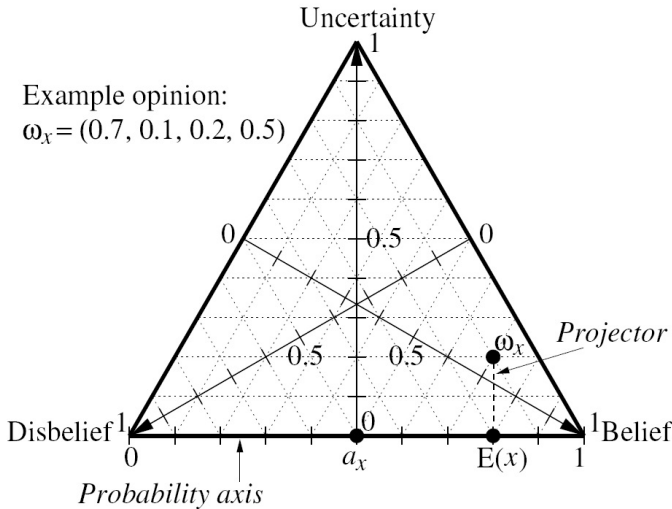


Figure 3.3: Opinion of 7 observations of  $x$  and 1 observations of  $\bar{x}$

The opinion can be used to define belief discounting as described in [Audun Jøsang and Kinaeter, 2005]. The opinion metric is mapped bijective to the beta function defined by:

$$\begin{aligned}
b &= \frac{r}{r+s+2}, \\
d &= \frac{s}{r+s+2}, \\
u &= \frac{2}{r+s+2},
\end{aligned} \tag{3.5}$$

These lead to the following definition of reputation discounting:

**Reputation Discounting**  $x, Y$  and  $T$  are all agents where  $\varphi(p|r_Y^X, s_Y^X)$  is  $Y$ 's reputation function by  $X$ , and  $\varphi(p|r_T^Y, s_T^Y)$  is  $T$ 's reputation by  $Y$ . Let  $\varphi(p|r_T^{X:Y}, s_T^{X:Y})$  be defined as:

1.  $r_T^{X:Y} = \frac{2r_Y^X r_T^Y}{(s_Y^X+2)(r_T^Y+s_T^Y+2)+2r_Y^X}$
2.  $s_T^{X:Y} = \frac{2r_Y^X s_T^Y}{(s_Y^X+2)(r_T^Y+s_T^Y+2)+2r_Y^X}$

It is also called  $T$ 's discounted reputation function by  $X$  through  $Y$ .

This way the system can now handle weigh the different opinions between the agents. The next problem to be solved is that old feedback is not always relevant for the actual reputation rating, since people can change behavior over time. One could go from good to bad and vice versa. The way Jøsang handles this is to let old feedbacks carry less weight than newer feedbacks[Jøsang and Ismail, 2002].

He introduces the feedback longevity factor  $\lambda \in [0, 1]$  which is used this way:  $Z_i$  is the resulting reputation parameters after  $i$  transactions denoted by  $(r_i(Z), s_i(Z))$ . Then after  $i+1$  transaction the parameters is  $(r(Z), s(Z))$  and the longevity is introduced:

$$\begin{aligned}
r_{i+1}(Z) &= ((r_i(Z) + r(Z)) * \lambda \\
s_{i+1}(Z) &= ((s_i(Z) + s(Z)) * \lambda
\end{aligned} \tag{3.6}$$

where  $0 \leq \lambda \leq 1$ . This is also recursive, so it makes it possible to keep reputation parameters on all agents without to store all the feedbacks given.

All of the above mentioned rules will give a feedback system, where it is mathematical formulas, which is the base for the trust ratings. In the next section an implementation of a reputation system is discussed.

### 3.3 Incremental Trust in Grid Computing

The paper by [Brinkløv and Sharp, 2006] defines that trust on a distributed system consists of two components, identity trust and behavioral trust. The identity trust is mainly used for authentication and gives an á priori trust value. The behavioral trust is also called the incremental trust, since it develops as more transactions has taken place. The behavioral trust consist of a direct trust and a reputational trust. The direct trust is agent A's personal experience of B's behavior where the reputation is discounted rating on B from different parties like agent C, etc.

The problem has always been how to combine the direct trust with the reputation trust. According to [Brinkløv and Sharp, 2006] there are two predominant ways of combining the trust: *discrete models* and *fuzzy logic*. The problem with a discrete model is that it uses linear combination of simple numerical values like -1 for total distrust and 1 for absolute trust. However trust is much more of linguistic concept, so it is proposed to use fuzzy logic instead.

The basic idea of fuzzy logic is that for each behavioral parameter, a set of fuzzy sets is defined in terms of so-called membership functions, which specify the degree of membership of the various sets for the possible values of the parameter. When the result of using several parameters  $(x_1, x_2, \dots, x_n)$  has to be combined to a result  $y$ , a set of rules is used of the form:

$$if(x_1 \in S_i \wedge x_2 \in T_j \wedge \dots) then y \in O_k$$

where  $S_i, T_j, \dots, O_k$  are all fuzzy sets and  $\in$  here indicates fuzzy membership. The membership functions  $\mu_k$  for the output fuzzy sets  $O_k, k = 1, \dots, r$  are very commonly given by the so-called min-max formula:

$$\mu_{O_k}(y) = max[min[\mu_{S_k}(x_1), \mu_{T_k}(x_2), \dots]] k = 1, 2, \dots, r \quad (3.7)$$

The actual value is found by defuzzification.

In the simulation in the article various methods are evaluated and the way to combine the trust score is shown in figure 3.4 on the next page

The results show that discrete methods tend to detect change in behavior faster than fuzzy logic, but to find the "right" parameters is much harder. The use of fuzzy logic combined with beta reputation offers a slower detection of behavioral change.

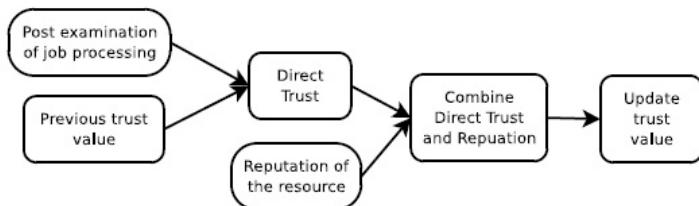


Figure 3.4: Trust score derivation

### 3.4 Summary

In this chapter I have shown what the trust notation can be. I showed how the Beta distribution can be used to build a reputation system as proposed by Jøsang. He proposes how to handle transactions with two possible outcome, success or failure. This is mapped to a Beta distribution, which enables how to predict possible outcomes. He then maps the Beta distribution to opinion containing the following 4 parameters:

1. Belief
2. Disbelief
3. Uncertainty
4. Base rate

From this a reputation system can be build up. Next it is showed how incremental trust can be applied in a Grid system. The incremental trust is a reputation system which is build on among other things the Beta distribution. It is shown that a trust score can be derived by combining a direct trust with a reputation trust using both fuzzy logic and beta reputation.

## CHAPTER 4

# Background: Economic Models

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### 4.1 Overview

This section gives an overview of financial models, that exist today, that can be used in e-Commerce and especially of a Grid system. First of there is the definition of the goods to be sold (assets). Since it is computer power and storage, things are only occupied for a limited period of time, before it can be sold again. This kind of behavior is similar to a hotel with rooms to rent and it is similar to how the tickets are sold for a flight fare. Theories for these industries are evaluated to see if they can apply to an economic model for accounting in a Grid system.

The last section deals with the situation where people buy the option to use a resource within a fixed time frame, but not a specific time. Here the theories of options is evaluated and the black-scholes theorem is investigated.

## 4.2 Perishable Assets

In order to define the goods on the grid I have looked at the following items, which is to be sold/rented. In a Grid system a job consists of CPU time, how many CPU's the job is expected to use, the amount of storage and memory. In other words the goods are:

1. CPU time
2. Disk capacity
3. Memory capacity

Common to all of these are that you normally have a fixed number of them. If you have a cluster you have certain amount of s, memory and storage. In revenue management there is a term called perishable asset in [Weatherford and Bodily, 1992] where one of the main features is that you have a fixed number. Another feature is the demand is uncertain. The following features are all common for perishable assets[Weatherford and Bodily, 1992]:

- A date from where the service or the product is available and a date from where it is no longer available. The service can not be stored unless it is at a high cost.
- There are a fixed number of units. Furthermore these units have a low variable cost, which enable multiple prices for the same unit. This also means it is always better to sell than let it go to waste. The fixed amount of units is often followed with a wide fluctuation on the demand, which shift the balance between supply and demand.
- It is possible to segment the customers into price-sensitive and quantity-sensitive. In the airline industries it is the time of purchase which is used for segmentation. In a Grid system it is the capacity of the resource site which can be segmented.

The aim with perishable assets must be to find the optimal trade off between average price and capacity utilization. The questions to be answered in such a system on a unit-to-unit basis is:

1. How many units should be made available initially at the various price levels?



2. How should this availability of units change over time?

In other words when is the optimal to change between prices for the same commodity. The theory of revenue management tries to answer this.

## 4.3 Revenue Management

Normally you want to sell at the time when the market conditions are most favorable and to the right price. This is a difficult business decision to make. In order to help you with this Revenue Management (RM) can provide an answer [Talluri and Ryzin, 2005], as it addresses three basic categories of decisions:

1. Structural decisions: What selling format to use (posted price, negotiations or auctions).
2. Price decisions: How to set a price over time, how to markdown over the product lifetime.
3. Quantity decisions: When to accept or reject an offer to buy, how to allocate capacity to different segments.

This also leads to a distinction between quantity-based RM and price-based RM. The quantity-based uses capacity-allocation decisions, the price-based uses price as a tool for managing demand.

In order to give an example of industries normally price-based RM is used in retailing and E-business and quantity-based RM is used in airlines.

In relation a Grid system will be quantity-based since it will be controlled by resource sites. However since the aim is to use trust as a parameter for discounting rates, this is price-based RM. The next sections will look into quantity-based RM, since it is highly relevant for the sites in a Grid system.

### 4.3.1 Single Resource Capacity Control

Single resource capacity control addresses the problem of optimally allocating capacity of different classes of demand for a single resource. This is similar to the sale of different fare classes on a single flight or the sale of hotel rooms for a given date. The term classes refers to different prices for the same good.

It is assumed that the firm sells its capacity in  $n$  distinct classes. It is also assumed that the product appeals to distinct and mutually exclusive segments of market. Customers in each class can afford only the class corresponding to their segment. The units of capacity are homogeneous.

The following control mechanisms can be used to control the amount of capacity.

**Booking Limits** Booking limits are used to limit how much capacity can be allocated within every class at any given time. These can either be *partitioned* or *nested*; Partitioning splits the capacity into separate blocks. *Nested* means that the capacity overlaps each other in a hierarchical manner. This implies that the highest classes have access to the entire capacity allocated for the lower classes. For instance if there is 12 units for class 1, 10 units for class 2 and 8 units for class 3 the total capacity is 30. This way you will always have units available for the highest revenue.

For nested booking limits:

- The nested booking limit for class  $j$  is  $b_j$
- Then  $b_j$  is the maximum units of capacity you are willing to sell to classes  $j$  and lower.

In table 4.1 it is shown.

Class 1	Class 2	Class 3
12	10	8
$b_1 = 30$	-	-
-	$b_2 = 18$	-
-	-	$b_3 = 8$

Table 4.1: Nested booking limits

**Protection Limits** Another control mechanism is a protection level. Protection level specifies how much of the capacity you can reserve for a specific class. It can be summarized as follows:

- Protection level for  $j$  is called  $y_j$

- $y_j$  is defined as the reserved capacity for the lower classes  $j, j-1, \dots, 1$  all together.

In table 4.1 the protection levels for class 1 be 12, for class 1+2 be 22 and for alle 3 classes it would be 30. This is also shown in table 4.2 The connection between protection limits and booking limits are:

$$b_j = C - y_{j-1}, \quad j = 2, \dots, n \text{ where } C = \text{capacity}, b_1 = C, \text{ and } y_n = C \quad (4.1)$$

Class 1	Class 2	Class 3
12	10	8
$y_1 = 12$	-	-
-	$y_2 = 22$	-
-	-	$y_3 = 30$

Table 4.2: Nested protection levels

### 4.3.2 Common Requirements

In order to solve optimal capacity control the following requirements have to be met:

1. Allocation of resources only happens if its revenue is greater than the value of the capacity required to satisfy it.
2. The value of capacity should be measured by its displacement cost, which is the expected loss in future revenue form using the capacity now than later.

The first requirement one basically states that you will only sell if you can profit. The second requirement is harder to control, but is important since it dictates that the value of capacity is more than just the cost of production.

### 4.3.3 Static Single-Resource Model

The most common model in revenue management [Talluri and Ryzin, 2005] is the static single-resource model. In order for it to be static the following is assumed:

1. Demand from each class arrives in non-overlapping intervals, but the order is not from low to high revenue.
2. Demand for different classes are independent random variables.
3. Demand for chosen class is not dependent on the capacity controls.
4. It suppresses many details about demand and control mechanism within each period.
5. Group bookings are not considered.

It also assumes risk-neutrality, since maximizing average revenue is most important.

#### 4.3.4 Littlewood's Two-Class Model

This model is a simple two-class model. It is build on Littlewood's work [Talluri and Ryzin, 2005]. It states that there are two classes with associated prices  $p_1 > p_2$ . The capacity is  $C$  and there is no cancellations or group bookings. Demand for each class  $j$  is denoted  $D_j$  and its distribution is  $F_j(\cdot)$ . Demand for class 2 arrives first. The question is now when to change from class 2 to class 1. It could be when to go from cheap price to normal price.

**Solution** Assume there are  $x$  units of capacity left and a request for class 2 arrives. If the request is accepted a unit of class 2 is sold. If the request is denied, then the unit will be sold to price  $p_1$  if and only if the demand for class 1 is  $x$  or higher. In other words:

$$D_1 \geq x \tag{4.2}$$

The expected gain by reserving the  $x^{\text{th}}$  unit of class 1 is<sup>1</sup>  $p_1 P(D_1 \geq x)$ , which is also the expected marginal value. This means that you should accept demands for class 2 as long as:

$$p_2 \geq p_1 P(D_1 \geq x) \tag{4.3}$$

---

<sup>1</sup>P is the probability function

An optimal protection limit exists, denoted  $y_1^*$ , such that class 2 requests are accepted as long the remaining capacity exceeds  $y_1^*$  and the request is rejected if the remaining capacity is  $y_1^*$  or less.  $y_1^*$  satisfies:

$$p_2 < p_1 P(D_1 \geq y_1^*) \text{ and } p_2 \geq p_1 P(D_1 \geq y_1^* + 1) \quad (4.4)$$

If the demand is assumed to follow a continuous distribution  $F_1(x)$  the the optimal protection level  $y_1^*$  is given by:

$$p_2 = p_1 P(D_1 > y_1^*), \text{ equivalent to } y_1^* = F_1^{-1}\left(1 - \frac{p_2}{p_1}\right) \quad (4.5)$$

If the demand  $D_1$  is a normal distribution with mean of  $\mu$  and standard deviation of  $\sigma$  then Littlewood's rule 4.5 gives:

$$F_1(y_1^*) = 1 - \frac{p_2}{p_1} \quad (4.6)$$

This leads to:

$$y_1^* = \mu + z\sigma \quad (4.7)$$

where  $z = \Phi^{-1}(1 - p_2/p_1)$  and  $\Phi(\cdot)^{-1}$  is the inverse cumulative normal function. This also states that if  $p_2/p_1$  is larger than 0.5 the optimal protection limit  $y_1^*$  is less than the mean value.

## 4.4 Hardware Costs

In a Grid system available resources represents a large investment. In order to determine the total investment all costs have to be located [S. Holm-Rasmussen, 2005]. A single server includes the following costs:

- Hardware cost
- Software costs
- Cooling cost
- Power cost
- Depreciation

In addition to this there are fixed costs such as rent, investment in network infrastructure, etc. In accounting this is done this way:

total turnover  
 - variable costs  
 = contribution margin  
 - fixed cost  
 =earnings before interest rate  
 - interest rate cost  
 =earnings before tax  
 - tax  
 =profit

In this thesis the main focus will be within the variable costs and particularly the servers costs, since this includes CPU's, storage capacity and memory.

When a server is bought, the normal procedure in Denmark is to let it depreciate over 3 years. Due to legislative regulations depreciations can be deducted from the tax later in the financial year. At the same time it is a recognition that the value of the server is decreasing due to usage. After three years the value of the server is zero, and you would buy a new one. The most common way of using depreciation in accounting is on an annual basis, so the invested amount is divided into 3 equally large amounts (linear depreciation), which is deducted annually.

Another parameter in the hardware cost is the profit. In order to make a profit, you use a markup percentage [S. Holm-Rasmussen, 2005] which is added on a yearly basis. When doing business there is always the risk of people not paying, which is accounted for by adding a percentage a year. In other words the total cost will be:

$$totalcosts = (initialcost) * (1 + depreciation)^{year} * (1 + markup)^{year} * (1 + risk)^{year} \quad (4.8)$$

This does not take into account how much energy the server uses. This has to be taken into account when calculating the total cost. This also applies to cooling.

## 4.5 Black-Scholes Model

Since large computational tasks easily can take many days to complete, there has to be a mechanism for charging the adequate amount for this task. From a providers point of view the earnings from a job taking days to be better than the displacement cost mentioned in section 4.3.2 on page 25. Seen from the consumers point of view it could be nice to reserve the capacity but not necessarily use it right away. If this was stock trade instead of computer economy one would issue a bond or option the stock. There are two types of options - the call option and the put option.

**Option** Gives the holder right, but not the obligation, to buy or sell a given asset on an agreed price before or on the exercise date<sup>2</sup>

**Call-option** The buyer has the right to buy the underlying asset. The seller has the obligation to sell the underlying asset at the agreed price

**Put-option** The buyer has the right to sell the underlying asset. The seller has the obligation to buy the underlying asset at the agreed price

In the context of a Grid system, the call option should be used for the service between consumer and resource and the put option for the relation between 2 resource sites. One could imagine that a site bought put options in other resource sites, in order to save its own capacity.

In this thesis only the call option will be investigated.

In order to put a price on the option there is a theorem called classical Black-Scholes Model [Frey, 1997] which can be used.

In the Black-Scholes model the market that is considered contains a risky asset, the stock, and a riskless bond  $B$  that is being traded. The price fluctuation on the stock are described by a Stochastic process  $X = (X_t)_{0 \leq t \leq \infty}$ , which is defined in the probability space  $(\Omega, F, P)$  where  $(F_t)_t \geq 0$  and  $B_t = 1$  for all  $t \geq 0$ <sup>2</sup>.

A typical model for price process of the stock is the *generalized Black-Scholes model*, where  $X$  is given by the solution to the following Stochastic differential

<sup>2</sup><http://en.wikipedia.org/wiki/Option>

equation:

$$dX_t = \mu(t, X_t)X_t dt + \sigma(t, X_t)X_t dW_t \quad (4.9)$$

$W$  is a standard Brownian Motion in the same probability space.  $\mu$  and  $\sigma$  are chosen so they are unique solution to 4.9.

The interpretation of 4.9 is that at a given point in time  $\mu(t, X_t)$  describes the instantaneous growth rate of the stock, where the *volatility*  $\sigma(t, X_t)$  measures the instantaneous variance of the return proces  $\ln X$ . In other words is volatility a metric for the risk taken by investing one unit of money into the stock.

As long as  $\sigma(t, X_t)$  is constant or follows at function of time, then model 4.9 is called the *classical Black-Scholes model*.

In order to put a price on the call option you need to know the trading stock price  $S$ , where the option has an exercise price  $K$  (the right to buy the share at price  $K$  at time  $T$  years in the future), the constant interest rate  $r$  and the constant stock volatility  $\sigma$ . This will give a price  $C^3$  given by:

$$C = SN(d_1) - Ke^{-rt}N(d_2) \quad (4.10)$$

where the rest of the parameters are:

**N** The cumulative standard normal distribution

$$d_1 = \frac{\ln(S/K) + (r + \frac{\sigma^2}{2})t}{\sigma\sqrt{t}}$$

$$d_2 = d_1 - \sigma\sqrt{t}$$

$\sigma$  standard deviation of stock returns

In order to locate  $K$  or the strike price, you need to have a forecast function of the stock. This means if the  $K$  is higher than  $S$  on the exercise day, the seller will earn on it. If on the other hand  $K$  is lower than  $S$  then the buyer will benefit from it. In terms of a Grid system the stocks will be CPU's, memory, storage and the strike price will be the production cost. If bonds were traded on a Grid system each site needs to have a forecast system to estimate their cost in the future.

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<sup>3</sup>definition by Black and Scholes at <http://bradley.bradley.edu/~arr/bsm/pg04.html>



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## 4.6 Summary

In this chapter I have shown what a perishable asset is, which is a service. Then I show how revenue management with single resource capacity can allocate optimally their resources according to a rule called Littlewood's rule. I then shown how production cost is handle in accounting and in the last section I elaborate on how bond theory with the theorem of Black-Scholes can be applied to a Grid system.



# Model formulation

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## 5.1 Overview

In this chapter the final parameters in the total price are selected. Furthermore the price formula is derived and the inputs for it are selected. Section 5.3 on page 38 shows the different steps the broker and resource site must perform in order to communicate.

## 5.2 Pricing Scheme for a Grid system

Goods in a Grid system can be defined as perishable assets according to the definition in 4.2 on page 22. On a resource site there are many servers connected in a cluster. This gives the site all its CPU's, all its storage and all its memory. So in order to calculate the price for one job, the price for one server has to be calculated. According to GRACE 2.2 on page 6 there are many parameters to be include in the price. I have selected the ones shown on figure 5.1 on the following page

The parameters are divided into three main categories:

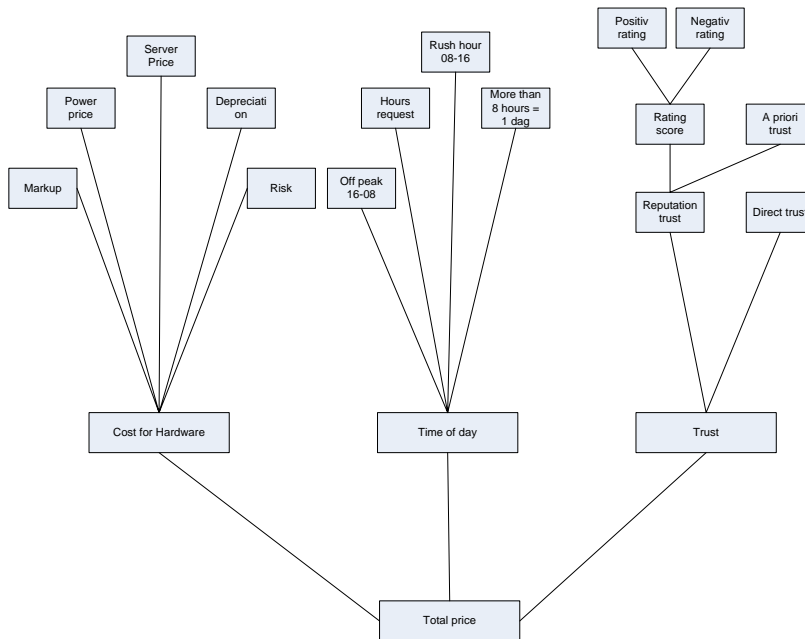


Figure 5.1: Price parameters in the total price

1. Cost for hardware
2. Time of day
3. Trust

The categories are elaborated in the following subsections.

### 5.2.1 Hardware Cost

In this section the hardware cost is analyzed and set. At figure 5.1 cooling is an variable. Cooling is here meant to be the power usage used for the cooling equipment in the different server rooms at the resource sites. Studies from The Electricity Conservation Fund and Technological Institute [S.Ø. Jensen and Viegand, 2004] show that cooling constitutes about 30% of the total power consumption in the

server room. This means that the only input which the resource site can control is the power price. The rest of the variables are business decisions which has to be taken by the managers of resource sites. The variables are then:

Name	Abbreviation	Formula
Purchase price	INVEST	
Power in kWh	PO	
Power usage per day	POU	$24 * PO$
Cooling	COOL	$0,3 * POU$
Depreciation	DEPR	6% p.a
Risk	RISK	10% p.a
Earning	MARKUP	10% p.a
The daily price	DAYCOST	Price per hours for the specific day
Server life time	INVESTPERIOD	1000 days

Notice that the investment period is 1000 days and is a approximation to 3 years which is 1092 days. The number 1000 is easier to work with and that is why it is chosen. The markup is chosen to be 10% and so is the risk. The depreciation is set to 6% which is a normal rate in Denmark. All these percentages are management decisions to take in order for the prices to follow the sites price policies.

In order to find total hardware cost with the current power price the following formula is used:

$$HWCOST = INVEST + (POU + COOL) * INVESTPERIOD \quad (5.1)$$

where the parameters come from 5.2.1.

$$TOTALCOST = HWCOST * (1 + DEPR)^2 * (1 + RISK)^2 * (1 + MARKUP)^2 \quad (5.2)$$

The daily price is found by:

$$DAYCOST = \frac{TOTALCOST}{1000} \quad (5.3)$$

This gives the daily price for day  $j$  with the power price  $PO_j$ . The power price is obtained from the power supply company or in Denmark it can be located on the Nordic Power Trade Market<sup>1</sup>. Since the power price fluctuates almost daily the resource sites have to calculate a new price every day.

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<sup>1</sup>www.nordipool.no

### 5.2.2 Trust in the Price Model

In order to connect price and trust, I will let the trust score from the reputation system be mapped to the risk parameter in the hardware cost. This way any change in the trust score will have immediate effect on the price according to the model 5.2 on the previous page. The risk parameter is set to go from 0 % to 20%, which is chosen since it is not too high a percentage nor too low percentage. If you are highly trusted the risk is almost 0 and if you are not trusted your risk is 20. It can be selected differently according to the sites price policies. Some sites will not conduct business with somebody who is not trusted, whereas other sites will take a higher percentage for doing transaction with one with a low trust score.

In the scenario where the risk fluctuates from 0% and 20 % and the trust score is taken from [Brinkløv and Sharp, 2006], which gives a range from [-1;1] the following mapping is used:

$$RISK = (1 - trustscore) * 10 \quad (5.4)$$

If you look at the formula 5.4 any change in the trust score is amplified by a decade. This means that small changes in trust scores will have a larger impact on the price.

### 5.2.3 Time of day in the Price Model

In order to segment after time, there has to be a standard node. HP Proliant DL380 G4 rack server <sup>2</sup> is chosen as many Danish IT-firms use this for hosting. In this standard node there is 2 CPU's, 4GB memory and 300GB storage.

The jobs on in Grid system contains information about how much CPU time is required, how many CPU's are to be used, memory usage and storage usage. These have to be combined with the daily price. This is done by weighting them equally:

$$\begin{aligned} x * 2CPU + y * 4GB RAM + z * 300GB HDD &= 1 \\ \Updownarrow \\ x * CPU + y * 2GB RAM + z * 150GB HDD &= 1 \end{aligned}$$

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<sup>2</sup>Specs. are available at <http://h18004.www1.hp.com/products/servers/proliantdl380/index.html>

Since they all carry the same weight:

$$\begin{aligned}x * CPU &= \frac{1}{3} \\y * 2GB RAM &= \frac{1}{3} \\z * 150GB HDD &= \frac{1}{3}\end{aligned}$$

Now the parameter  $\alpha$  is introduced as a scalar dependent on the job information on hardware resources:

$$\alpha = \left( \frac{1}{3} * x CPU + \frac{1}{6} * y RAM + \frac{1}{450} * z STORAGE \right) \quad (5.5)$$

From this the hourly rate is deducted as:

$$HourlyRate = \alpha * DAYCOST \quad (5.6)$$

The hourly rate only holds for one day or until the power price changes. If a job request is longer than 8 hours the user has to buy an option for the job and the price for that is calculated using Black-Scholes from 4.10 on page 30. The price will rise since you have to pay for the option and for the usage of the resource.

### 5.2.4 Discount Pricing

To attract jobs to sites, which otherwise not receive many jobs, the site can use discounting as a strategy. It makes sense since the units have low variable cost and since the market is homogeneous where the consumers only choice of comparison between sites is the to use prices. If the site have a lower price than the rest, then this site will raise its probability of receiving jobs. If the resource sites have a discount price and normal price for a capacity  $C$  of goods, then according to Littlewood's rule 4.3.4 on page 26 the optimum protection level is

$$F_1(y_1^*) = 1 - \frac{p2}{p1} \quad (5.7)$$

Since it is assumed the demand follows a normal distribution this equals to:

$$y_1^* = \mu + z\sigma \text{ where } z = \Phi^{-1}(1 - p2/p1) \quad (5.8)$$

Since discount often is taken from the end price, the discount price is  $Dayprice - (discountpercentage * dayprice)$ , which leads to the following prices:

- $p2 = (1 - discountpercentage) * dayprice$

- $p1 = \text{dayprice}$

Since  $z$  in the protection limit comes from the fraction of  $p2/p1$ , this leads to:

$$\begin{aligned}
 z &= \Phi^{-1}\left(1 - \frac{(1 - \text{discountpercentage}) * \text{dayprice}}{\text{dayprice}}\right) \\
 \Downarrow \\
 z &= \Phi^{-1}(\text{discountpercentage})
 \end{aligned}
 \tag{5.9}$$

Then for the demand  $D$  with mean  $\mu$  and standard deviation  $\sigma$  and the discount price is dependent of the dayprice in 5.3 on page 35 the protection limit is:

$$y_1^* = \mu + \Phi^{-1}(\text{discountpercentage})\sigma \tag{5.10}$$

How to set  $C$  it's up to the site managers, but a good indicator is the Demand Factor, which is the ratio of the total demand to capacity. If this is less than one, the capacity is a not sufficient for very long.

### 5.3 Process Flows for Job Submission in an Economy Grid

Another part of the model formulation of a just price on in Grid system is how the process flow should be. GRACE 2.2 on page 6 gives a good indication of how this could be done. There are the point of view from the consumer and the point of view from the resource site. What is new it's the trust evaluation compared to the methodology in GRACE. In this system there has be an entity, which has absolute trust from both the consumer and the resource site, in order to introduce money as a concept. The entity will be referred to as a bank, but it is the whole monetary system it refers to.

Figure 5.3 on the next page shows the process the user has go through in order to submit a job. The user will accept the offer from the site if the user thinks it it the best offer at the time. In the end of the process the trust between the site and user is evaluated and money is transferred from the user to the site.

In figure 5.4 on page 41 the site is asked to give a price for a job. But before the site calculates the price the user has to pass three checks:



1. Credit check. There is no need to use time on a user who does not have any money.
2. Is the current users trust score above the sites threshold? In other words does the site wish to conduct business with a user with the current trust.

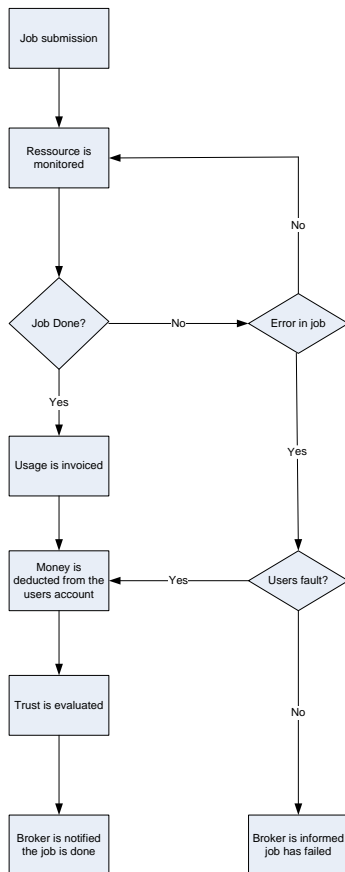


Figure 5.2: The process of a job done at the site.

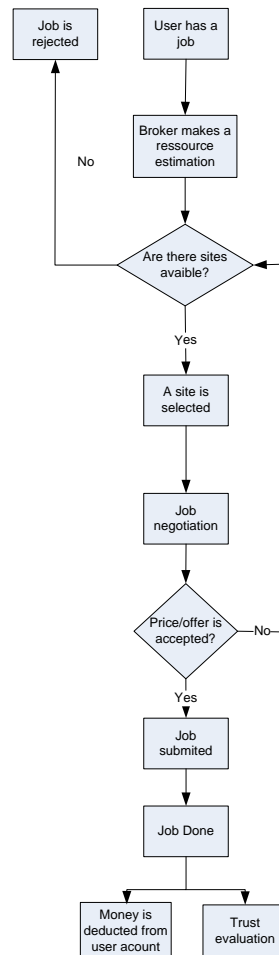


Figure 5.3: The process of a users job submission.

3. Does the site have enough available resources to meet the demands in the user's job? This is build on the assumption that the sites are fair and will not take jobs when they don't have enough resources.

If all three checks add up, the site is willing to do the job. In the last figure 5.2 on the previous page the site is deducting the money for a job done. If there is an error and the job is not completed, the site will decide whether it is a technical mistake on the site or a user mistake. Dependent on the outcome the following would happen: If it is the user's mistake the money will be deducted, and if it is the site's mistake the job would re-run on the site.

## 5.4 Summary

In this chapter I derived my price model and I derived how the user and sites interact with each other. In the section Pricing scheme for a Grid system I derive my calculation model for prices and I show how discount prices can be handled. In the process flows I show the different scenarios to be handled.

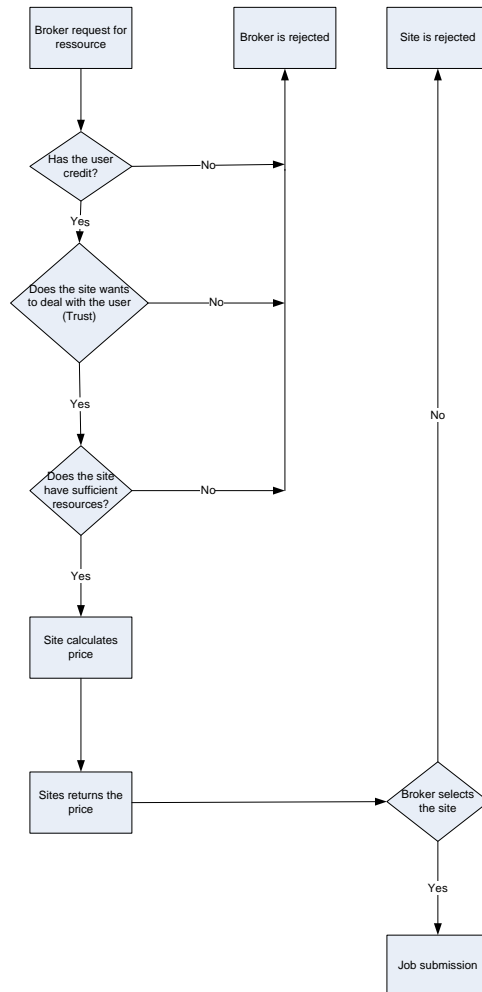


Figure 5.4: The process of a users job submission to a site from the sites point of view.



## CHAPTER 6

# Simulation

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## 6.1 Overview

In this section the presumptions and delimitations for the simulation is defined. The market behavior is defined as well as the user and site behavior. In the two last sections two test cases are formulated.

## 6.2 Analysis

### 6.2.1 Presumptions and Delimitations

Since the simulation has to simulate a economic computer grid which is highly complex, I have made certain assumptions and delimitations for the simulation.

The following items apply:

- The market is a competitive market, since there is little entry cost to join

the market. In theory everybody can offer their computational resources to a Grid system.

- Since the goods are computational power and specific CPU's, it is also a homogeneous market.
- The depreciation model is linear and not exponential. I selected this since it is easier to handle.
- In the simulation the sites are not allowed to overbook. This ensures that all jobs will be executed on some site.
- In the simulation only CPU time requests and request for CPU's is investigated. The cross elasticity between CPU, memory and storage is not simulated.
- The user cannot cancel a job once it is submitted to the sites.
- The general rule of optimization is price and not time.
- The I/O operations and the network performance is assumed to be perfect in this simulation. In the real world this may be a major parameter.
- The price elasticity is perfectly inelastic, so any change price will not have an effect on the quantity
- The cross elasticity for CPU, memory and storage is defined by 5.5 on page 37.

## 6.2.2 Market Behavior

To determine the market behavior, the Nordugrid was investigated for demand and supply. The log files from January 2005 to November 2005 was obtained and from that the demand for job requests and the distribution of the jobs arrival was found.

The log files from the Nordugrid contained 187724 successful jobs, but there was no logging of memory used, CPU time used and how many CPU's used. In order to get this the end time of a job was deducted from the start time which I use a CPU time. This is not the exactly CPU time, but an approximation to it. I believe the approximation is appropriate.

The data showed that demand arrived with a Poisson distribution whereas the request for CPU time followed a normal distribution.

In order to see money flow in the system and the market a bank is introduced. The banks purpose is to check creditability of the users, to keep track of accounts with money and to show transaction history. It does not have any influence on the market with the traditional financial instruments such as to raise the short-termed interest rate or the long-termed interest rate.

### 6.2.3 Derivation of Demand

To derive the demand the data set from the Nordugrid was analyzed. The data indicated that the arrivals between job requests followed Poisson distribution. The arrivals was derived by subtracting one jobs start time from the next job start time. The result can be seen in figure 6.1

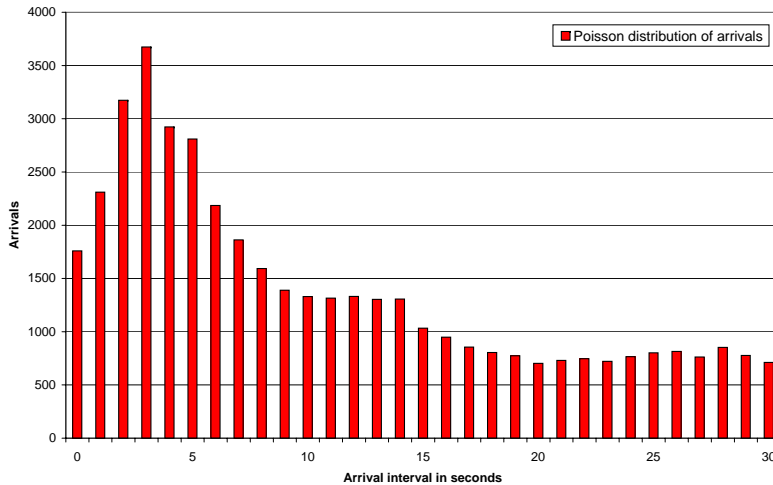


Figure 6.1: CPU request arrival interval in seconds taken from the NorduGrid

Since Poisson distribution is used to model the number of events occurring within a given time interval this model was examined to see if it match the data. After deriving the mean from the sample from 0 to 30 seconds, where the mean is 10,73 seconds this was used as  $\lambda$  in the Poisson distribution. The derived Poisson distribution shown in equation 6.1 where  $X$  is the arrival time and its graph is shown in figure 6.2 on page 47.

$$p(X, \lambda) = \frac{e^{-\lambda} \lambda^x}{x!} \text{ for } X = 0, 1, 2, \dots, 30 \quad (6.1)$$

In order for it to be a Poisson distribution the samples have to satisfy 3 criteria [J.D. Petrucelli and M.Chen, 1999]. To state those criteria it is assumed that  $\lambda$  equals the mean number of counts in the entire interval. Next it is assumed that the entire interval equals  $M$  units. The last assumption is that the entire interval can be partitioned into subunits of size  $\delta$ . The criteria is then:

1. The counts in disjoint units are independent.
2.  $\frac{1}{\delta}P(\text{There is exactly 1 count in } U(\delta)) \rightarrow \frac{\lambda}{M}$  as  $\delta \rightarrow 0$
3.  $\frac{1}{\delta}P(\text{There is more 1 count in } U(\delta)) \rightarrow 0$  as  $\delta \rightarrow 0$

In the data set the entire time interval is 30 seconds with a mean of 1485 counts and a mean number of counts per second to be  $0,696^{-3}$ . It is assumed that the 30 second interval can be partitioned into less and less units. Then the criteria is:

1. The counts in disjoint units are independent.
2.  $\frac{1}{\delta}P(\text{There is exactly 1 count in } U(\delta)) \rightarrow \frac{1485}{30} = 0,696^{-3}$  as  $\delta \rightarrow 0$
3.  $\frac{1}{\delta}P(\text{There is more 1 count in } U(\delta)) \rightarrow 0$  as  $\delta \rightarrow 0$

This shows that the Poisson distribution can be used to model the arrival time in the simulation.

## 6.2.4 Request for CPU time and CPU's

When I displayed this on a graph it showed that the request from 0 to 5 minutes followed a normal distribution which can be seen in figure 6.3 on page 48.

The statistical analysis showed that the requests follow a normal distribution with  $\mu = 156,75$  seconds and a variance of  $\sigma^2 = 3198,633$  and is shown in figure 6.4 on page 48.

Since there was no data on CPU requirements, I had to choose this parameter. I have chosen it to resemble the distribution for CPU time requests. However I have selected a lower mean to be  $\frac{1}{10}$  of the mean of CPU time and setting a variance of 25. This can be seen in figure 6.5 on page 48.



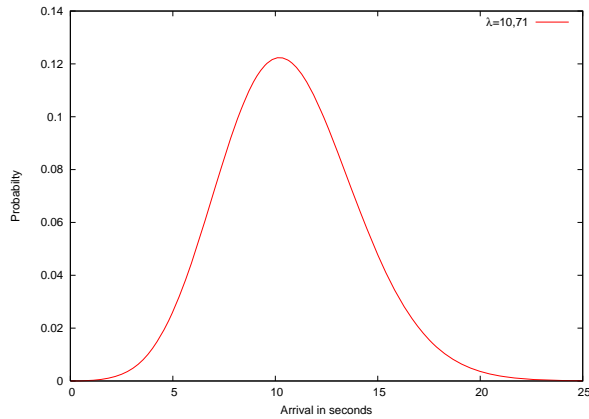


Figure 6.2: Poisson distribution with  $\lambda = 10,71$  used to model the users requests for CPU time

### 6.2.5 User behavior

The users are assumed to be myopic customers [Talluri and Ryzin, 2005], which means they will buy the goods as soon as the offered price is below their willingness to pay. Their willingness is defined to be if they have money on their account in the bank. However the user still checks the price with other sites in order to get the cheapest price. If the user encounters a site with a lower price, it will select this site instead as long as the site in question have capacity to handle the job. By definition the cheapest site for the user will be the site with least occupied CPU's and where the trust score between the user and the site is best and where the sites is trustworthy. The users do not have any preference for if which site is most trusted just along as the site is trusted. Summing up the cardinal rule for the users behavior is:

1. Buy at the cheapest price from a resource that is trustworthy.

### 6.2.6 Site behavior

The cardinal rule for the site is:

1. Capacity should be allocated to a request if and only if its revenue is

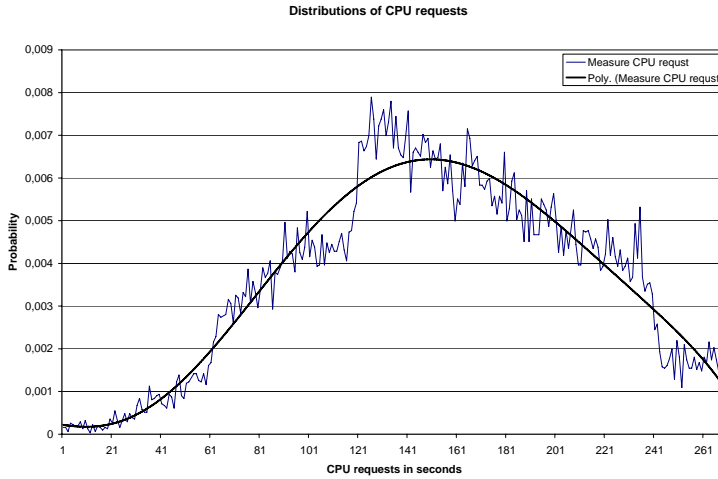


Figure 6.3: CPU request from the NorduGrid.

greater than the capacity cost

Moreover a site is interested in having the highest profit and not particularly the highest amount of jobs. In order for sites to attract jobs they can have a discount price for a certain amount of the capacity. The amount is governed by the site. How to change between discount price and normal price has to be

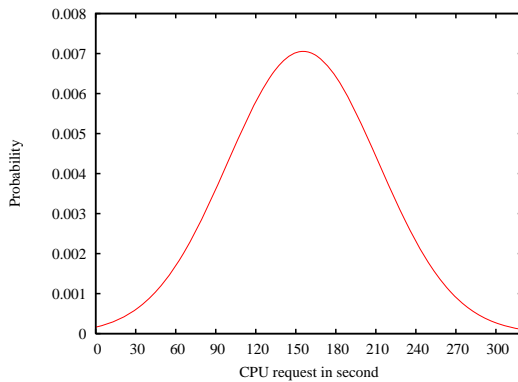


Figure 6.4: Request for CPU time with  $\mu = 156,75$  and  $\sigma = 56.56$ .

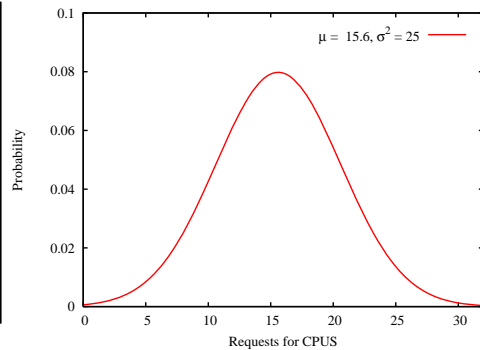


Figure 6.5: Requests for CPU with  $\mu = 15.6$  and  $\sigma = 5$ .

done dynamically.

Furthermore the site calculates the price to a user from the users trust score. If a user's trust score is below a threshold value determined by the site, the site will raise the risk parameter to a high percentage. That way business will still be conducted but a high price.

### 6.2.7 Power Market Behavior

The behavior of the power market is modeled after the Danish Power Market. The power prices fluctuate according to the demand and supply on the power market, where a Grid system is located. I obtained the monthly power price from 2003-2005 for the usage of 4000kWh in Denmark<sup>1</sup> which has to be piped into the trust simulator. A normal server uses approximately 3500 kWh a year, the figures (see figure 6.6) from Danish Energi can be used.

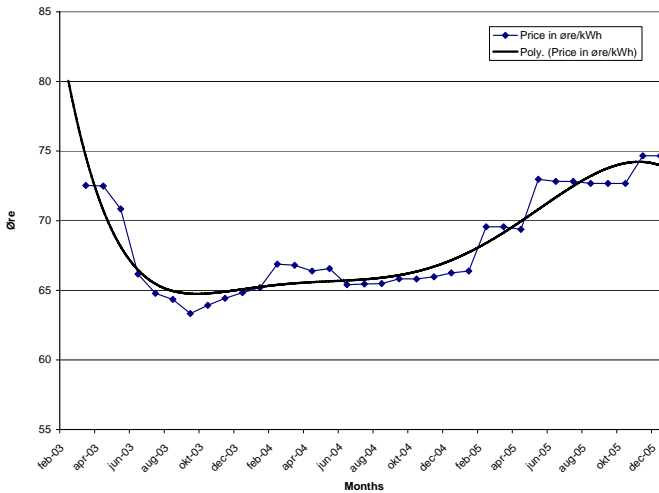


Figure 6.6: The price fluctuation of kWh on the danish power market.

<sup>1</sup>Available at <http://www.danskenergi.dk/webtech/statistik.nsf/fWEB?ReadForm&Load=KJEN-5B8F7R>

## 6.3 Design

### 6.3.1 The Trust Simulator

In order to run the simulations a framework is provided by Michael Brinkløv [Brinkløv, 2005], which is the base of my simulations. His system is a discrete event simulator designed to examine the progression of trust and managing this in a grid system.

The requirements for the Trust simulator is to keep track and provide trust scores for sites and users. Furthermore the system has to handle beta reputation.

### 6.3.2 Elements extending the Trust Simulator

In order to use the Trust simulator for economic simulation the following classes has to be added:

- Pricechange event. This is used when power prices changes. See the appendix D.1 on page 99 for source code.
- Salary event. This is used to simulate that the user receives a monthly salary or an annual budget. See the appendix D.2 on page 100 for source code
- Creditrejected event. Used to check if the user has money or not. See the appendix D.3 on page 102 for source code
- bank. Controls all money flow. See the appendix D.4 on page 103 for source code

The following elements has to modified:

- Modifications to resource user. This is particularly how to select a site for the users job. See the appendix D.5 on page 107 for source code
- Modifications to resource site. The resource site gets some more methods to calculate prices. See the appendix D.6 on page 116 for source code

- Modifications to configuration class. Since economy is introduced, some new data structures are made. See the appendix D.7 on page 127 for source code

In the following sections the above mentioned changes are elaborated.

### 6.3.3 Pricechange event

An event is something that happens at a separate time in the simulator. The event is made from the input file "price.txt" which contains a list of power prices. It also contains the time for the power price to change. When the event is triggered in the simulation it runs through the array of sites and asks them to update their prices with the method *update\_price(price)* where the argument *price* comes from "price.txt". The site updates their internal power price, with *price*.

### 6.3.4 Salary event

This event is to simulate the user receiving money. The reason for this is that firms, universities have a budget for each year and receive funding for a period of years. When the event is triggered it asks the Bank entity in the simulation to update all the users' accounts with the amount specified in the "salary.txt". The text file contains a point in time where an amount is added.

### 6.3.5 Creditrejected event

This event is used for tracking how many jobs that have been rejected by the site, due to the user having no money. It also shows if the rejection is due to the site not having capacity to meet the user's job request.

### 6.3.6 Bank

The bank should keep track of sites and users' accounts. It is a single object which is created when the simulation starts and is created by the configuration object. The configuration is inherited by all the users and sites. In order to

manipulate the bank, there are dummy function in the configuration file, that mappes the functions in the bank.

**Data Elements** The data elements in the bank are:

- A list of accounts

Every resource needs to have an account in order to track the money flow.

- A list of jobs awarded

Since every job interacts with the bank, the bank can count the amount of jobs given to the sites. This makes sense when testing if the site with the most jobs also is the site with the highest revenue.

- A start amount for accounts.

Every user account starts of with an initial amount, while no sites have money on their account. The idea is that site accounts grows as the user accounts decreases.

**Methods** The following methods are needed for the bank:

- To put a salary on a user account

Since the system operates with salary at a given time to the user, the account need to be modified with the amount.

- Transfer money from the user account to the site account.

When the job is done, the site needs to be paid. The site uses this function in the bank where the price deducted is the price informed to the user.

- To setup the bank

When the bank is initialized all user account are created with the amount specified in the bank section of the configuration file 6.3.9 on page 55 and the sites' accounts are created with the amount of zero.

- To check the credibility of the user

This method looks at the user account. If the account exists and there is more than zero money in the account it will return true. Else the user has no money and it returns false.

### 6.3.7 Modifications to Resource User

In the resourceuser class the major changes were in the method "submit\_job", where the user now has to select the cheapest site. Since the only method that is changed is "submit\_job" I will only elaborate this.

The first two changes is how the CPU time and CPU's are requested. Now the request for CPU time and CPU requests are taken from a random normal distribution with a  $\mu$  and  $\sigma$  defined in the configuration file. The random generator is taken from GNU Scientific Library <sup>2</sup>.

The next major change is how the user selects the site.

The users will ask all sites for a price where the sites will return the price based on the current trust score between the site and user. The user will then take the site with the cheapest price. In pseudo code that can be seen in 6.1 on page 58. The behavior for the user if it encounters a site, which behaves badly, is not award the site with a job. The extreme situation is all sites are ill behaving, which will lead to the user refusing to send any jobs to any sites.

### 6.3.8 Modifications to Resource Site

In the existing resource site class, the concept of cost had to be introduced. In the configuration file 6.3.9 on page 55 the Economy section specifies the markup, depreciation, discount, protection level and the hardware cost. In order for the sites to know this, they all have protected variables which is set in

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<sup>2</sup>See the documentation at <http://www.gnu.org/software/gsl/>

the constructor, when the site object is created. Another major change in the constructor that calculates the protection level for the capacity as mentioned in 4.3.4 on page 26.

**methods** The following functions needs to be added to the resourcesite:

- To calculate the protection level for the given capacity

When discount is specified in the configuration file for the site, the protection level has to be calculated. The method calculate the  $z$  of the formula in 4.7 on page 27 where it returns the inverse cumulative normal distribution for the discount rate and returns the protection level. This is used to find the booking level, which is the *capacity - protectionlevel*.

- To answer if the site is fully booked.

When a user asks for a site, he first asks if the site is full. This method returns true if the sites private variable "current available CPU's" deducted the users CPU request is less than zero.

- To update the power price, when the power price changes.

When the pricechangeevent happens this method is called. It takes the price specified in the events and coping it to the site's own private power variable.

- To calculate the price for a job

This function calculates the users price for a job. It takes the user and his request for CPU's. First of it is checked if the user and site have ever conducted business together.

If this is not the case, the site assumes that the user has a neutral trust, which is the trust value of 0. If the user and the site have conducted the business, the trust score is calculated and used. The site then has to check if it finds the user trustworthy according to the sites threshold for trust. If the user is not trustworthy the site will still conduct business, but now the site has to guard itself from a economical loss. It should handle this by raising the risk parameter



to a high percentage such as 80%. However if the site trust the user the the risk is calculated as specified in 5.4 on page 36.

The total cost has to be calculated and an internal price, which is the production cost at the given time. When the internal price is known the maximum discount has to be calculated in order to make sure that no discount given is below the maximum. It it meant to be a safety valve, which will, if triggered, raise the price to 10 times the normal level. This way the user will not select this site, as it is to expensive. And this way the site will not sell with loss.

When discounting is used, there has to be a check that the site does not sell above the booking limit. The psedo code can be seen in tabel 6.2 on page 59.

- To calculate the price for a job during the call "add\_job"

In the method "add\_job" the private variable "cpuasked" is set from the Job object. The method creates an entity object (see figure 6.7 on the following page) if it has not been created. Next it calculates the price for that specific user with the site to user trust score as an argument. Then in the end the method checks that the site will not be overbooked if it accepts the job. If not the counter of available CPU's is decreased and the counter of occupied CPU's is increased.

- To write out the price in method "job\_done" and deduct money from the user in the bank

When the job is done at the site, the price history is updated with a time stamp and the price, the user has used. This prices is also used to deduct the correct amount from the user's account to the site's account.

### 6.3.9 Modifications to Configuration Class

In configuration file there has to be a new economy data structures and there has to be methods to handle the communication between the sites/user and the bank.

## 6.4 Implementation

The Trust Simulator is implemented in C++ and my modifications are also made in C++ and can be seen in appendix D on page 99. The event simulator follows the flow chart on figure 6.7:

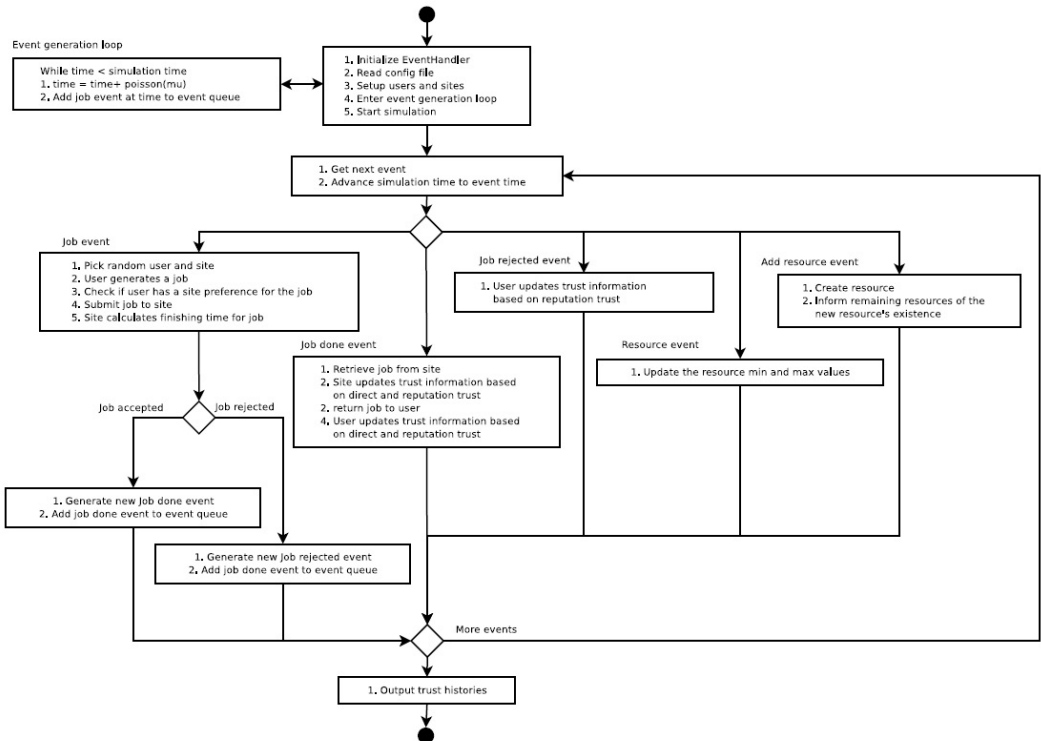


Figure 6.7: Trust simulator flow chart

This system is modified to simulate grid economy based on trust.

Since it is a discrete system it models a system by a representation of state variables that changes instantaneously at separate points in time. The simulator is designed as an object-oriented application, so all objects are inherited. In order for me to extend it, I have to add some new objects and modify existing objects. Now the sites have to calculate a price for each job and the users have to make a conscious decision about which site to use. The system should also handle external input from the world in terms of changing power prices. To

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simulate an economy there has to be a bank, which handles monetary issues. The main properties of the trust simulator is not changed in the sense that it will still be a discrete system.

I added a new structure to the configuration files. It is called economy and contains all the relevant parameters for each site. Here is the markup rate, the depreciation, the hardware cost etc defined in.

The rest of the changes has been implemented according to the specifications of section 6.3.2 on page 50 and the source codes can be seen in D on page 99.

## 6.5 Summary

In this chapter I have shown what it will require to extend the Trust simulator to handle economy according to the Model formulation chapter. I have defined the behavior of the players in the simulation and what component that needed to be made in order to extend the Trust simulator. In the end I show how it is implemented.

Listing 6.1: Pseudo code for how a user selects a site to submit a job

```
1
2     The site id given randomly assumes to be the cheapest.
3
4     Site id is copied to a temporary site rs
5     Price is calculated for site id and put in a temporary
6     price
7     for{siteiterator=start of site array; siteiterator!=end of
8     array; siteiterator++}
9     {
10        Select site from siteiterator.
11        if(site have CPU capacity)
12            if(current sites price is cheaper than
13                temporary price)
14                check the users trust to the site
15
16                If(first time talking to site)
17                    assume site has a trust value of a stranger
18
19                else(not first time communicating)
20                    temporary trust = the current trust between
21                    user and site
22
23                if(temporary trust is above the current
24                    trust threshold for the user=
25                    this is now the cheapest site
26                    site id=temporary id
27                    price=temporary price
28        }
29        Check to see if the for loop has found a cheaper site than
30        the randomly site?
31        if(random site and user have communicated before)
32        check trust against threshold
33
34        site found.
```

Listing 6.2: Pseudo code for how a site calculates the price

```

1 function calculatesite price(user , request for CPU's)
2 {
3   if(first time communicating with site)
4     trust=the trust for stranger
5   else
6     trust=trust value for user
7
8
9   risk=trust mapped to risk
10
11  if(site trust to the user is above the current treshold)
12    risk=80
13
14  calculate hardware cost
15  calculate total cost
16  calculate the internal cost for day
17  calculate the maximum discount
18
19      if(the site has a discount rate and this is below maximum
20         discount)
21      {
22          if(the current occupied CPU's + the amount
23             requested by the user < the booking limit)
24          {
25              user price=day price-discount
26          }
27          else
28              if(the discount is higher than maximum discount)
29                  user price=10 * user price
30      }
31  Return the user price
32
33 }
```

Listing 6.3: Economy struct from Configuration.cpp

```

1 //Parameters to calculate prices for the sites
2 struct economy{
3     unsigned int markup;
4     unsigned int depreciation;
5     unsigned int discount;
6     unsigned int hardwarecost;
7     double protectionlevel;
8     unsigned int sitecpus;
9 };
10
11 //map of parameters and sites.
12 typedef map<string ,economy> ecmmap;
```



## CHAPTER 7

# Simulation Experiments

---

## 7.1 Overview

In the first two sections I define my test cases. In each of these sections the different setups for testing the simulator is elaborated. In the last two section the results of all the tests are presented and elaborated.

## 7.2 Test case 1: Price is Dependent on Trust

The main idea in this part of the simulation is to show the correlation between the users trust score and the price the site will offer to the user. In other words there is a mapping between  $\Delta trust$  and  $\Delta price$ . The test was done with 10 repeats with different seeds. Seed going from 100 and incremented by 10 every time. The number of sites and users is never higher than 100 for each.

### **7.2.1 Price Variation due to Different Price Parameters**

### **7.2.2 Test Case 1: Experiment 1 - Price Change due to Trust**

In this setup the power price is kept constant and all sites have the same economic parameters. All that can influence the price is the site-user trust score.

### **7.2.3 Test Case 1: Experiment 2 - Price Change due to Trust with Varying Power Prices**

In this setup the power price follows the power consumption data and all sites have the same economic parameters. All that can influence the price is the site-user trust score.

### **7.2.4 Test Case 1: Experiment 3 - Different Prices due to Different Parameters on the Sites**

In this setup the power price follows the data found in and all sites have the different economic parameters. The variable input is still the site-users trust score.

### **7.2.5 Test Case 1: Experiment 4 - How Threshold Affects Site Capacity and User's Price**

In the setup both the trust from site to user and vice versa is tested. First the situation were a site is ill behaving and the users refusing to send jobs to this site. The a user will behave badly and his price will rise at the site to a very high price.



## **7.3 Test Case 2: Site have 2 prices and Switches Dynamically Between Them**

The main objective of this simulation is for sites to dynamically switch between a discount price and normal price based on the capacity. Furthermore to check that the user selects the cheapest site available. Now the sites have to handle request for CPU time as well as CPU's. The requests follow the distribution from 6.2.4 on page 46. In the end the sites with the right capacity and usage of discount prices will have the highest earning.

### **7.3.1 Test case 2: Experiment 5 - User Selects Cheapest Site**

In this test it is checked that users selects the cheapest site. The setup starts out with the sites all behaves good and then one site choses to behave bad. The specific site is also the cheapest site, since it is the only one that gives discount. When the user realizes this, he stops to send jobs to the specific site and instead of send it to other sites even though it is more expensive. If in the end there is no sites the user trust, the system stops, since then all sites are untrustworthy.

### **7.3.2 Test Case 2: Experiment 6 - Site Calculates Protectionlevel**

In this setup some sites have discount prices and calculates dynamically the prices. For this test only 3 sites have discount prices and the discount percentages is 10%, 20% and 30%. The capacity is set to 300 CPU's and users can ask all sites for price. Furthermore the users have enough money in the bank to buy.

### **7.3.3 Test Case 2: Experiment 7 - Site Checks that Discount in not Below Cost Price**

In this setup it it verified that no matter how high the discount percentage is set to, it can never be higher than the production cost.

### **7.3.4 Test Case 2: Experiment 8- Site Capacity is Limited**

In this setup the capacity is set low, in order to generate many overbookings. This way the site with the highest revenue is the site who handles the most jobs.

### **7.3.5 Test Case 2: Experiment 9 - Users can Only Select Few Sites**

In this setup the user cannot ask all sites, but only a certain amount of sites.

## 7.4 Results From the Simulation

The sections containing the results from test case 1 and test case 2 are presented.

### 7.5 Results from Test Case 1

The main objective in this test case is to test the price calculation in a Grid system. The price calculation is dependent on user's trust scores and energy prices. The different setups shows how the price changes when a user is ill behaving and vice versa. It is also tested that user's always selects the cheapest site to submit their jobs.

#### 7.5.1 Results from Experiment 1

- Aim: To show price change due to trust

The experiment was carried out with 10 sites and 10 users and repeated 10 times with different random seeds to test the stability of the simulator. User zero acts nice in the beginning but changes behavior and becomes bad for a short while. Then in the end User zero act good again. There is no threshold set for the trust score and it is assumed that strangers have the trust score of 0, which is neutral. As seen on figure 7.1 and at figure 7.2 the price rises when the trust decreases.

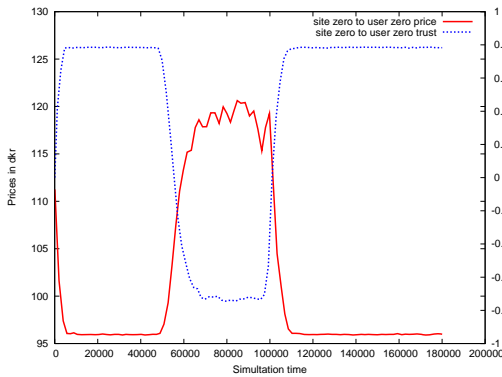


Figure 7.1: The price vs trust for user zero on site zero.

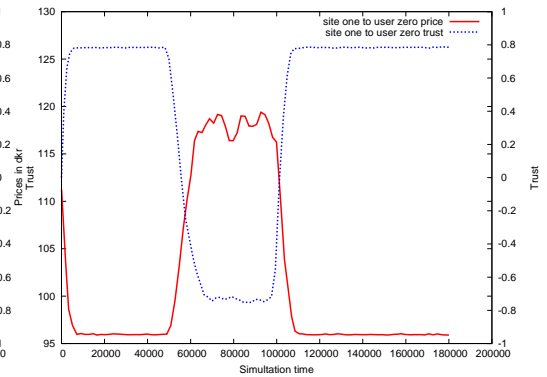


Figure 7.2: The price vs trust for user zero on site one.

The price does so at the same rate as the trust. What is more interesting is when the user is bad, there is a longevity factor that makes the system forget the user's ill behavior. If you look at the price when it is highest, it fluctuates more than when it is at its lowest. The reason why the changes in price are more visible is the mapping of trust to risk (equation 5.4 on page 36 applies) is an amplification by a factor of ten. In figure 7.3 and figure 7.4 the trust of User zero is seen from

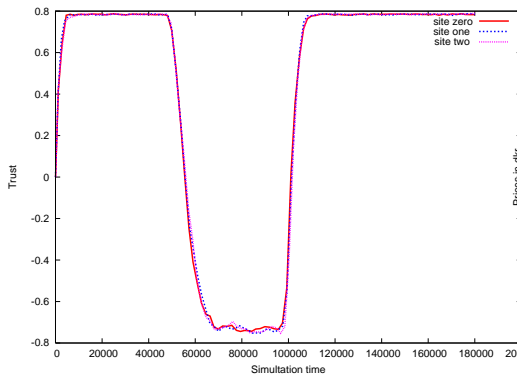


Figure 7.3: The trust for user zero on 3 different sites.

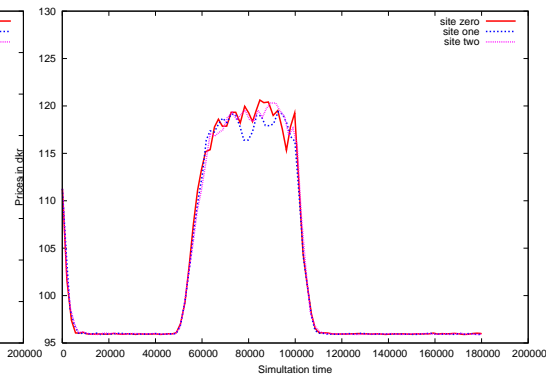


Figure 7.4: The price for user zero on 3 different sites.

3 different sites and 3 different prices. Again it shows the price fluctuates more when it is high than low.

## 7.5.2 Results from Experiment 2

- Aim: To show price change due to trust with varying power prices.

It is the same setup as in experiment one, except now the power price fluctuates. There is no threshold set for the trust score and it is assumed that strangers have the trust score of 0, which is neutral. In the beginning User zero is ill behaving but after a while he behaves good again as seen on figure 7.5. User Two behaves good all the time. The price for User Two falls quickly and has some small peaks in the beginning as seen on figure 7.6 on the next page. This is due to the changes in power price. Another noticeable feature that power is changing is in the end of the experiment the prices keep on lowering.

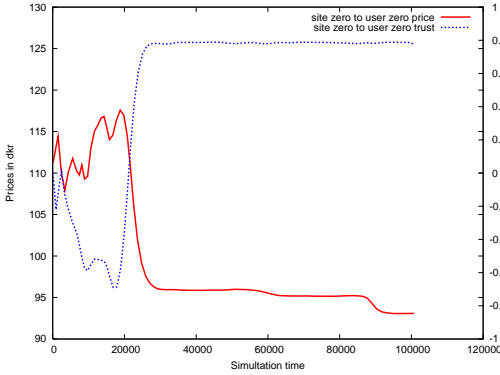


Figure 7.5: The price vs trust at different power rices.

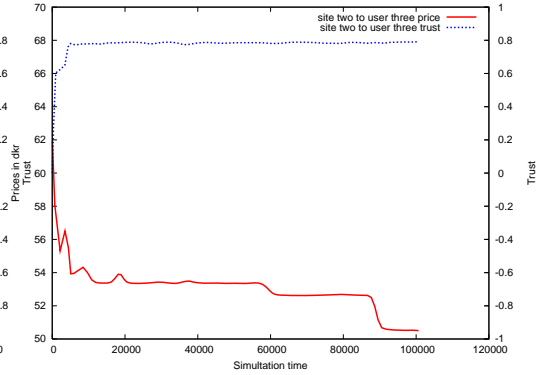


Figure 7.6: The price vs trust at different power prices

### 7.5.3 Results from Experiment 3

- Aim: To show different prices due to different parameters on the sites

The test setup is the same as in experiment three except that now all the parameters is different. Now Site zero, one and two all have different hardware costs. As expected the prices is very different with the lowest price at the site

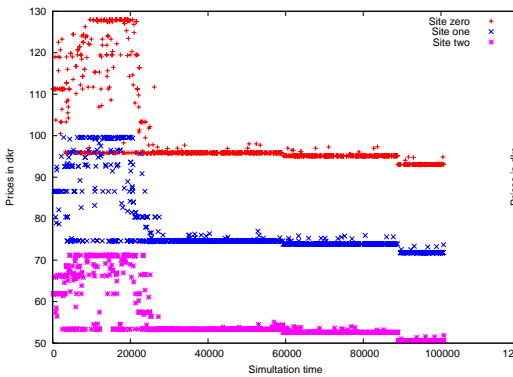


Figure 7.7: The price for different hardware prices.

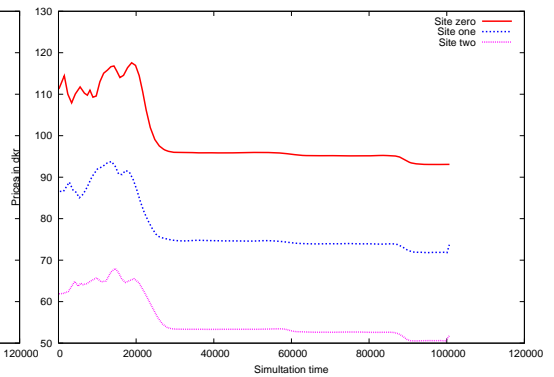


Figure 7.8: The price approximated for different hardware prices.

with the lowest hardware cost. Since the hardware cost is the largest item in the equation 5.1 on page 35 and used to add the percentages in 5.2 on page 35 it makes good sense.

### 7.5.4 Results from Experiment 4

- Aim: To show how threshold affects site capacity and users price

In this setup the threshold was set for both users and sites. Each site is assigned 3000 CPU's and there is 15 users and 20 sites as can be seen C.3 on page 90. If the sites does not trust the user, which is the users trust score is lower than the threshold, the site will return a higher price, since the risk is set to 80 percent. It is the strategy of we will do business with you but you will pay overprice for it, since we lack trust in you. If the site ill behave, the user will now not select that site. It will simply as another site. The test starts out with site one behaving good and giving discounts. Site one it the most popular site, until the site ill behaves. Then the users will not assign jobs to the site. As seen on figure 7.9 the users stop Too send job when the site is not behaving very good. On figure 7.10 the users trust To the site is shown.

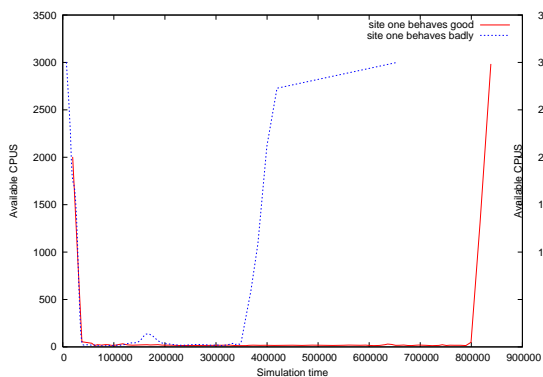


Figure 7.9: Site one behaves good at all time and Site One behaves badly.

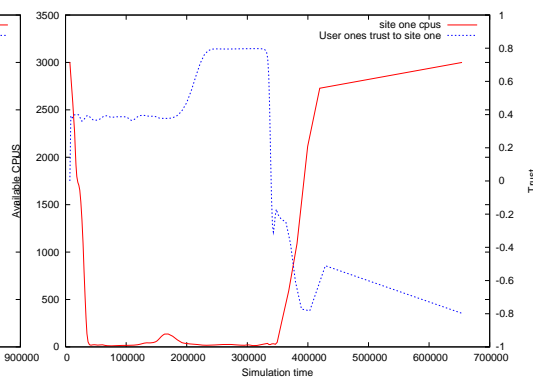


Figure 7.10: When the trust declines To the site, the site does not receive new jobs. This leads To an increase in available CPU's.

In the case where the user ill behaves and he violates the site's threshold, it can be seen that the price rises to 300 Dkr. from approximately 100 dkr. The user can still submit job as long there is money on his account in the bank, but as soon as this is empty no site will perform a job for him. The price behavior is seen on figure 7.11 on the next page.

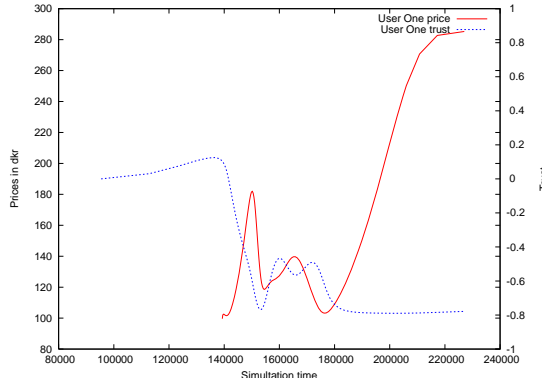


Figure 7.11: User One is not liked and the price rises.

## 7.6 Results from Test Case 2

The primary objective of test case two is to implement the revenue management on the system. Now all the sites have CPU's to account for and they need to track how many CPU's are used at all time. There is also a bank now who has to keep track of the money flows from user to sites accounts, keep track of credit rejections. The bank also have an extra function of counting how many sites who had to turn down a job, if the sites is overbooked. This happens only if all sites is fully booked at the same time. Furthermore the bank keep track of which sites has the most jobs and which site had the most revenue.

### 7.6.1 Results from Experiment 5

- Aim: To show the users selects cheapest sites.

In this setup there is 10 user and 10 sites. The sites now have a capacity of 500 CPU's each and site zero have a lower hardware cost than the other sites. As seen on figure 7.13 on the following page all sites is filling up quickly since they all within the 20000 first units of time is fully occupied. Notice the slope in the end of the simulation. Notice how site zero is more reluctant To let go, which indicates that the last jobs is given To site zero. The price difference is very clear in figure 7.13 on the next page, where the prices offered To user zero from site zero and site one. In figure 7.14 on the following page it is shown the site zero receives most jobs at any test trial.

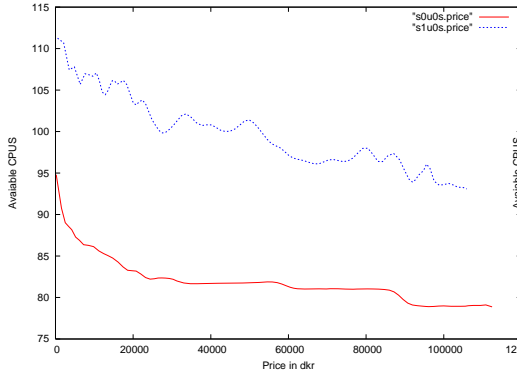


Figure 7.12: The price for user zero is cheapest at site zero.

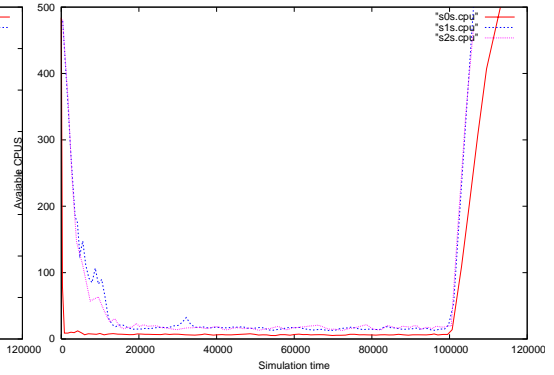


Figure 7.13: All the sites is occupied fast. How ever site zero is occupied immediately.

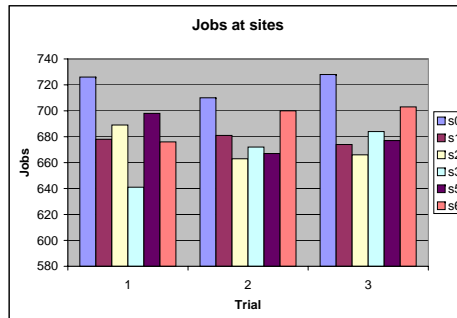


Figure 7.14: Out of three re-runs site zero has the most jobs.

### 7.6.2 Results from Experiment 6 and 7

- Aim for experiment six: To show that sites calculates protection levels.
- Aim for experiment seven: To show that sites checks that discount in not below the cost price.

The setup contained 4 out of 100 sites which has discount rates. The site was site zero, one, two and three. The special thing about three discount rate was it was set very high. In fact set so high that the site tries to sell with loss, since the price is below production cost. The discount rates ranged from 10% to 25%. When the simulation starts it calculates the protection level for the site if dis-



count rate is set. As seen in listing in 7.1 for site zero the protection level is set To 228, for site one 203, site two 193 and for site three it is set To 170.

Listing 7.1: Output from a simulation with discountsrate

```
1 BANK OPRETTET
2 site id: s0
3 my discount is: 10
4 my CPU's are is: 600
5 my name is s0 and my protectionlevel is: 228
6 my name is s0 and my bookinglimit is: 372
7 site id: s1
8 my discount is: 20
9 my CPU's are is: 600
10 my name is s1 and my protectionlevel is: 203
11 my name is s1 and my bookinglimit is: 397
12 site id: s2
13 my discount is: 25
14 my CPU's are is: 600
15 my name is s2 and my protectionlevel is: 193
16 my name is s2 and my bookinglimit is: 407
17 site id: s3
18 my discount is: 40
19 my CPU's are is: 600
20 my name is s3 and my protectionlevel is: 170
21 my name is s3 and my bookinglimit is: 430
22 site id: s4
23 my discount is: 0
24 my CPU's are is: 600
25 site id: s5
26 my discount is: 0
27 my CPU's are is: 600
28 site id: s6
29 my discount is: 0
30 my CPU's are is: 600
```

It allows the discount sites, which all have the capacity of 600 CPU's, to sell 600 – *protectionlevel* cheap. The rest has to be sold at normal price. Since the users always will select the cheapest site, all the discount sites are fast booked, which should leave many jobs to the 96 other sites. However the result show that the discount sites receives a total of 20% of all jobs, which is 1,4 million jobs in this setup. In other words the 96 other sites had to share 80% of all jobs, which on average leaves 0,83% of jobs To each one. The distribution of earnings and jobs is shown in figure 7.16 on the following page and 7.15 on the next page.

Notice how site three does not show any where. It makes sense since the price is 10 time higher than the normal price, since the discount rate is to high. No user have selected to send a job since it is to expensive. Another finding is that

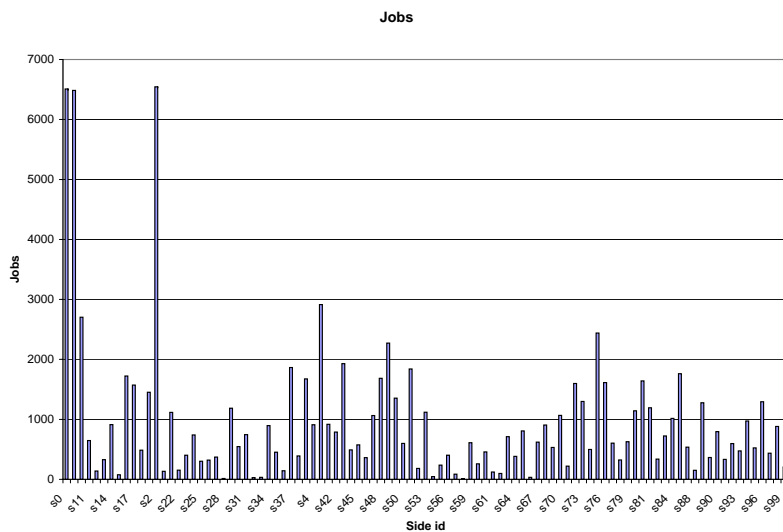


Figure 7.15: The job distribution for 1,4 million jobs.

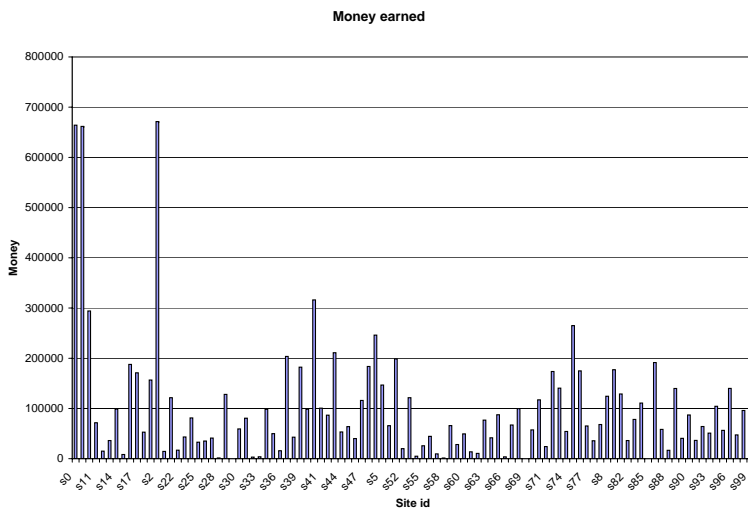


Figure 7.16: The earnings distribution for 1,4 million jobs.

the sites with most earnings also are the sites with the lowest average price for jobs as shown in figure 7.17 on the facing page.

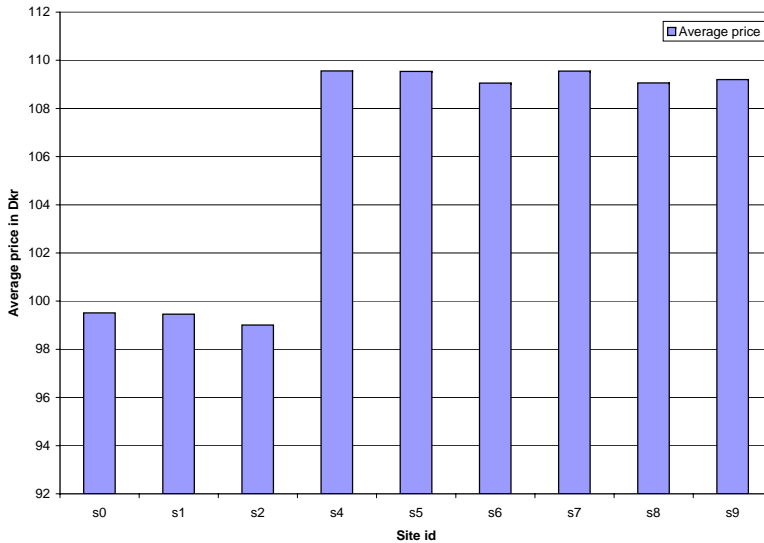


Figure 7.17: The average price for 10 first sites. Site 3 is not in since it had no jobs because of to high price.

### 7.6.3 Results from Experiment 8

- Aim: To show site capacity is limited.

The setup is similar To the setup in 7.6.2 on page 70 also with discount sites. To test it throughly I made three subtest with the following setup:

**Test 1** 20 users and 10 sites. Sites have 600 CPU's.

**Test 2** 20 users and 10 sites. Sites have 1200 CPU's.

**Test 3** 20 users and 20 sites. Sites have 600 CPU's.

The first test is to see what happens when the environment has limited capacity compared To the jobs it has To serve. In the second test I have the same amount of sites as in test one, but now the capacity has doubled. In the last test I have increased the number of sites and kept the CPU's at the same level as in test one.

The result of test one is shown in 7.18 on the following page and it shows an almost even distribution of jobs expect from site one and six. Its hard to tell

which sites is the discount sites and who is not.

The result of test two is shown in 7.19 which show a little increase in jobs per site in overall and site one and six still does not receives many jobs.

The result of test three is shown in 7.20 on the facing page which show an even distribution of jobs and each site has less jobs than in the other test.

Another way of looking on the three test is to see how many times there were overbookings. The overbooking event indicates that *all sites* is fully booked. This will tell more about how the system handles capacity issues. The result is shown in figure 7.21 on the next page and it shows that a system with many sites and few CPU's handles better than any other system. The test 4 is result from experiment nine which is located 7.6.4 on the facing page.

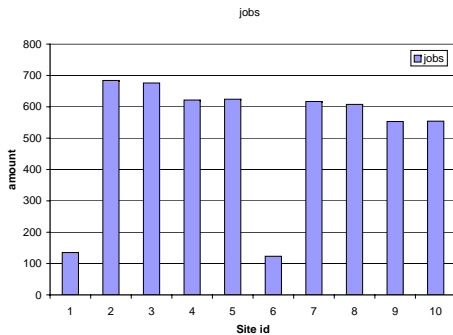


Figure 7.18: Test 1: Amount of jobs with few sites with few CPU's.

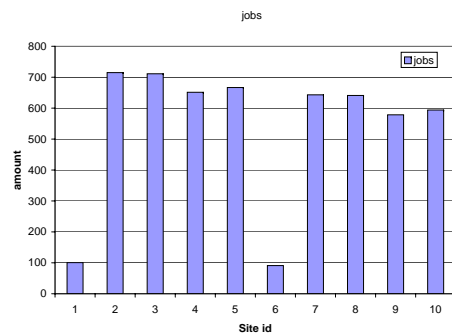


Figure 7.19: Test 2: Amount of jobs with few sites with many CPU's.

It can be noticed that the average jobs conducted per site is: The difference

Test	jobs
1	519
2	539
3	424
4	383

Table 7.1: The average jobs for the 4 test per site.

between test two and three is 120 jobs on average per site. In test two the sites had To handle 120 more jobs within the same time frame as in test three. Since the amount of CPU's follow a normal distribution with a mean of 15 CPU's, this lead To an increase in CPU's of 1800 per site. In test two the sites have

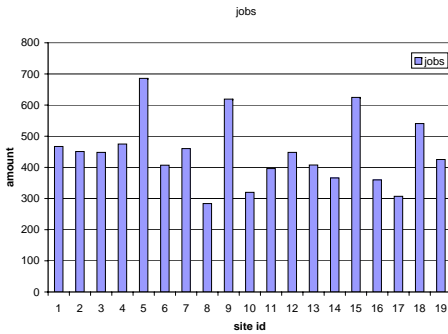


Figure 7.20: Test 3: Amount of jobs with many sites with few CPU's.

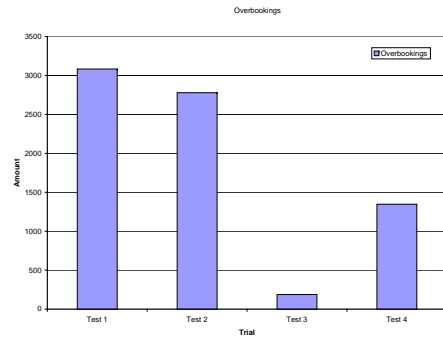


Figure 7.21: Amount of overbookings in the four subtest conducted.

600 CPU's more than the sites in test three, but the increase in work load is tripled.

### 7.6.4 Results from Experiment 9

- Aim: To show Users can only select few sites.

The test setup consisted of:

- 20 users and 20 sites. Sites have 600 CPU's. Different user pattern.

The different user pattern is instead of asking all the sites for a price, the user is only allowed to ask a random site which holds a position in the sites array and from that position and to the end of the sites array. The change in code was done in Resourcesite.cpp and consisted of the change in 7.2.

Listing 7.2: Now the user have To select from the random id to end of site array.

```
1 for (rIter=rSites.find(tempsite); rIter!=rSites.end(); rIter++)
```

The result of the test is shown in figure 7.22 on the following page and it resembles the pattern of test one and two in the previous section.

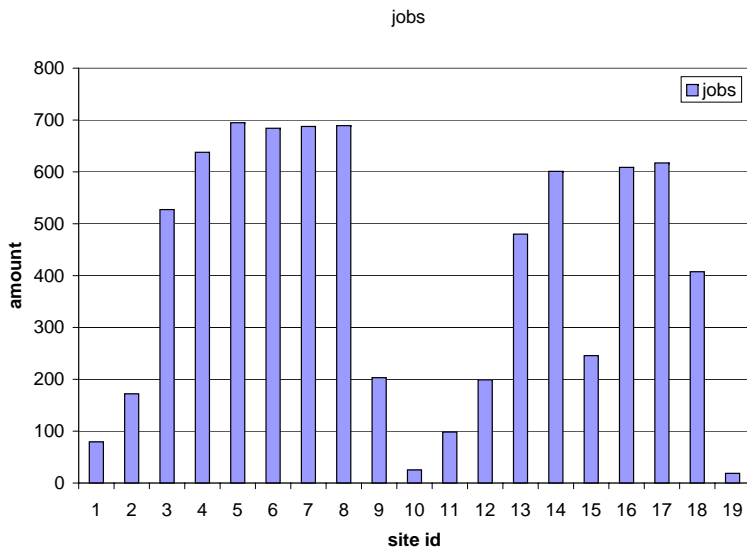


Figure 7.22: The distribution of jobs when users selection is limited.

## 7.7 Summary

In this chapter I start out with defining my test cases. These test cases are meant to show how prices fluctuates with the users trust and how the sites handles the capacity in a way that ensure an optimum. In the results I show how the different components interact according to specification.

# Discussion

---

## 8.1 Overview

In the first section of this chapter I summarize my results and discuss them. In the second section I will discuss possible future works for trust based accounting in grid systems.

## 8.2 Archived Results

I have archived to find a price for jobs in a Grid system. The price is dependent on many varying parameters which changes all the time due evolution of hardware. One of the parameters in the price was trust value derived from the system, which was mapped to an economic risk. This risk was set in the range of 0 to 20 percent, but it was not investigated what would happen if another mapping was used, say from range 0 to 40 percent. This mapping also assumed the trust score from the Grid system to be between -1 and 1. Another metric could have been used, but this was not analyzed.

The price was found with in relation to CPU requests, memory and storage. In

the price the CPU, memory and storage was assumed to have an equal weight of  $1/3$  for each. The cross relation was not simulated due to lack of relevant data and therefore it was difficult to give a qualified estimate of the use of memory and storage. The use of CPU time requests was found for the Nordugrid site, but there was no data of the amount of CPU's requested. I used a fraction of the CPU time request, since it followed normal distributed, to simulate the request of CPU's. I am confident that the request for CPU's follow the normal Distribution in a Grid system, but I do not have data to support my claim. Another issue with the relation between CPU's, memory and storage is that today storage is cheap compared to the price for CPU's and memory. The price for 2 CPU's compared to 4 GB of memory in a industrial server is almost the same and therefore it has to taken into account when calculating a just price. In my simulation this was not investigated, since I only looked at CPU's and CPU time.

In my assumptions I assume that the user always will buy. In the real world some users might want to wait to see if the prices decreases over time. The user is allowed to ask all sites in simulation, but in a Grid system with many sites this could turn out to be not such a good idea since it will take much time. Imaging the user talking to a American site which seems to be the cheapest, he will buy at that place since for the user it looks to be the best deal. However if there is network latency it not the best deal. The network latency have to be taken into consideration when pricing for the price to be more realistic.

The price was calculated according to the trust between the site and the user. If the site and user never had talked to each other before I used a default value. My strategy was that everybody is not bad and not good, but neutral. The after some transactions the site and user would have developed a trust between them, which indicates if you are behaving good or bad. Another strategy could have been that every user is guilty until otherwise proven. This could have been implemented instead in order for the site to have the highest possible prices at all time.

The usage of capacity management for resource sites proved very powerful. The sites that used discount rates where fully booked most of the time and they had the highest revenues. If the other sites, which did not have discount rates, were allowed to dynamically to change the prices to get more jobs, it would start a price war. It is not necessarily a bad thing with price war just as long as nobody can sell below their production cost, which is also called price dumping.

It was surprising to see that it was not the amount of CPU's at each site com-



pared to the amount of sites which had an impact on how many overbookings that occurred. When I looked at the average amount of jobs per site it was clear that an increase in CPU's only gave an a small deduction average job per site whereas in the case with more site there was an significant deduction. This lead to an triple increase of workload for each site, but they only had the double capacity. Also the time the CPU's are occupied plays a role and since there are more jobs, the CPU's stays occupied longer. If the arrival rate of CPU's request arrived with another Poisson distribution it could have resolved the problem, since the request arrive with a longer time interval.

The actual price I find is is realistic for one node in a Grid site. The reason for this is the model is build up with the cost for only one server, but can easily be expanded to include multiple nodes. The parameters that changes are the cost price for the hardware and amount of kWh used.

The price is only a price the variable cost and the fixed cost is not accounted for. The fixed cost of building the server room and the network infrastructure is not included. In order for this to be used in an enterprise this has be taken into account.

## 8.3 future work

My thesis investigates how to find a price and how to manage the resource sites. What is missing is how to handle jobs that run for many days. The Black-Scholes theorem for bond is introduced in the background sections on 4.5 on page 29 and seems to be a very interesting mechanism to calculate the price for a job that runs for more days. The idea is that the user and the site in a Grid system negotiates the time limits for the jobs. The user might want to run it within 5 days but is not sure when. The site needs to have a forecast function for the production cost over the next 5 days in order to calculate price for the bond and the price for performing the job. This way the total price to the user consist now of the price for the bond and the strike price, which is the price the user pays when the job is submitted.

Another way of using the Black-Scholes theorem is in business, where one firm can have call-options to users and by put-options of other sites. This way the site can have a fixed amount of capacity and when over capacity is needed they use put-options at other sites.

My simulation is build upon the notation of a bank. In the future there is a need for an introduction of a monetary currency on the Grid. There has be economic entities such as banks on the Grid who have high trust both from the users and from the sites. In order for the Grid to expand beyond the academic communities and into the commercial world this is a must.

# Conclusion

---

In my thesis I have investigated the different economic schemes for a Grid system and I have investigated how reputation trust can be utilized in calculating a price for a resource. Furthermore the field of revenue management with focus on capacity optimization have been investigated to control the resources at a Grid site.

I found that trust is a good way of controlling the price of resource to a user, since it awards a user with high trust with a low price compared to the price for a user with low trust.

Another discovery was that discount prices is a very powerful way of getting users to buy your resources since the resources utilizing this was fully booked. Discount pricing also lead to the highest earnings even though the average price was low compared to the sites, who did not have discount prices.

To find a just price for a Grid system it has to be just for both the user and the resource site and I believe that my model gives such a price.



APPENDIX A

# Simulation flow chart

---

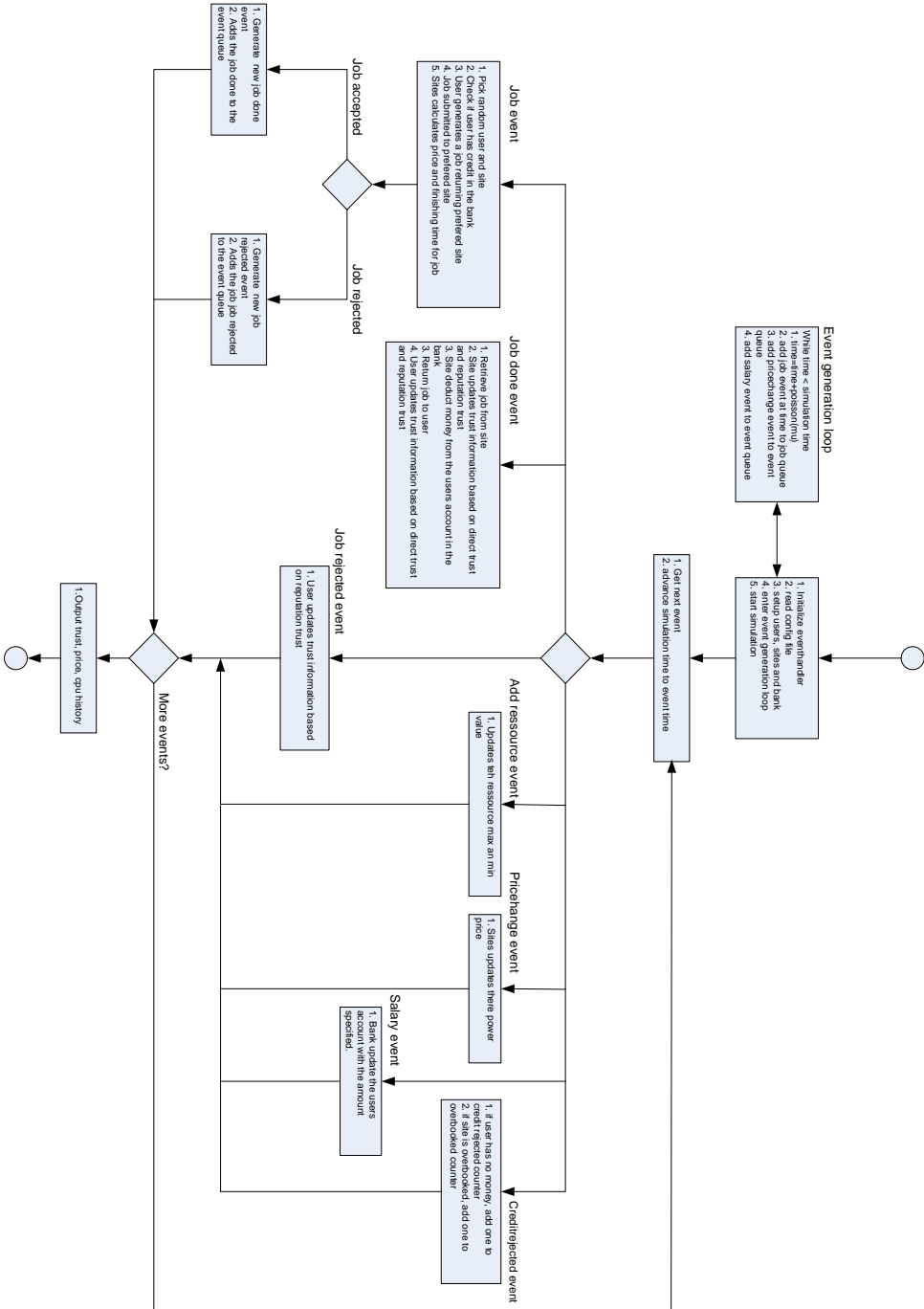


Figure A.1: The simulator flow chart

## APPENDIX B

# Contents on CD

---

The enclosed cd contains the following directories and files:

**sourcecode** Contains the source code for the simulation

**experiments** Contains the configuration files for the different experiments and their figures.

**scripts** Different perl and bash scripts to handle the data output. There is also a gnuplot script.

**report** The master thesis as pdf.

**Nordugrid log database** Contains the logfiles from nordugrid from January 2005 to November same year. It is contained both as a Mysql database and a Microsoft access database. It also contains the different queries and output from this. It also contains figures from these queries.





# Configuration files

---

In this chapter the different configuration files is specified. For experiment one all the items is specified, but for the rest of the experiments only the relevant parameters is specified.

## C.1 Experiment 1

```
1 ; Configuration file for grid-simulator
2 ; a semicolon indicates a comment
3
4 [Config]
5 users=11
6 sites=10
7 time=180000
8 mu=10.73
9 normalmu=156.75
10 normalsigma=56.55
11
12 path=/s973487/masterthesis/simulation/simul/exp4/
13
14 ; Various optional values for the [Config] section.
15 ; If not specified the the default is used.
16 ;
17 ; random number seed (might get overwritten at the commandline)
18 seed=100
```

```
19 ;
20 ; print a summary of the read config. This is very usefull when
21 ; debugging "strange" simulation output (default is no)
22 ;printconfigsummary=yes
23 ; prefix for the output dir (created in path). default is
24 ; first letter of trust_eval + first letter of rep_eval
25 ; for example with fuzzy-beta the prefix is fb
26 dirprefix=fbpricevstrust
27
28 ; write the the full output path to std.err
29 ; this can be useful when making scripts (default is no)
30 stderrpath=yes
31
32 ; reuse the jobevent objects (default is yes)
33 ;reuseevents=no
34
35 ; reuse the job objects (default is yes)
36 ;reusejobs=no
37
38 ; use references when calculating reputation (default is yes)
39 ;reputationrefs=yes
40
41 ; number of references to provide and use (default is 4)
42 ;numberofrefs=5
43
44
45 [Default]
46 i=1
47 j=1
48 groups = 0
49 threshold=-0.1
50 trust_eval=fuzzy
51 rep_eval=beta
52 direct_weight=0.6
53 smoothing=1.0
54 sigma=3
55 stranger=0.0
56 ;discrete_ratio=0.5
57 forgetting=0.4
58 base_weight=6
59 ;rep_ratio=0.0
60
61
62
63 [Economy]
64 markup=10
65 depreciation=6
66 discount=0
67 protectionlevel=0.0
68 hardwarecost=45000
69
70
71 ;(optional) section for specifying which resources save combined
    trust info
72 [CtSave]
```

```

73 default=no
74 s0=u0 u1
75 s1=u0 u1 u2 u3
76 s2=u0 u1 u3
77
78 [Resources]
79 <u0>
80 time=50000 i=20 j=20
81 time=110000 i=1 j=1

```

## C.2 Experiment 2+3

Since it the power that changes no need for the configuration file to change.

```

1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=20
5 sites=10
6 time=180000
7 mu=10.73
8 normalmu=156.75
9 normalsigma=56.55
10
11
12 dirprefix=fbpricevstrust
13
14
15 [Default]
16 i=1
17 j=1
18
19 threshold=-0.1 ; default threshold
20 trust_eval=fuzzy
21 rep_eval=beta
22 direct_weight=0.6
23 smoothing=1.0
24 sigma=3
25 stranger=0.0
26 ; discrete_ratio=0.5
27 forgetting=0.4
28 base_weight=6
29
30
31 [Economy]
32 markup=10
33 depreciation=6
34 discount=0
35 protectionlevel=0.0
36 hardwarecost=45000
37
38 section for specifying which resources save combined trust info

```

```
39 [CtSave]
40 default=no
41 s0=u0 u1
42 s1=u0 u1 u2 u3
43 s2=u0 u1 u3
44 u0=s0 s1 s2
45
46
47 [Resources]
48 <u0>
49 time=50000 i=20 j=20
50 time=110000 i=1 j=1
51 <s1>
52 time=80000 i=20 j=20
```

## C.3 Experiment 4

### C.3.1 Site ill behaves

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=15
5 sites=20
6 time=800714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 ;
12 ; random number seed (might get overwritten at the commandline)
13 seed=100
14
15
16 [Default]
17
18 threshold=-0.2 ; default threshold
19
20
21 [Economy]
22 cpus=3000 ; The capacity of the site in terms of cpus
23 <s1>
24 discount=20
25
26
27
28 [Bank]
29 initialmoney=1000000
30
31
```

```
32
33
34 section for specifying which resources save combined trust info
35 [CtSave]
36 default=yes
37
38 [Resources]
39 <s1>
40 time=300000 i=10 j=10
```

### C.3.2 User ill behaves

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=15
5 sites=20
6 time=800714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 ;
12 ; random number seed (might get overwritten at the commandline)
13 seed=100
14
15
16 [Default]
17
18 threshold=-0.2 ; default threshold
19
20
21 [Economy]
22 cpus=3000 ; The capacity of the site in terms of cpus
23 <s1>
24 discount=20
25
26
27
28 [Bank]
29 initialmoney=1000000
30
31
32 section for specifying which resources save combined trust info
33 [CtSave]
34 default=yes
35
36
37 [Resources]
38 <u1>
39 time=50000 i=10 j=10
```

## C.4 Experiment 5

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=10
5 sites=10
6 time=100714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11
12 dirprefix=cpuusage
13
14
15 [Economy]
16 markup=10
17 depreciation=6
18 discount=0
19 protectionlevel=0.0
20 hardwarecost=45000
21 cpus=500
22 <s0>
23 hardwarecost=35000
24
25
26 [Bank]
27 initialmoney=100000
28
29
30 ;(optional) section for specifying which resources save combined
   trust info
31 [CtSave]
32 default=no
33 s0=u0 u1 u2 u3
34 s1=u0 u1 u2 u3
35 s2=u0 u1 u2 u3
36 s3=u0 u1 u2 u3
```

## C.5 Experiment 6

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=100
6 time=1007140
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
```

```
11
12 dirprefix=manydiscounts
13
14 threshold=-1.0
15
16
17
18 [Economy]
19 markup=10
20 depreciation=6
21 discount=0
22 protectionlevel=0
23 hardwarecost=45000
24 cpus=600
25 <s0>
26 discount=10
27 <s1>
28 discount=20
29 <s2>
30 discount=25
31 <s3>
32 discount=40
33
34 [Bank]
35 initialmoney=300000
36
37
38 section for specifying which resources save combined trust info
39 [CtSave]
40 default=no
41 s2=u4
42 s3=u4
```

## C.6 Experiment 7

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=100
6 time=1007140
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11
12 dirprefix=manydiscounts
13
14 threshold=-1.0
15
16
17
18 [Economy]
```

```
19 markup=10
20 depreciation=6
21 discount=0
22 protectionlevel=0
23 hardwarecost=45000
24 cpus=600
25 <s0>
26 discount=10
27 <s1>
28 discount=20
29 <s2>
30 discount=25
31 <s3>
32 discount=40
33
34 [Bank]
35 initialmoney=300000
36
37
38 section for specifying which resources save combined trust info
39 [CtSave]
40 default=no
41 s2=u4
42 s3=u4
```

## C.7 Experiment 8

### C.7.1 Test 1

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=10
6 time=100714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 [Economy]
12 markup=10
13 depreciation=6
14 discount=0
15 protectionlevel=0
16 hardwarecost=45000
17
18 cpus=600
19
20 <s1>
21 discount=10
22 <s2>
```



```
23 discount=20
24 <s3>
25 discount=25
26 <s4>
27 discount=40
28
29
30
31 [Bank]
32 initialmoney=300000
33
34 default=yes
```

### C.7.2 Test 2

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=10
6 time=100714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 [Economy]
12 markup=10
13 depreciation=6
14 discount=0
15 protectionlevel=0
16 hardwarecost=45000
17
18 cpus=1200
19
20 <s1>
21 discount=10
22 <s2>
23 discount=20
24 <s3>
25 discount=25
26 <s4>
27 discount=40
28
29
30
31 [Bank]
32 initialmoney=300000
33
34 default=yes
```

### C.7.3 Test 3

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=20
6 time=100714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 [Economy]
12 markup=10
13 depreciation=6
14 discount=0
15 protectionlevel=0
16 hardwarecost=45000
17
18 cpus=600
19
20 <s1>
21 discount=10
22 <s2>
23 discount=20
24 <s3>
25 discount=25
26 <s4>
27 discount=40
28
29
30
31 [Bank]
32 initialmoney=300000
33
34 default=yes
```

## C.8 Experiment 9

```
1 ; Configuration file for grid-simulator
2
3 [Config]
4 users=50
5 sites=20
6 time=100714
7 mu=10.73
8 normalmu=155.75
9 normalsigma=56.56
10
11 [Economy]
12 markup=10
13 depreciation=6
14 discount=0
15 protectionlevel=0
16 hardwarecost=45000
```

```
17
18  cpus=600
19
20  <s1>
21  discount=10
22  <s2>
23  discount=20
24  <s3>
25  discount=25
26  <s4>
27  discount=40
28
29
30
31  [Bank]
32  initialmoney=300000
33
34  default=yes
```



## APPENDIX D

# Source codes

---

## D.1 Pricechange event

### D.1.1 pricechangeevent.h

```
1  #if !defined(PRICECHANGEEVENT.H)
2  #define PRICECHANGEEVENT.H
3
4
5  #include <iostream>
6  #include <string>
7  #include "event.h"
8  #include "resourceuser.h"
9  #include "resourcesite.h"
10
11
12  using namespace std;
13  /** The purpose of the Pricechange event is to provide an event
14      where each site receives a new power price. The event will
15      read the price in the "price.txt" file and distribute the
16          current day price to the sites.
17      */
17  class PriceChangeEvent : public Event {
18  private:
19      double price;
20  public:
21      /** Constructs the PriceChangeEvent
```

```

22     \param time is Time of the PriceChangeEvent
23     \param p is the power price
24     */
25     PriceChangeEvent(unsigned int time,double p) : Event(time),price(
        p){}
26
27     bool action(rmap& rUsers,rmap& rSites);
28
29
30     Event* get_event();
31 };
32 #endif

```

## D.1.2 pricechangeevent.cpp

```

1  #include "pricechangeevent.h"
2
3  bool PriceChangeEvent::action(rmap& rUsers,rmap& rSites){
4      rmap::iterator it;
5
6          ResourceSite* rs;
7
8              for (it=rSites.begin();it!=rSites.end();it++)
9                  {
10                     rs=dynamic_cast<ResourceSite*>(rSites[it->first]);
11                     //Update price in the site
12                     rs->update_price(price);
13
14                 }
15
16         return false;
17 }
18
19
20 Event* PriceChangeEvent::get_event(){
21     throw string("Trying_to_get_event_from_PriceChangeEvent");
22     return e;
23 }

```

## D.2 salary event

### D.2.1 salaryevent.h

```

1  #if !defined(SALARYEVENT.H)
2  #define SALARYEVENT.H
3
4
5  #include <iostream>
6  #include <string>
7  #include "event.h"

```

```

 8 #include "resourceuser.h"
 9 #include "resourcesite.h"
10 #include "bank.h"
11
12
13 using namespace std;
14 /** The purpose of the Salary event is to provide an event
15     where each site receives a new power price. The event will
16     read the price in the "salary.txt" file and distribute the
17     current day price to the sites.
18 */
19 class SalaryEvent : public Event {
20 private:
21     double salary;
22 public:
23     /** Constructs the PriceChangeEvent
24         \param time is Time of the SalaryEvent
25         \param p is the salary price
26     */
27     SalaryEvent(unsigned int time, double p) : Event(time), salary(p) {}
28
29     bool action(rmap& rUsers, rmap& rSites);
30
31     Event* get_event();
32 };
33 #endif

```

### D.2.2 salaryevent.cpp

```

 1 #include "salaryevent.h"
 2
 3 bool SalaryEvent::action(rmap& rUsers, rmap& rSites){
 4     rmap::iterator it;
 5
 6
 7     ResourceUser* rs;
 8
 9     for (it=rUsers.begin();it!=rUsers.end();it++)
10     {
11         rs=dynamic_cast<ResourceUser*>(rUsers[it->first]);
12         //Update price in the site
13         rs->salary(salary);
14     }
15
16     return false;
17 }
18
19
20 Event* SalaryEvent::get_event(){
21     throw string("Trying_to_get_event_from_SalaryEvent");
22     return e;
23 }

```

## D.3 creditrejected event

### D.3.1 creditrejectedevent.h

```
1  #if !defined(CREDITREJECTEDEVENT_H)
2  #define CREDITREJECTEDEVENT_H
3
4  #include <iostream>
5  #include <string>
6  #include "event.h"
7  #include "resourcesite.h"
8  #include "resourceuser.h"
9
10
11 using namespace std;
12
13 class CreditRejectedEvent: public Event{
14 private:
15     string uid;
16     Configuration* c;
17     bool condition;
18
19 public:
20     /** Constructs the CreditRejectedEvent.
21         \param u The id of the user (job creator)
22         \param t The time of the event
23     */
24     CreditRejectedEvent(string u, const int t, Configuration* conf, bool
25         overbooked) : Event(t), uid(u), c(conf), condition(overbooked)
26     {
27
28     /** The method will extract the job from the site and
29         return it to the job creator.
30         \param rUsers A reference to the map containing pointer to
31             the users
32         \param rSites A reference to the map containing pointer to
33             the sites */
34     bool action(rmap& rUsers, rmap& rSites);
35     Event* get_event();
36 };
37 #endif
```

### D.3.2 creditrejectedevent.cpp

```
1  #include "creditrejectedevent.h"
2
3  bool CreditRejectedEvent::action(rmap& rUsers, rmap& rSites){
4  //cout << "Creditcheck failed for : " << uid << endl;
5
```



```
6  if(condition)
7  c->countcredit(true);
8  else
9  c->countcredit(false);
10 return false;
11 }
12
13 Event* CreditRejectedEvent::get_event(){
14     throw string("Trying_to_get_event_from_JobdoneEvent");
15     return e;
16 }
```

## D.4 Bank

### D.4.1 Bank.h

```
1  #if !defined(BANK_H)
2  #define BANK_H
3
4  #include <iostream>
5  #include <fstream>
6  #include <map>
7  #include <list>
8  #include <vector>
9  #include <string>
10 #include <sstream>
11 #include <iterator>
12
13 using namespace std;
14
15 class Bank{
16
17 private:
18     typedef map <string,double> bankaccounts;
19     map <string,int> jobsawarded;
20     map <string,int>::iterator countjobs;
21
22     bankaccounts accounts;
23
24     double iniusermoney;
25
26     int countrejections,countoverbooked;
27
28 public:
29     Bank();
30
31     //Destructor
32     ~Bank(){}
33
34     void PutSalaryOnAccount(string id, double amount);
35 }
```

```

36 void Transfer_money_user_to_site(string sid,string uid,double
    amount);
37
38 void Set_Accounts(string id, double amount, bool initial);
39
40 void Set_init_money(double inimoney);
41
42 bool Has_Credit(string uid);
43
44 void count_creditrejected(bool condition);
45
46 void Print_accounts();
47
48
49
50
51
52
53 };
54
55 #endif

```

## D.4.2 Bank.cpp

```

1 #include "bank.h"
2
3 Bank::Bank() {}
4
5 void Bank::PutSalaryOnAccount(string id,double amount){
6     map <string,double>::iterator it;
7     //cout << "salary kaldt med amount " << amount << endl;
8     //search the map full of accounts
9     for(it=accounts.begin();it!=accounts.end();it++)
10    {
11        //Only use the user accounts
12        if((it->first)==id)
13            {
14                //Add the salary to the account
15                it->second+=amount;
16                //cout << "Salary put on : " << (it->first) << "
17                //cout << "Total amount on account " << (it->second) << endl;
18                }
19    }
20 }
21
22 void Bank::Transfer_money_user_to_site(string sid,string uid,double
    amount){
23     accounts[uid]-=amount; //Deduct the amount from the user
        account
24     accounts[sid]+=amount; //Add the amount to the site account
25
26     countjobs=jobsawarded.find(sid);
27

```

```

28     if(countjobs==jobsawarded.end())
29     {
30         jobsawarded[sid]=1;
31     }
32     else
33         jobsawarded[sid]+=1;
34
35 }
36
37 void Bank::Set_init_money(double inimoney){
38     iniusermoney=inimoney;
39 }
40
41 void Bank::Set_Accounts(string id, double amount, bool initial){
42
43     //cout << "Set_Accounts kaldt med: " << id << " og " <<
44         amount << " og "
45     //<< initial << endl;
46     if(initial)
47     {
48         //cout << "inde i initial " << endl;
49         countrejections=0;
50         countoverbooked=0;
51         //Does accounts exist
52         if(accounts.find(id)==accounts.end())
53         {
54             //cout << "accounten findes i
55                 mappet " << endl;
56             //is a user account
57             if(id[0]=='u')
58             {
59                 accounts[id]=iniusermoney;
60                 //cout<<"No account for "
61                     << id <<" existed"<<
62                     endl;
63                 //No accounts for id with money and
64                 the account is created with
65
66                 inimoney
67             }
68             //is a site account
69             if(id[0]=='s')
70             {
71                 //cout<<"No account for "
72                     << id <<" existed"<<
73                     endl;
74                 //No accounts for id with money and
75                 the account is created with
76
77                 zero, since it is a
78                 site
79                 accounts[id]=0;
80             }
81         }
82     }
83 }

```

```
70     else
71     accounts[id]=amount;
72
73     //cout << "for id : " << id << " står der " << accounts[id]
74         << endl;
75 }
76
77
78 bool Bank::Has_Credit(string uid)
79 {
80     if(accounts.find(uid)!=accounts.end() && uid[0]=='u'){
81         //there is an account
82         if(accounts[uid]>0)
83         {
84             //cout << "Credit check for user " << uid << "
85                 passed" << endl;
86             return true;
87         }
88         if(accounts[uid] < 0){
89             //cout << "The current account " << uid <<
90                 "has not enough money" << endl;
91             return false;
92         }
93     }
94     else
95     cout << "The current account " << uid << " does not exist"
96         << endl;
97     return false;
98 }
99 void Bank::Print_accounts()
100 {
101     map <string,double>::iterator it;
102     int maxjobs=0;
103     int sum=0;
104     string winner;
105
106     countjobs=jobsawarded.begin();
107     //Returns the site with the most jobs earned
108     for(countjobs=jobsawarded.begin();countjobs!=jobsawarded.end();
109         countjobs++)
110     {
111         cout << "Site:_" << countjobs->first << "_has_had_" << countjobs->
112             second << "_job" << endl;
113         sum=sum+jobsawarded[countjobs->first];
114         if(jobsawarded[countjobs->first]>maxjobs)
115         {
116             winner=countjobs->first;
117             maxjobs=jobsawarded[countjobs->first];
118         }
119     }
```

```

119 }
120
121 cout << endl;
122 cout << "The site:_" << winner << "_has won with most jobs with_"
    << maxjobs << "_jobs" << endl;
123 cout << endl;
124 cout << "The amount of credit rejections are:" << countrejections
    << endl;
125 cout << endl;
126 cout << "The amount of overbookings are:" << countoverbooked<<
    endl;
127 cout << endl;
128 cout << "The amount of jobs tracked in the bank:_" << sum <<endl;
129 cout << endl;
130
131 for(it=accounts.begin();it!=accounts.end();it++)
132     {
133         cout << "Konto:_" << (it->first) << "_har indestående_" <<
            (it->second) << endl;
134         double temp=(it->second);
135         string testid=(it->first);
136         if(testid[0]=='s')
137             {
138                 double revenue=temp/jobsawarded[testid];
139                 cout << "Det giver et samlet gennemsnit pris på_"
                    << revenue << endl;
140             }
141     }
142
143
144
145
146 }
147
148 void Bank::count_creditrejected(bool condition)
149 {
150     if(condition)
151         countoverbooked++;
152     else
153         countrejections++;
154 }

```

## D.5 resourceuser

### D.5.1 resourceuser.h

```

1 #if !defined(RESOURCEUSER_H)
2 #define RESOURCEUSER_H
3 #include <iostream>
4 #include <map>
5 #include <deque>

```

```

6  #include <string>
7
8  #include "job.h"
9  #include "entity.h"
10 #include "resource.h"
11 #include "resourcesite.h"
12
13 using namespace std;
14
15 /** This class contains the specific methods
16     needed by users. This class inherits the Resource class */
17 class ResourceUser : public Resource{
18     private:
19         /** Overwrites default assignment operator. Eliminates
20             the risk of accidentally assigning a ResourceUser
21             object to another ResourceUser object */
22         ResourceUser& operator = (ResourceUser&);
23
24         /** Overwrites default copy constructor. Eliminates the
25             risk of accidentally copying a ResourceUser object.
26             The ResourceUser objects must be unique. */
27         ResourceUser(ResourceUser&);
28
29         ///boolean the denotes if jobs should be reused.
30         bool reusejobs;
31     public:
32         /** Constructs a ResourceUser object. The constructor will also
33             specify
34             the resource is of type user and check if job objects should
35             be
36             re-used.
37             \param i The unique id of the user
38             \param c Pointer to the configuration object */
39         ResourceUser(string i, Configuration* c) : Resource(i, c)
40         {
41             myType=user;
42             reusejobs=true;
43             if(c->valid_key("[Config]", "reusejobs")){
44                 string reuse=c->get_parameter<string>("[Config]", "reusejobs
45                 ");
46                 if(reuse=="no" || reuse=="false"){
47                     reusejobs=false;
48                     //Only print once
49                     if(i=="u0")
50                         cout<<"Info: _Not_re-using_job_objects"<<endl;
51                 }
52             }
53         }
54     }
55
56     /// Destructor
57     ~ResourceUser(){
58     }
59
60     /** specific resource user methods */

```

```

58  /** Creates a job and returns it. If there is a rejected
59  job in the job queue then this is used. Otherwise a new job
60  is created. If the
61  proposed site is below the trust threshold then it is decided
62  if the site should be retried. This depends on what the
63  current simulation time is. The actual CPU time for the job
        will
64  be a random number between 20 and 40. The informed CPU time
        will
65  depend on the behavioural values.
66  \param sid The proposed site
67  \param rSites The map of sites
68  \param ctime The current simulation time
69  \return Pointer to the job to process */
70  Job* submit_job(string sid,rmap& rSites,const unsigned int ctime)
71  ;
72
73  /** Returns the job to the user. The job is examined and
74  a trust calculation is performed. First it is checked that
75  the site is equal to the site that processed the job. Then
76  the direct trust value is calculated. Afterwards the users
77  are asked to rate the site leading to the calculation of a
        reputation
78  trust value. These trust values are combined and the
79  trust value for the site is updated. All the trust values are
        saved
80  in sequences.
81  \param rUsers A reference to the map containing pointers to
        users
82  \param rSites A reference to the map containing pointers to
        sites
83  \param sid The site that processed the job
84  \param j Pointer to the processed job
85  */
86  void job_done(rmap& rUsers,rmap& rSites,const string& sid,Job* j)
87  ;
88
89  /** Returns a rejected job to the user. The job can later be
90  resubmitted to another site. First the other users are asked
91  to rate the rejecting site in order for a reputation trust
        value
92  to be calculated. If the user's combined trust value for the
93  site is higher than the reputation trust value, then the
        combined
94  trust value is set to the reputation value. Otherwise 0.1 is
95  subtracted from the combined trust value (unless it results
        in a value
96  below -1).
97  \param rUsers A reference to the map containing pointers to
        the users
98  \param rSites A reference to the map containing pointers to
        the sites
99  \param j Pointer to the rejected job

```

```

100     \param ctime The current simulation time
101     */
102     void receive_rejected(rmap& rUsers ,rmap& rSites ,Job* j ,const
        unsigned int ctime);
103
104     void salary(double amount);
105
106 };
107
108 #endif

```

## D.5.2 resourceuser.cpp

```

1  #include "resourceuser.h"
2
3  Job* ResourceUser::submit_job(string s ,rmap& rSites ,const unsigned
    int cTime){
4      double tempprice=0.0;
5      ResourceSite* rs;
6      rmap::iterator rIter;
7      mandatory* m=config->get_mandatory();
8
9      //ask 10 other
10     stringstream ss;
11
12         int cpuTime;
13         int cpus;
14         cpuTime=(int) rg.normal(m->normalmu , m->normalsigma);
15         cpus=(int) rg.normal((m->normalmu)/10,(m->normalsigma)/10);
16
17
18         //id of the purposed site
19         string sid= s;
20
21         //tempstring to locate cheapest site id
22         string tempsite=sid;
23
24
25
26         //Assume that current site sid is the cheapest
27         rs=dynamic_cast<ResourceSite*>(rSites[sid]);
28         tempprice=rs->calculatesiteprice(id ,cpus);
29
30
31         //Now we search from the site ID to the end to locate the
            cheapest site
32         for(rIter=rSites.begin();rIter!=rSites.end();rIter++)
33             {
34                 rs=dynamic_cast<ResourceSite*>(rSites[(
                    rIter->first)]);
35
36
37                 if(!rs->overbooked(cpus))
38                     {

```



```

39         //if the current site is
           cheapen then sid, we
           will use this site
40         if(rs->calculatesiteprice(
           id,cpus) < tempprice)
41         {
42             float tmptrust;
43             if(entities.find(
           rIter->first)==
           entities.end())
44             {
45                 tmptrust=(config->
           get_properties(
           id)->stranger;
46             }
47             else
48             {
49                 tmptrust=
           get_trust_value
           (rIter->first);
50             }
51
52             if(tmptrust>=threshold)
53             {
54                 tempsite=rIter->
           first;
55                 tempprice=rs->
           calculatesiteprice
           (id,cpus);
56             }
57         }
58     }
59 }
60
61
62     /*if(entities.find(tempsite)==entities.end() ) {
63         if(threshold>(config->get_properties(id))->stranger
64         )
65         throw string("No cheap trustworhty sites");
66     }
67     else if(threshold>get_trust_value(tempsite))
68         throw string(tempsite+" is not trustworthy and
           nobody are");
69
70     */
71     //Setting the id of the cheapest site found
72     sid=tempsite;
73
74
75
76
77
78     //test if site is known. If not an entity object will be created
79     emap::iterator eit;

```

```

80  eit=entities.find(sid);
81  if(eit==entities.end()){
82      properties* p=config->get_properties(id);
83
84      // Get entity options
85      saveOptions* save=config->get_save_options(id,sid);
86      float sv=p->stranger;
87      float fv=0;
88      //forgetting value is only used with the beta reputation system
89      if(repEval==beta)
90          fv=p->forgetting;
91
92      //Create entity object
93      Entity *e = new Entity(sid,sv,fv,cTime,save);
94
95      //Insert into map
96      entities[sid]=e;
97
98      //Set boolean if entity saves some trust sequences
99      if(entities[sid]->saves())
100          savesSomething=true;
101
102
103  }
104
105
106
107  /* Make a job and submit it .
108   * Check if there is preferences */
109
110
111      Job* j;
112
113
114
115  // Check if an old job has to be re-submitted
116  if(!jobs.empty()){
117      j=jobs[0];
118      jobs.pop_front();
119      // The actual cpuTime needed
120      if(j->rejected())
121          cpuTime=j->get_cpu_time();
122      //else
123  //{
124      //cpuTime=rg.irand(20,40);
125  //    cpuTime=(int) rg.normal(m->normalmu, m->normalsigma);
126  //    cout << "Cputime calculated by a normal demand function "
127  //        <<cpuTime << endl;
128  // }
129  j->set_cputime(cpuTime);
130  j->set_rejected(false);
131  j->set_prefered_site(tempsite);
132  }else{
133  // The actual cpuTime needed. Random number between 20 and 40
134      //cpuTime=(int) rg.normal(m->normalmu, m->normalsigma);

```

```

134     j = new Job(cpuTime, id);
135 }
136 /* This determines the informed cputime from a laplace
137 * distribution. This is determines the users behaviour.
138 * The width of the laplace distribution is determined by
139 * abs(min)+abs(max). The number la returned from the laplace
140 * distribution is then shifted. informed =la-min+max */
141 int la=(int)rg.laplace(abs(min)+abs(max));
142 int icputime=cpuTime+la-min+max;
143
144 /* The minimum informed cputime is 1
145 * Anything below that would not be creadable */
146 if(icputime<=0)
147     icputime=1;
148 j->set_informed_cpu_time(icputime);
149
150 /* Check if the site meets the minimum trust
151 * level. If this is not the case then
152 * select another site or check if the site
153 * should be retried. This should only be done
154 * if the job is not a high security job meaning
155 * that there already is a preferred site. */
156 if(get_trust_value(sid)<threshold && j->get_prefered_site()=="
    none"){
157     /* Finding out how many times the site's jobs has been
158     * rejected and when the last time the trust level
159     * was updated */
160     unsigned int rejected=entities[sid]->get_rejections();
161     unsigned int retryTime = entities[sid]->get_retry_time();
162     unsigned int retries =entities[sid]->get_retries();
163     /* If its the first time the site is rejected
164     * Set the time to retry the site to 2 * current time */
165     if(rejected==0){
166         entities[sid]->set_retry_time(2*cTime);
167         entities[sid]->set_rejections(rejected+1);
168     }
169     /* If the site already has been retried but
170     * did not improve set the next retrytime
171     * to (number of retries + 1) * retry time
172     * if its not ready again to retry */
173     if(retries>rejected && retryTime<cTime){
174         entities[sid]->set_rejections(rejected+1);
175         entities[sid]->set_retry_time((retries+1)*retryTime);
176     }
177     /* If its not time to retry the site
178     * pick a random site with a trust level
179     * above the threshold */
180     if(retryTime>cTime)
181         sid=random_entity(threshold, site);
182     else{
183         entities[sid]->set_retries(retries+1);
184     }
185
186     /* If no sites are above the threshold
187     * then pick the most trusted site */

```

```

188     if(sid=="none"){
189         sid=most_trusted(site);
190         j->set_prefered_site(sid);
191     }
192 }
193 // provide references
194 strset ref;
195 deque<string>::iterator it;
196 for(it=references.begin();it!=references.end();it++)
197     ref.insert(*it);
198
199 j->set_user_references(ref);
200 j->set_prefered_site(tempsite);
201
202 // Set the amount of requested cpus between 1 and 10
203 //cpus=cputime;
204 j->set_cpus(cpus);
205
206 //cout << "User: " << id << " has requested " << cpus << " cpus "
    <<endl;
207
208
209 // for debug
210 if(0){
211     cout<<id<<"_submitted_a_job_that_takes_"<<cpuTime<<endl;
212     cout<<"_told_that_it_would_take_"<<icputime<<endl;
213     cout<<"_Told_me_to_submit_job_to_"<<sid
214         <<"_who_i_trust:_"<<get_trust_value(s)<<endl;
215     if(get_trust_value(s)<threshold)
216         cout<<"_Below_my_threshold_for_sites!"<<endl;
217 }
218
219 //cout << "Job returned " << endl;
220 // Returning job
221 return j;
222 cout << endl;
223
224
225
226 }
227
228
229
230 void ResourceUser::salary(double amount)
231 {
232     config->putsalaryonaccount(id,amount);
233 }
234
235 void ResourceUser::job_done(rmap& rUsers,rmap& rSites,const string&
    s,Job* j){
236     if(s!=j->get_site()){
237         string error = id + "_wrong_site_specified_";
238         throw(error);
239     }
240     //Find current time

```

```

241     unsigned int cTime=j->time_finished();
242     // Calculate the direct trust value
243     float dt=direct_trust(j);
244
245     // Calculate the reputation
246     float rep = reputation(rUsers,rSites,s,user,j);
247
248     // Finally calculated the trust value
249     float trust = combine_trust(dt,rep);
250
251     //Write trust values
252     entities[s]->update_values(trust,dt,rep,cTime);
253     if(id=="u0" && s=="s0" && 0){
254         cout<<"time:_"<<j->time_finished()<<endl;
255         cout<<"ct:_"<<trust<<"_dt:_"<<dt<<"_rep:_"<<rep<<endl;
256     }
257
258     //Delete job
259     //delete j;
260
261     //reuse job
262     if(reusejobs)
263         jobs.push_back(j);
264     else
265         delete j;
266
267     //Save references
268     deque<string>::iterator it;
269     it=find(references.begin(),references.end(),s);
270     if(it!=references.end())
271         references.erase(it);
272
273     //if more the numberofrefs are provided the remove oldest
274     if(references.size()>=numberofrefs)
275         references.pop_front();
276
277     references.push_back(s);
278
279     /*Debug
280     if(0 && id=="u0"){ //&& s=="s0"){
281         char c;
282         cout<<"I am "<<id<<" and I got job back from: "<<s<<endl;
283         cout<<"My behaviour is :["<<min<<","<<max<<"]"<<endl;
284         cout<<"Job parameters: Cputime informed: "<<j->
                get_informed_cpu_time()
285                 <<" actual cpu time : "<<j->get_cpu_time()<<" time given "
286                 <<j->get_given_cpu_time()<<endl;
287         cout<<"Direct trust was : "<<dt<<endl;
288         cout<<"Reputation was: "<<rep<<endl;
289         cout<<"combined trust was: "<<trust<<endl;
290         cout<<"My references are :";
291         copy(references.begin(), references.end(), ostream_iterator<
                string>(cout, " "));
292         cout<<endl;
293         strset* sref=j->get_site_references();

```

```

294     if(sref->size()>=4){
295         cout<<"Received the references from site "<<s<<endl;
296         copy((*sref).begin(), (*sref).end(), ostream_iterator<string>
                >(cout, " "));
297         cout<<endl;
298     }
299     //cin>>c;
300     *
301     }*/
302 }
303
304 void ResourceUser::receive_rejected(rmap& rUsers,rmap& rSites,Job*
        j,const unsigned int cTime){
305     /* Update trust information
306     * When a job is rejected i cant calculate
307     * a direct trust value nor would I if i could.
308     * I will rely on the reputation of the site */
309     string sid =j->get_site();
310     //float trust = reputation(rUsers,rSites,sid,user);
311     float trust = reputation(rUsers,rSites,sid,user,j);
312     /* If I trust the site more then everybody else
313     * then set my trust to what the surroundings
314     * believe. If I trust the site more then everybody
315     * else the I just decrease my trust with 0.1 */
316     if(trust<get_trust_value(sid)){
317         entities[sid]->update_values(trust,cTime);
318     }
319     else{
320         trust=get_trust_value(sid)-0.1;
321         if(trust<-0.8)
322             trust=-0.8;
323         entities[sid]->update_values(trust,cTime);
324     }
325     /* Put job back into queue and re-submit
326     * at next opportunity */
327     if(reusejobs)
328         jobs.push_front(j);
329     else
330         delete j;
331     // Or Don't reuse. Delete job
332     //delete j;
333 }

```

## D.6 resourcesite

### D.6.1 resourcesite.h

```

1  #if !defined(RESOURCESITE_H)
2  #define RESOURCESITE_H
3  #include <iostream>
4  #include <map>

```

```
5 #include <deque>
6 #include <string>
7
8 #include "job.h"
9 #include "entity.h"
10 #include "resource.h"
11 #include "resourceuser.h"
12
13
14 using namespace std;
15
16 /** This class contains the specific methods
17     needed by sites. This class inherits the Resource class. */
18 class ResourceSite: public Resource{
19     private:
20
21         double siteprice, usersprice, powerprice, trustscore,
22             internalprice;
23
24         int markup, depreciation, discount, hardwarecost, jobtime,
25             protectionlevel, cpuasked;
26
27         int sitecpus, currentavailablecpus, occupiedcpus, bookinglimit;
28
29         economy* ep;
30         mandatory* m;
31
32         /** Overwrites the default assignment operator. Eliminates
33             the risk of accidentally assigning a ResourceSite
34             object to another ResourceSite object */
35         ResourceSite& operator = (ResourceSite&);
36
37         /** Overwrites the default copy constructor. Eliminates the
38             risk of accidentally copying a ResourceSite object.
39             The ResourceSite objects must be unique. */
40         ResourceSite(ResourceSite&);
41
42         /// The queue is used to store the rejected jobs before they are
43             returned to the user.
44         deque<job> rejectQueue;
45
46         //double calculatesiteprice(double trust);
47         double power(double number, int power);
48
49         //Calculates the protectionlevels
50         double calculateprotectionlevel(double disco);
51
52     public:
53         /** Constructs a ResourceSite object. The constructor will also
54             specify that the resource
55             is of type site.
56             \param i The unique id of the site
```

```

56     \param c Pointer to the configuration object */
57 ResourceSite(string i, Configuration* c): Resource(i, c)
58 {
59     myType=site;
60     ep=c->get_economy(i);
61     markup=ep->markup;
62     depreciation=ep->depreciation;
63     hardwarecost=ep->hardwarecost;
64     sitecpus=ep->sitecpus;
65     protectionlevel=ep->protectionlevel;
66     discount=ep->discount;
67     currentavaiblecpus=sitecpus;
68
69     cout << "site_id:_ " << id << endl;
70     cout << "my_discount_is:_ " << discount << endl;
71     cout << "my_CPUS_are_is:_ " << sitecpus << endl;
72     //cout << "my hardwarecost is: " << hardwarecost << endl;
73
74
75     usersprice=0;
76     powerprice=0;
77     if(discount>0)
78     protectionlevel=calculateprotectionlevel(discount);
79     cout << "my_name_is_" << id << "_and_my_protectionlevel_is:_ "
        << protectionlevel << endl;
80 ;
81
82 }
83 /** Destructor. Checks that rejected job queue is empty
84 */
85 ~ResourceSite()
86 {
87     if(!rejectQueue.empty())
88         throw string(id+"_had_a_job_that_was_not_returned_to_job_
            creator");
89 }
90 /* specific site functions */
91
92 /** Will (maybe) add a job to the site's job queue. This
93 depends on job creators trust value. If the trust value is
94 below the threshold then the job creator might get
95 retried depending on the current time. If it is rejected
96 then 0 is returned. Otherwise the finishing time for the
97 job is returned.
98 \param j The job to be processed
99 \param rUsers The map of users
100 \param ctime The current simulation time
101 \return The time the job will finish. */
102 int add_job(Job* j, rmap& rUsers, const unsigned int ctime);
103
104 /** Returns a processed job to the job creator. First the
105 user id is checked to make sure it matches the creator of the
106 job. Then the direct trust value for the job is calculated.
    Afterwards
107 the reputation of the job creator is calculated. These are

```



```

108         combined and the trust value for the user is updated. Finally
109         a
110         pointer to the finished job is returned.
111         \param rUsers A reference to the map containing pointers to
112         the users
113         \param rSites A reference to the map containing pointers to
114         the sites
115         \param uid The id of the user
116         \return Pointer to the processed job */
117 Job* job_done(rmap& rUsers, rmap& rSites, const string& uid);
118
119 /** Returns a rejected job. It is checked that the rejected
120 job in the queue actually is created by the specified user
121 before returning
122 a pointer to the job.
123 \param uid Id of the job creator
124 \return The rejected job */
125 Job* get_rejected(const string& uid);
126
127 void update_price(double price);
128
129 double calculatessiteprice(string uid, int cpusreq);
130
131 bool ResourceSite::overbooked(int cpurequested);
132
133 };
134 #endif

```

## D.6.2 resourcesite.cpp

```

1 #include "resourcesite.h"
2
3 int ResourceSite::add_job(Job* j, rmap& rUsers, const unsigned int
4 cTime) {
5     //Variables used by addJob
6     int cputime, timeGiven, timeDone;
7     double usertrust;
8     unsigned int time;
9     string uid=j->get_creator();
10    cpuasked=j->get_cpus();
11
12
13    //test if user is known. If not an entity object will be created
14    emap::iterator eit;
15    eit=entities.find(uid);
16    if(eit==entities.end()){
17        properties* p=config->get_properties(id);
18        //Get entity options
19        saveOptions* save=config->get_save_options(id, uid);
20        float sv=p->stranger;
21        float fv=0;

```

```

22     //forgetting value is only used with the beta reputation system
23     if(repEval==beta)
24         fv=p->forgetting;
25
26     //Create entity object
27     Entity *e = new Entity(uid,sv,fv,cTime,save);
28     //cout << "Entity obejkt lavet for : " << uid << " i tid: " <<
        cTime << endl;
29 //string c;
30 //    cin>>c;
31     //Insert into map
32     entities[uid]=e;
33     //entities[uid]=new Entity(uid,sv,fv,cTime,save);
34
35
36     //Set boolean if entity saves some trust sequences
37     if(entities[uid]->saves())
38         savesSomething=true;
39 }
40
41
42 /* If the user misbehaved earlier but bettered
43 * himself the record of the misbehaviour
44 * should not be kept. The term "once a crook
45 * always a crook" does not apply! */
46 if(get_trust_value(uid)>threshold &&entities[uid]->get_rejections
    (>0 ){
47     entities[uid]->set_retry_time(0);
48     entities[uid]->set_rejections(0);
49 }
50 /* Does the user have the required trust level
51 * or should it be tested again?
52 * If not send back timeDone=0. This means
53 * that the job was rejected! */
54 if(get_trust_value(uid)<threshold){
55     /* Finding out how many times the user's jobs has been
56     * rejected and when the last time the trust level
57     * was updated */
58     unsigned int rejected=entities[uid]->get_rejections();
59     //unsigned int lastUpdate = entities[uid]->updatedAt();
60     unsigned int retryTime = entities[uid]->get_retry_time();
61     unsigned int retries =entities[uid]->get_retries();
62     /* If its the first time the user is rejected
63     * Set the time to retry the user to 2* current time */
64     if(rejected==0){
65         entities[uid]->set_retry_time(2*cTime);
66         entities[uid]->set_rejections(rejected+1);
67     }
68     /* If the user already has been retried but
69     * did not improve set the next retrytime
70     * to (number of retries + 1) * retry time
71     * if its not ready again to retry */
72     if(retries>rejected && retryTime<cTime){
73         entities[uid]->set_rejections(rejected+1);
74         entities[uid]->set_retry_time((retries+1)*retryTime);

```

```

75     }
76     /* If its not time to retry the user
77     * simply drop the job. Meaning that 0 is returned */
78     if(retryTime>cTime){
79         /* Not that it is given 0 cputime
80         * and that I was the site that rejected the job */
81         j->set_given_cpu_time(0);
82         j->set_site(id);
83         // Storing the job temporarily
84         rejectQueue.push_back(j);
85         return 0;
86     }
87     else{
88         entities[uid]->set_retries(retries+1);
89     }
90 }
91
92
93
94 //Price is being calculated
95 calculatesiteprice(uid,cpuasked);
96
97
98
99     /* Finding out when job is going to start
100     * meaning finding out when the last job in
101     * the job queue finishes */
102     time=0;
103     if(job_count(>0)
104         time=jobs[jobs.size()-1]->time_finished();
105
106     // job queue is empty start job now
107     if(time<cTime)
108         time=cTime;
109
110     //Needed cpu time
111     cputime=j->get_informed_cpu_time();
112     /* This determines the given cputime from a laplace
113     * distribution. This is determines the sites behaviour.
114     * The width of the laplace distribution is determined by
115     * abs(min)+abs(max). The number la returned from the laplace
116     * distribution is then shifted. timegiven =la-min+max */
117     int la=(int)rg.laplace(abs(min)+abs(max));
118     timeGiven = cputime+la-min+max;
119     //Give job at least 1 cpu second
120     if(timeGiven<=0)
121         timeGiven=1;
122
123     //finish time
124     j->set_given_cpu_time(timeGiven);
125     timeDone =time+timeGiven;
126
127     jobtime=timeGiven;
128
129     //Adding site id to job

```

```

130     j->set_site(id);
131
132     //provide references
133     strset ref;
134     deque<string>::iterator it;
135     for(it=references.begin();it!=references.end();it++)
136         ref.insert(*it);
137
138     j->set_site_references(ref);
139
140     //Adding Job
141     j->start(time);
142     j->set_finish_time(timeDone);
143     jobs.push_back(j);
144
145
146     //for debug
147     if(0){
148         cout<<"---\nCurrent_time_is:_"<<cTime<<endl;
149         cout<<"There_are_"<<jobs.size()<<"_jobs_in_"<<id<<"_queue"<<
150             endl;cout << " called_in_jobsubmitted" << endl;
151     cout << "My_name_is_" << id << "_and_I_have_currently_" <<
152         currentavailablecpus << "_free_cpus_out_of_" << sitecpus << endl;
153
154     cout<<"job_will_start_at:_"<<time<<endl;
155     cout<<"job_needs_"<<cputime<<"_cputime"<<endl;
156     cout<<"will_give_it:_"<<timeGiven<<"_cputime"<<endl;
157     cout<<"Job_will_finish_at:_"<<timeDone<<endl;
158     cout<<"job_was_created_by:_"<<j->get_creator()<<endl;
159     cout<<id<<"_share_"<<config->compare_groups(id,uid)
160         <<"_with_"<<j->get_creator()<<endl;
161     cout<<"my_prefered_user_is:_"<<most_trusted(user)<<endl;
162     cout<<"Got_a_job_from:_"<<j->get_creator()
163         <<"_who_I_trust:_"<<get_trust_value(uid)<<endl;
164     if(get_trust_value(uid)<threshold)
165         cout<<"Below_my_threshold_for_users"<<endl;
166     }
167
168     if(!overbooked(cpuasked))
169     {
170         currentavailablecpus=currentavailablecpus-cpuasked; //removes
171             the cpus from the capacity
172         occupiedcpus=occupiedcpus+cpuasked; //adds the
173             cpus from the capacity
174         //if(currentavailablecpus<=sitecpus)
175         entities[uid]->update_cpus(currentavailablecpus,cTime);
176     }
177
178     //returning the time when the job is done
179     return timeDone;
180 }
181
182 }
183
184 }

```

```

180 Job* ResourceSite::job_done(rmap& rUsers, rmap& rSites, const string&
    u){
181     if(jobs.empty()){
182         string error=id+"Job_queue_was_empty_when_trying_to_get_"+u+"_
            job";
183         throw(error);
184     }
185
186     //get job creator
187     Job* j=jobs[0];
188     string uid=j->get_creator();
189     if(uid!=u){
190         string error=id+"_job_was_created_by_"+uid
191             +"_but_jobdoneevent_was_for_"+u;
192         throw(error);
193     }
194
195     unsigned int cTime=j->time_finished();
196
197     float dt=direct_trust(j);
198     // save dt
199     // entities [uid]->setDtUpdate(dt, cTime);
200     // float rep = reputation(rUsers, rSites, uid, site);
201     float rep = reputation(rUsers, rSites, uid, site, j);
202     // save rep
203     // entities [uid]->setRepUpdate(rep, cTime);
204
205     float trust = combine_trust(dt, rep);
206
207     //Write new trust value
208     // entities [uid]->setValue(trust, j->timeFinished());
209     //Write trust values
210     entities [uid]->update_values(trust, dt, rep, cTime);
211     entities [uid]->update_price(usersprice, cTime);
212
213
214     // Deduct amount bought from user...
215     config->deductmoney(id, uid, usersprice);
216     //cout << "actualprice " << testprice << " for id : " << dayratio
        << endl;
217
218     //remove finished job from queue
219     jobs.pop_front();
220
221     // Release all the cpus
222     //cout << "inde i jobdone" << endl;
223     //cout << "mit navn er site: " << id << " og jeg har " <<
        currentavaiblecpus << " cpus " << endl;
224     currentavaiblecpus=currentavaiblecpus+(j->get_cpus()); //
        removes the cpus from the capacity
225     occupiedcpus=occupiedcpus-(j->get_cpus()); //
        adds the cpus from the capacity
226     // if (currentavaiblecpus<=sitecpus)
227     entities [uid]->update_cpus(currentavaiblecpus, cTime);
228

```

```

229
230 //Update references
231 deque<string>::iterator it;
232 it=find(references.begin(),references.end(),uid);
233 if(it!=references.end())
234     references.erase(it);
235
236 if(references.size()>=numberofrefs)
237     references.pop_front();
238
239 references.push_back(uid);
240 //Debug
241 if(0 && id=="s0"){ // && uid=="u0"}{
242     char c;
243     cout<<"I am_"<<id<<"_and_I_finished_a_job_created_by:"<<uid<<
        endl;
244     cout<<"My_behaviour_is:["<<min<<","<<max<<"]"<<endl;
245     cout<<"Job_parameters:_Cputime_informed:"<<j->
        get_informed_cpu_time()
246         <<"_actual_cpu_time:"<<j->get_cpu_time()<<"_time_given_"
        <<j->get_given_cpu_time()<<endl;
247     cout<<"Direct_trust_was:"<<dt<<endl;
248     cout<<"Reputation_was:"<<rep<<endl;
249     cout<<"combined_trust_was:"<<trust<<endl;
250     strset* sref=j->get_user_references();
251     if(sref->size()>=4){
252         cout<<"Received_the_references_from_user_"<<uid<<endl;
253         copy((*sref).begin(), (*sref).end(), ostream_iterator<string>
            >(cout, "_"));
254         cout<<endl;
255     }
256 }
257 //cin>>c;
258 }
259
260
261
262 return j;
263
264 }
265
266
267
268 double ResourceSite::calculatesiteprice(string uid, int cpusreq){
269
270     double trust;
271     double risk, costs, totalcost, maxdiscount;
272
273
274     //cout << "Calculateprice kaldt af id :" << uid << endl;
275     emap::iterator pit;
276     pit=entities.find(uid);
277
278     //if there is no user entity use a default value
279     if(pit==entities.end()){
280         //cout << uid <<" er ikke funde i entity mappet" << endl;

```

```

281     properties* pi=config->get_properties(id);
282     trust=0; //pi->stranger;
283     //cout << "trust er lig med stranger i calculatestring ,
           hvor id er lig="<< id<< endl;
284     }
285     else
286     {
287     //cout << " inde i else stamentent i calculatesite hvor uid
           ="<< uid<< endl;
288     trust=get_trust_value(uid);
289
290     }
291
292     risk=(1-trust)*10;
293
294     //Test to see if the uses trustscore is above the threshold
295     if(trust<threshold)
296     {
297     risk=80;
298
299     }
300
301     costs=hardwarecost+(24*1.3*((double)powerprice/100)
           )*1000;
302     double tempdepre=(double) (100+depreciation)/100;
303     double temprisk=(double) (100+risk)/100;
304     double tempmarkup=(double) (100+markup)/100;
305
306     totalcost=costs*power(tempdepre,2)*power(temprisk
           ,2)*power(tempmarkup,2);
307
308
309
310     usersprice=totalcost/1000;
311     //cout << " lige foer return userprice " <<
           usersprice<< endl;
312     internalprice=hardwarecost/1000+(24*1.3*((double)
           powerprice/100));
313
314     //cout << "Usersprice er "<< usersprice << " og
           Internalprice er " << internalprice << " med
           trust " << risk << " og med power pris pa " <<
           powerprice <<endl;
315
316     //This is the maximum discount where there is a
           breakeven between cost and price
317     maxdiscount=100*((usersprice-internalprice)/usersprice);
318
319
320     if(discount>0 && discount<maxdiscount)
321     {
322
323
324     if((occupiedcpus+cpusreq)<bookinglimit)
325     {

```

```
326         usersprice=(1-((double)discount/100))*usersprice;
327
328     }
329 }
330 else
331 {
332     if(discount>maxdiscount)
333         usersprice=10*usersprice;
334
335 }
336
337     return usersprice;
338
339 }
340 }
341
342 void ResourceSite::update_price(double price)
343 {
344     powerprice=price;
345     //cout<< " powerprice i resource site kaldt af id " << id
346         << " med pris " << powerprice << endl;
347
348 }
349
350
351 Job* ResourceSite::get_rejected(const string& u){
352     /* Exit program if the queue of rejected jobs
353     * is empty. This should not happen. Throw exception */
354     if(rejectQueue.empty()){
355         string error= id + "_is_missing_a_rejected_job_created_by_" +
356             u + "_rejectQueue_was_empty";
357         throw(error);
358     }
359     jobqueue::iterator jit;
360     for (jit=rejectQueue.begin(); jit != rejectQueue.end(); jit++){
361         if((*jit)->get_creator()==u){
362             rejectQueue.erase(jit);
363             return (*jit);
364         }
365     }
366     /* Error. Job created by u was not found
367     * as a rejected job. Should not happen
368     * throw exception */
369     string error= id + "_is_missing_a_rejected_job_created_by_" + u;
370     throw(error);
371 }
372
373
374 double ResourceSite::power(double x, int n)
375 {
376     double product = 1.0;
377     int i;
378     if ( n == 0)
379         return 1.0;
```



```

380     else
381     {
382         for (i = 1; i <= n; i++)
383     {
384         product= product*x;
385     }
386         return product;
387     }
388 } // end of power
389
390
391 double ResourceSite::calculateprotectionlevel(double disco)
392 {
393 m=config->get_mandatory();
394 double normalmu=m->normalmu;
395 double z;
396 double normalsigma=m->normalsigma;
397 z=rg.inversecdf(disco/100);
398
399 //cout << " z er lig :." << z << endl;
400
401 return normalmu/5+z*(normalsigma)/5;
402 }
403
404 bool ResourceSite::overbooked(int cpurequested)
405 {
406
407 if((currentavaiblecpus-cpurequested)<0)
408 return true;
409 else
410 return false;
411 }

```

## D.7 configuration

### D.7.1 configuration.h

```

1 #if !defined(CONFIGURATION_H)
2 #define CONFIGURATION_H
3 extern "C" {
4     double parse_this(char *);
5     void setup();
6     int error_flag();
7     //double parseresult;
8     //int ERRORFLAG;
9 }
10
11 #include <iostream>
12 #include <fstream>
13 #include <queue>
14 #include <map>

```

```
15 #include <list>
16 #include <vector>
17 #include <set>
18 #include <string>
19 #include <sstream>
20 #include <iterator>
21 #include <ctime>
22 #include "saveoption.h"
23 #include "bank.h"
24
25 using namespace std;
26
27
28 typedef set<string> strset;
29 typedef list<int> intlist;
30 typedef set<int> intset;
31 typedef vector<string> strvec;
32 typedef map<string, strvec> submap;
33 typedef map<string, submap> cmap;
34 typedef map<string, string> strmap;
35
36 struct properties{
37     int i;
38     int j;
39     intset groups;
40     double threshold;
41     string trust_eval;
42     string rep_eval;
43     double direct_weight;
44     double sigma;
45     double smoothing;
46     double stranger;
47     //Only used if discrete combination
48     double discrete_ratio;
49     //Only used if beta
50     double forgetting;
51     double base_weight;
52     //Only used if discrete reputation
53     double rep_ratio;
54 };
55 //Map for resource options
56 typedef map<string, properties> romap;
57
58 struct mandatory{
59     unsigned int users;
60     unsigned int sites;
61     unsigned int time;
62     double mu;
63     double normalmu;
64     double normalsigma;
65     string path;
66     unsigned long int seed;
67     properties rDefaults;
68 };
69
```

```
70 //Parameters to calculate prices for the sites
71 struct economy{
72     unsigned int markup;
73     unsigned int depreciation;
74     unsigned int discount;
75     unsigned int hardwarecost;
76     double protectionlevel;
77     unsigned int sitecpus;
78 };
79
80 //map of parameters and sites.
81 typedef map<string,economy> ecmmap;
82
83
84 ///Structure containing id and time where a resource is to be added
85 struct addition{
86     string id;
87     unsigned int time;
88 };
89 //Structure used to order addition according to event time
90 struct addcomparison
91 {
92     bool operator () (addition * left, addition * right)
93     {
94         int l=(*left).time;
95         int r=(*right).time;
96         return l > r;
97     }
98 };
99
100 typedef priority_queue<addition*, vector< addition*,allocator<
      addition*> >,
101                    addcomparison > addqueue;
102 typedef list<addition*> addlist;
103
104 //Structure containing the id and time where the resource will
      change behavior
105 struct change{
106     string id;
107     unsigned int time;
108     int i;
109     int j;
110 };
111 //Structure used to order changes according to event time
112 struct changecomparison
113 {
114     bool operator () (change * left, change * right)
115     {
116         int l=(*left).time;
117         int r=(*right).time;
118         return l > r;
119     }
120 };
121
```

```
122 typedef priority_queue<change*, vector< change*,allocator<change*>
123     >,
124     changecomparison > changequeue;
125 typedef list<change*> changelist;
126
127 /** The purpose of this class
128     is to read the .conf file and set up the simulation.
129     The objects can query the configuration object for
130     the specifics of the simulation.*/
131 class Configuration{
132
133 private:
134     //map containing the config file as strings
135     cmap cfile;
136
137     economy ep;
138
139     //structure the contain the bare minimum of options
140     mandatory simOptions;
141     //map that contain properties for individual resources
142     romap resources;
143     /* flag that denotes that the simulation uses a
144     cmdline specified random number seed */
145
146     ecmap economyresources;
147
148     Bank nationaltreasure;
149
150
151     bool commandlineseed;
152
153     //Queue of additions
154     addqueue additions;
155
156     //list of additions (used for summary )
157     addlist listofadds;
158
159     //Time(s) where additions will happen
160     intset addtimes;
161
162     //Queue of changes
163     changequeue changes;
164
165     //list of changes (used for summary)
166     changelist listofchanges;
167
168     //Time(s) where changes will happen
169     intset changetimes;
170
171     ///Default saving options object
172     SaveOption* sOptions;
173
174     //the heuristics map (might not be used)
175     strmap heuristic;
```

```

176     //bool that denotes if custom heuristics are used
177     bool customHeuris;
178
179     /** Method that reads the entire configuration
180         file into a map. Lines beginning with ;
181         are skipped. When a line starts with [something]
182         is is a section. When a line starts with
183         <something> is is a subsection (in the section).
184         \param file The configuration file name
185     */
186     void read_whole_config(const string& file);
187
188     /* Boolean functions for checking various inputs */
189
190     ///Checks if line starts with '[' and has a ']' in it
191     bool valid_section(string& line);
192
193     ///Checks if line starts with '<' and has a '>' in it
194     bool valid_subsection(string& line);
195
196     ///Checks if line appears to be an id (must start with 'u' or 's'
197         and have a number
198     bool valid_id(const string& line);
199
200     /** Gets the next elemet (from the iterator position) from a
201         string vector. It checks that end of the vector has not
202         been reached. */
203     string& get_next_token(const strvec&,strvec::iterator& );
204
205     /** Tokenize a string. The string str is tokenized, using
206         the delimiters, and the result is inserted into the
207         vector.
208         \param str The string to be parsed
209         \param tokens The result vector
210         \param delimiters The delimiters to use */
211     void tokenize(const string& str,
212                 vector<string>& tokens,
213                 const string& delimiters = "_,=\t");
214
215     ///Removes beginning and ending chars from a string
216     string trim(string& s,const string& drop = "_");
217
218     ///Check if to properties structures contain equal values
219     bool compare_properties(const properties&,const properties&);
220
221     /** Read the mandatory options from the cmap. Also the default
222         properties are read. */
223     void set_mandatory_options();
224
225     /** Set the properties from a section and subsection. If
226         not the default properties cmap["[Default]"]["none"]
227         then only the values that are there are read
228         \param p The properties struct where the read values will be
229             saved
230         \param section The section to read from

```

```

229     \param subsection The subsection to read from */
230 void set_properties(properties& p,
231                  const string& section = "[Default]",
232                  const string& subsection = "none");
233
234 /*set the different economic parameters for sites and subsites*/
235 void set_economies();
236
237 // set the parameters for the economy
238 void set_economy(economy& test, const string& section="[Economy]",
239                 const string&
240                 subsection="none");
241
242
243 //set up the bank
244 void set_bank();
245
246
247
248 bool compare_economy(const economy& e1, const economy& e2);
249
250
251
252 /** Reads the section cmap[section] for subsections.
253     If a subsection is there then the set_properties method
254     is called to check for custom properties. It is assumed
255     that a subsection is a <resource_id>. */
256 void set_resources(const string& section = "[Resources]");
257
258 /** Reads the section cmap[section] for possible additions.
259     Meaning it looks for a format of:
260     time=some_unsigned_int some_resource_id
261     Example: time=15000 u50
262
263     The behavior can also be specified (optional)
264     Example: time=15000 u50 i=10 j=10
265
266     If no behavior is specified then the default values
267     are used. It does not matter if its written as i,j or j,i*/
268 void set_additions(const string& section = "[Add]");
269
270 /** Reads the section cmap[section] for subsections
271     These subsections are searched for possible changes
272     in the resource's behavior.
273     Meaning it looks for a format of:
274     time=some_unsigned_int i=some_int j=some_int
275     Example: time=15000 i=10 j=5
276     It does not matter if its written as i,j or j,i */
277 void set_changes(const string& section = "[Resources]");
278
279 /** The method scans the sections [CtSave],[DtSave] and [RtSave]
280     (in the cmap) for possible saving info. If a section is
281     left out the default action is save nothing */
282 void set_saveoptions();

```

```

283
284  /** Helper function for set_saveoptions.  */
285  void set_savetype(const string& section, const string& type);
286
287  /** Checks if there is a section [Heuristics] in the cmap
288      If its there then a heuristics map is created
289      with the values specified.  */
290  void set_heuristic();
291
292  /** Finds key in cmap[section][subsection] and returns
293      the tokenized result as a string vector using the
294      delimiters specified. This method is used by the
295      template functions.
296      \param section The section to search for
297      \param key The key to search for
298      \param subsection The subsection to look in
299      \param delimiters The delimiters to use when tokenizing  */
300  strvec find_and_tokenize(const string& section,
301                          const string& key,
302                          const string& subsection="none",
303                          const string& delimiters="_=#;\t");
304
305  void init();
306  public:
307  /** Constructs and automatically loads the configuration
308      file specified.
309      \param file The configuration filename.
310      */
311  Configuration(const string& file);
312  /** Constructs and automatically loads the configuration
313      file specified. Also the value s overwrites any randomnumber
314      seeds specified in the config file
315      \param file The configuration filename.
316      \param s The random number seed to use in the simulation
317      */
318  Configuration(const string& file, unsigned long int s);
319
320  /// Destructor
321  ~Configuration(){
322      //Check for unprocessed additions
323      if(!additions.empty()){
324          cout<<"There are additions that was not processed!"<<endl;
325          while(!additions.empty()){
326              addition* a= additions.top();
327              delete a;
328              additions.pop();
329          }
330      }
331      //Check for unprocessed changes
332      if(!changes.empty()){
333          cout<<"There are changes that was not processed!"<<endl;
334          while(!changes.empty()){
335              change* c= changes.top();
336              delete c;
337              changes.pop();
338          }
339      }

```

```

338     }
339     //Delete saveOptions
340     delete sOptions;
341
342 }
343
344
345 /** Template function that returns the value of a key in a
346     particular section. The type should explicitly defined
347     when using function. For example :
348     double value=configuration.get_parameter<double>("section",
349         "key");
350     The function is defined in the header file for portability
351     between different compilers.
352     \param section The section to search in (cmap[section])
353     \param key The key to look for (in cmap[section][subsection])
354     \param subsection The subsection to search in.
355 */
356 template <class T>
357     T get_parameter(const string& section,
358         const string& key,
359         const string& subsection="none"){
360     //Get the tokenized line in configuration file.
361     strvec tokenized=find_and_tokenize(section,key,subsection);
362     stringstream ss;
363     T result;
364     //The first string in tokenized must be key
365     if(tokenized[0]!=key)
366         throw string("lost_keyword:_"+key+"_when_tokenizing_line");
367     //vs[1] is the value of type T that should be returned
368     ss<<tokenized[1];
369     if(!(ss>>result))
370         throw string(key+"_has_invalid_type:_"+tokenized[1]);
371     //clear stringstream
372     ss.clear();
373     ss.str("");
374     return result;
375 }
376
377 /** Template function that returns a set of values from a key in
378     a
379     particular section. The type should explicitly defined
380     when using function. For example :
381     set<string> stringset=configuration.get_parameterset<string>
382         >("section","key");
383     The function is defined in the header file for portability
384     between different compilers.
385     \param section The section to search in (cmap[section])
386     \param key The key to look for (in cmap[section][subsection])
387     \param subsection The subsection to search in. */
388 template <class T>
389     set<T> get_parameterset(const string& section,

```



```

390                                     const string& key,
391                                     const string& subsection="none"){
392     set<T> result;
393     //Get the tokenized line in configuration file.
394     strvec tokenized=find_and_tokenize(section,key,subsection);
395     strvec::iterator it;
396     stringstream ss;
397
398     //First element in the list should be key
399     if(tokenized[0]!=key)
400         throw string("An_error_occured_when_getting_the_parameter_
401                       list_for_"+key);
402
403     //Erasing first element
404     tokenized.erase(tokenized.begin());
405
406     //The rest of the element should be of type T
407     for(it=tokenized.begin();it!=tokenized.end();it++){
408         ss<<*it;
409         T tmp;
410         if(!(ss>>tmp))
411             throw string(*it+"is_of_the_wrong_type..Noticed_when_
412                           getting_parameter_set_for_"+key);
413         result.insert(tmp);
414         ss.clear();
415         ss.str("");
416     }
417     return result;
418 }
419
420 /** Template function that will read a file that
421     contain pairs. Example content of file:
422
423     string1 int1
424     string2 int2
425     ....
426
427     The function will return a list of pointers
428     to pairs. For example for the above file
429     list<pair<string,int>*> will be returned.
430
431     Important: Remember to delete the pairs
432     you are done using them since this function
433     allocates memory to the pairs but it does
434     not free the memory when you are have finished
435     using the pairs. Another thing to remember
436     is that the reading will stop when one of the
437     values had an incorrect type. For example
438     when reading the file :
439
440     string1 int1
441     string2 int2
442     string3 string4
443
444     only the first two pairs will be read. Modify

```

```

443     the function to throw an exception if its needed. */
444     template <class A, class B>
445     list<pair<A,B>* > read_file(const string& file){
446         //open file
447         ifstream inFile(file.c_str());
448         //if file cant be opened
449         if(!inFile)
450             throw string("Invalid_file:_" + file);
451         //the list of pairs returned
452         list<pair<A,B>* > result;
453         //while there is still data
454         while(inFile){
455             A a;
456             B b;
457             if(inFile>>a && inFile>>b){
458                 pair<A,B>* p= new pair<A,B>;
459
460                 p->first=a;
461                 p->second=b;
462                 //push back the pointer to the pair
463                 result.push_back(p);
464             }
465         }
466         //close file
467         inFile.close();
468
469         return result;
470     }
471     /** Check if the key is valid. It looks in the config map
472         (the read config file) and checks if the key is there
473         \param section The section to search in (cmap[section])
474         \param key The key to look for (in cmap[section][subsection])
475         \param subsection The subsection to search in */
476     bool valid_key(const string& section,
477                  const string& key,
478                  const string& subsection="none");
479
480     ///Return a pointer to the default properties structure
481     properties* get_default_properties(){
482         return &simOptions.rDefaults;
483     }
484     ///Check if there is custom heuristics
485     bool custom_heuristics(){
486         return customHeuris;
487     }
488     ///Return a pointer to the heuristics map
489     strmap* get_heuristics(){
490         if(!customHeuris)
491             throw string("trying_to_get_custom_heuristics_but_none_have_
492                 been_read");
493         return &heuristic;
494     }
495     /** Return a pointer to the resource's properties structure.
496         If the resource does not have a unique properties
497         structure associated with it then a pointer to the

```

```

497     default properties structure is returned.
498     \param id The resource id of the properties to search for
499     */
500 properties* get_properties(const string& id);
501 /** Return a pointer to the saveOptions structed associated
502     between the two resources. The method returns what id1
503     save about id2.
504     \param id1 The resource that want to know what it should save
505     \param id2 The resource who's information might get saved by
506         id1 */
507
508 economy* get_economy(const string& id);
509 /*
510 Returns a pointer to the economy properties
511 */
512 saveOptions* get_save_options(const string& id1,const string& id2
513 );
514
515 /** Returns a pointer to a string set contain the resource id's
516 that a resource (id) should use as its (initial) reputation
517 group. The reputation goup is defined as:
518
519 if there are 100 resource users in the simulation,
520     numberofrefs=4
521 and id = u54 then the string set containing u55, u56, u57,
522     u58 is returned.
523
524 if there are 100 resource if numberofrefs=5 and id = u98 then
525     the
526 string set containing u99, u0, u1, u2, u3 is returned.
527
528 The string set is created (using new) so remember to
529 delete it when done using it.
530 \param id The resource whom then reputation group is intended
531 for. */
532 strset* get_rep_group(const string& id);
533
534 /** Returns a pointer to the mandatory structure. This structure
535     number users, sites, save path ect. */
536 mandatory* get_mandatory(){
537     return &simOptions;
538 }
539
540 /** The Bank Manipulating functions */
541
542 void Configuration::countcredit(bool check);
543
544 void deductmoney(string sid,string uid,double amount);
545
546 void setaccounts(string uid, double amount, bool initial);
547
548 //Manipulate the bank object and in particular the users accounts
549 void putsalaryonaccount(string id, double amount);

```

```

545
546 bool hascredit(string userid);
547
548 //printsout the accounts
549 void print_accounts();
550
551 /** Frontend function to the expression parser. This function
552     adds a newline to the string '\n' (for the C flex/bison calc.
553     l/calc.y)
554     and creates a char[] that is submitted to the C expression
555     parser.
556     Check the file calc.y - function setup() - for the math
557     functions
558     available as default. The result is then returned as a double
559     IMPORTANT : math.h uses the name log to the natural logarithm
560     (ln)
561     and log10 to the 10 base version. This is done to avoid
562     confusion with
563     the *nix program ln.
564     For example :
565
566     if string exp="ln(10)";
567     then parse_exp(exp)==log(10);
568
569     if string exp="log(10)";
570     then parse_exp(exp)==log10(10);
571
572     Make sure that you only use log when you mean log10
573     \param exp The expression to parse */
574 double parse_exp(string exp);
575
576 /** Calculated the number of shared groups between resource id1
577     and resource id2. */
578 int compare_groups(const string& id1,const string& id2);
579
580
581
582
583 /* methods for changing and adding resources to the simulation */
584 ///Check if there are more resources to be added at the current
585     time
586 bool more_additions(unsigned int ctime);
587
588 ///Return a pointer to the next addition structure(remember to
589     delete it)
590 addition* get_next_add();
591
592
593 ///Return a set of times where resources should be added
594 intset& get_addtimes(){
595     return addtimes;
596 }
597
598 ///Check if there are more resources to be change at the current
599     time
600 bool more_changes(unsigned int ctime);
601

```

```

591     ///Return a pointer to the next change structure(remember to
        delete it)
592     change* get_next_change();
593
594     ///Return a set of times where resources should change
595     intset& get_changetimes(){
596         return changetimes;
597     }
598
599     /*Various test functions*/
600
601
602     ///Prints a summary of the read config file
603     void print_summary();
604     ///print the entire read config (the cmap – used for testing)
605     void print_config();
606     ///print the content of a properties structure (used for testing)
607     void print_properties(const properties* p);
608     ///print the content of the resources map (containing the
        properties structures)
609     void print_resources();
610     ///print a string vector
611     void print_strvec(const strvec&);
612     ///print the difference between resource (id) properties p and
        the default properties
613     void print_difference(const string& id,const properties* p);
614 };
615 #endif

```

## D.7.2 configuration.cpp

```

1  #include "configuration.h"
2
3  Configuration::Configuration(const string& file){
4      commandlineSeed=false;
5      read_whole_config(file);
6      init();
7  }
8
9  Configuration::Configuration(const string& file,unsigned long int s
    ){
10     ///A commandline seed is being used
11     commandlineSeed=true;
12     ///Setting seed
13     simOptions.seed=s;
14     read_whole_config(file);
15     init();
16 }
17 void Configuration::init(){
18     set_mandatory_options();
19
20     set_resources();
21
22     set_economies();

```

```

23
24     set_additions();
25
26     set_changes();
27
28     set_saveoptions();
29
30     set_bank();
31
32     //if there is specified custom heuristics then used those
33     set_heuristic();
34     //Setup exp parser
35     setup();
36
37     //print summary
38     if(valid_key("[Config]", "printconfigsummary")){
39         string printsum=get_parameter<string>("[Config]", "
           printconfigsummary");
40         if(printsum=="yes" || printsum=="true"){
41             print_summary();
42         }
43     }
44 }
45
46 void Configuration::read_whole_config(const string& file){
47     string line;
48     string section="none";
49     string subsection="none";
50
51     //Open config file
52     ifstream inFile(file.c_str());
53
54     //Throw exception if can open file
55     if(!inFile)
56         throw string("Invalid_config_file:_"+ file);
57
58     //Read whole config into configfile map.
59     while(! inFile.eof() ){
60         //Read line
61         getline(inFile, line);
62         //First trim the line
63         line=trim(line);
64
65         //if line is not a comment
66         if(line[0]!=';'){
67             if(valid_section(line)){
68                 //line is a [Section]
69                 section=line.substr(line.find_first_of('[',0),
70                                   line.find_first_of(']',0)+1);
71                 //Reset subsection to "none"
72                 subsection ="none";
73             }
74             else if(valid_subsection(line)){
75                 //line is a <Subsection>
76                 subsection=line.substr(line.find_first_of('<',0),

```

```

77         line.find_first_of('>',0)+1);
78         //cout<<"subsection: "<<subsection<<endl;
79     }
80     else if(line.size()>0){
81         //Push back line if size>0. No reason to store empty lines
82         cfile[section][subsection].push_back(line);
83     }
84 }
85 }
86 //close file
87 inFile.close();
88 }
89 bool Configuration::valid_subsection(string& s){
90     string line=s;
91     line=trim(line);
92     /* The requirements for a valid subsection
93        is the it starts with the char '<'
94        and that there is a ending '>' somewhere. */
95     if(line[0]!='<')
96         return false;
97     if(line.find(">",0)!= string::npos)
98         return true;
99     else
100        return false;
101 }
102
103 bool Configuration::valid_section(string& s){
104     string line=s;
105     //remove whitespace
106     line=trim(line);
107     /* The requirements for a valid section
108        is the it starts with the char '['
109        and that there is a ending ']' somewhere. */
110     if(line[0]!='[')
111         return false;
112     if(line.find("]",0)!= string::npos)
113         return true;
114     else
115         return false;
116 }
117
118 bool Configuration::valid_key(const string& section,
119                               const string& key,
120                               const string& subsection){
121     //Check if section is defined in cfile map
122     if(cfile.find(section)==cfile.end())
123         return false;
124
125     //Check if subsection is defined in cfile map
126     if(cfile[section].find(subsection)==cfile[section].end())
127         return false;
128
129     //Find key in the string vector at cfile[section][subsection]
130     strvec::iterator it=cfile[section][subsection].begin();
131     while(it!=cfile[section][subsection].end()){

```

```

132     //if the line *it starts with key.
133     string line=*it;
134     if(trim(line).substr(0,key.size())==key)
135         return true;
136
137     it++;
138 }
139 return false;
140 }
141
142 bool Configuration::valid_id(const string& id){
143     //id must start with u or s
144     if(id[0]!='u' && id[0]!='s')
145         return false;
146     //the rest of the id must be a number
147     string rnumber=id.substr(1,id.size());
148     //Check number
149     int number;
150     stringstream ss;
151     ss<<rnumber;
152     if((ss>>number)){
153         return true;
154     }
155     //Clear stringstream
156     ss.clear();
157     ss.str("");
158     return false;
159 }
160
161
162 string& Configuration::get_next_token(const strvec& v, strvec::
    iterator& it){
163     /* Simple function that advances the iterator one step and checks
164        if the end is reached. If it is then a exception is thrown.
165        Else
166        a reference to the element is returned.
167
168        (The reason for this function is just to avoid explicitly
169        checking
170        if the end is reached each time.)
171     */
172     it++;
173     if(it==v.end())
174         throw string("trying_to_get_next_token ,_but_there_are_no_more")
175         ;
176     return *it;
177 }
178
179 void Configuration::set_economies(){
180     //Setting default
181     set_economy(ep);
182     string section="[Economy]";
183     // cout << "inde i set_economies" << endl;

```



```

183
184     for(submap::iterator it=cfile[section].begin();it!=cfile[
185         section].end();it++)
186         if(it->first!="none")
187             {
188                 economy test=ep;
189                 set_economy(test,section,it->first);
190
191                 if(compare_economy(test,ep)!=true)
192                 {
193                     string eid = (it->first).substr((it->first)
194                         .find("<",0)+1,
195                         (it->first).find(">",0)-1);
196
197                     //save the properties struct for
198                     economyresource eid
199                     economyresources[eid]=test;
200                 }
201             }
202
203
204 void Configuration::set_economy(economy& e, const string& section,
205     const string& subsection){
206     if(section!="[Economy]")
207     {
208         throw string ("Economy_parameters_was_not_specified
209             _in_section_"+section+"_(Might_not_be_a_default
210             _value_to_use_either)");
211     }
212
213     if(subsection=="none"){
214         e.markup=get_parameter<int>(section,"markup",
215             subsection);
216         e.depreciation=get_parameter<int>(section,"
217             depreciation",subsection);
218         e.discount=get_parameter<int>(section,"discount",
219             subsection);
220         e.protectionlevel=get_parameter<double>(section,"
221             protectionlevel",subsection);
222         e.hardwarecost=get_parameter<int>(section,"
223             hardwarecost",subsection);
224         e.sitecpus=get_parameter<int>(section,"cpus",
225             subsection);
226         //cout<<"Economy parameters set";
227     }
228
229     else
230     {
231         //Setting properties for a specific resources
232         if(valid_key(section,"markup",subsection))
233         {
234             e.markup=get_parameter<int>(section,"markup",subsection);

```

```

225     cout << "inde_i_" << section << "_med_" << subsection << "
        hvor_special_markup_kaldt" << endl;
226 }
227 if(valid_key(section,"depreciation",subsection))
228     e.depreciation=get_parameter<int>(section,"depreciation",
        subsection);
229
230 if(valid_key(section,"discount",subsection))
231     e.discount=get_parameter<int>(section,"discount",
        subsection);
232
233 if(valid_key(section,"protectionlevel",subsection))
234     e.protectionlevel=get_parameter<double>(section,"
        protectionlevel",subsection);
235
236 if(valid_key(section,"hardwarecost",subsection))
237     {
238         e.hardwarecost=get_parameter<int>(section,"
            hardwarecost",subsection);
239         cout << "inde_i_" << section << "_med_" <<
            subsection << "hvor_special_hardwarecost_kaldt"
            << endl;
240     }
241
242 if(valid_key(section,"cpus",subsection))
243     e.sitecpus=get_parameter<int>(section,"cpus",subsection);
244 }
245
246
247
248 }
249
250 void Configuration::set_properties(properties& p,
251                                 const string& section,
252                                 const string& subsection){
253     //sets property struct p with the values in section,subsection
254
255     if(section=="[Default]" && subsection=="none"){
256         //Setting default properties
257         p.i=get_parameter<int>(section,"i");
258         p.j=get_parameter<int>(section,"j");
259         p.groups=get_parameterset<int>(section,"groups");
260         p.threshold=get_parameter<double>(section,"threshold");
261         p.trust_eval=get_parameter<string>(section,"trust_eval");
262         p.rep_eval=get_parameter<string>(section,"rep_eval");
263         p.direct_weight=get_parameter<double>(section,"direct_weight");
264         p.smoothing=get_parameter<double>(section,"smoothing");
265         p.sigma=get_parameter<double>(section,"sigma");
266         p.stranger=get_parameter<double>(section,"stranger");
267         //Only used if discrete
268         if(p.trust_eval=="discrete")
269             p.discrete_ratio=get_parameter<double>(section,"
                discrete_ratio");
270         //Only used if beta
271         if(p.rep_eval=="beta"){

```

```

272     p.forgetting=get_parameter<double>(section,"forgetting");
273     p.base_weight=get_parameter<double>(section,"base_weight");
274 }
275 if(p.rep_eval=="discrete")
276     p.rep_ratio=get_parameter<double>(section,"rep_ratio");
277 }
278 else{
279     //Setting properties for a specific resources
280     if(valid_key(section,"i",subsection))
281         p.i=get_parameter<int>(section,"i",subsection);
282
283     if(valid_key(section,"j",subsection))
284         p.j=get_parameter<int>(section,"j",subsection);
285
286     if(valid_key(section,"groups",subsection)){
287         p.groups=get_parameterset<int>(section,"groups",subsection);
288     }
289
290     if(valid_key(section,"threshold",subsection))
291         p.threshold=get_parameter<double>(section,"threshold",
292             subsection);
293
294     if(valid_key(section,"trust_eval",subsection))
295         p.trust_eval=get_parameter<string>(section,"trust_eval",
296             subsection);
297
298     if(valid_key(section,"rep_eval",subsection))
299         p.rep_eval=get_parameter<string>(section,"rep_eval",
300             subsection);
301
302     if(valid_key(section,"direct_weight",subsection))
303         p.direct_weight=get_parameter<double>(section,"direct_weight",
304             subsection);
305
306     if(valid_key(section,"smoothing",subsection))
307         p.smoothing=get_parameter<double>(section,"smoothing",
308             subsection);
309
310     if(valid_key(section,"sigma",subsection))
311         p.sigma=get_parameter<double>(section,"sigma",subsection);
312
313     if(valid_key(section,"stranger",subsection))
314         p.stranger=get_parameter<double>(section,"stranger");
315
316     //Only used if discrete combination
317     if(p.trust_eval=="discrete")
318         if(valid_key(section,"discrete_ratio",subsection))
319             p.discrete_ratio=get_parameter<double>(section,"
320                 discrete_ratio",subsection);
321         else if(simOptions.rDefaults.trust_eval!="discrete")
322             throw string("discrete_ratio_was_not_set_in_subsection_"+
323                 subsection+"_(Might_not_be_a_default_value_to_use_
324                 either)");
325
326     //Only used if beta
327     if(p.rep_eval=="beta")

```

```

319     if (valid_key(section, "forgetting", subsection))
320         p.forgetting=get_parameter<double>(section, "forgetting",
            subsection);
321     else if (simOptions.rDefaults.rep_eval!="beta")
322         throw string("forgetting_was_not_specified_in_subsection_"+
            subsection+"_(Might_not_be_a_default_value_to_use_
            either)");
323
324     //Only used if beta
325     if (p.rep_eval=="beta")
326         if (valid_key(section, "base_weight", subsection))
327             p.base_weight=get_parameter<double>(section, "base_weight",
                subsection);
328         else if (simOptions.rDefaults.rep_eval!="beta")
329             throw string("base_weight_was_not_specified_in_subsection_"+
                subsection+"_(Might_not_be_a_default_value_to_use_
                either)");
330
331     //Only used if discrete reputation
332     if (p.rep_eval=="discrete")
333         if (valid_key(section, "rep_ratio", subsection))
334             p.rep_ratio=get_parameter<double>(section, "rep_ratio",
                subsection);
335         else if (simOptions.rDefaults.rep_eval!="discrete")
336             throw string("rep_ratio_was_not_set_in_subsection_"+
                subsection+"_(Might_not_be_a_default_value_to_use_
                either)");
337
338     }
339 }
340 void Configuration::set_resources(const string& section){
341     if(section!="[Resources]")
342         throw string("Trying_to_set_resources_from_section:_" +section)
            ;
343
344     for(submap::iterator it=cfile[section].begin();it!=cfile[section
        ].end();it++)
345         if(it->first!="none"){
346             //Setting properties for <it->first>
347             properties p= simOptions.rDefaults;
348             set_properties(p, section, it->first);
349             if(!compare_properties(p, simOptions.rDefaults)){
350                 //<id> has custom properties. Finding id
351                 string rid = (it->first).substr((it->first).find("<" ,0)+1,
352                     (it->first).find(">" ,0)-1);
353                 //save the properties struct for resource rid
354                 resources[rid]=p;
355             }
356         }
357 }
358 void Configuration::set_changes(const string& section ){
359     if(section!="[Resources]")
360         throw string("Trying_to_set_resources_from_section:_" +section)
            ;
361

```

```

362     string key ="time";
363     strvec tokenized;
364     stringstream ss;
365     submap::iterator subit;
366     strvec::iterator vit;
367     for(subit=cfile[section].begin();subit!=cfile[section].end();
368         subit++){
369         if(subit->first!="none"){
370             /* Processing subsection <it->first>. Will check if the
371             key ("time") is there. If the key is found then i and j
372             is also mandatory. I and j are the change in behaviour. */
373             for(vit=subit->second.begin();vit!=subit->second.end();vit++)
374                 if(trim(*vit).substr(0,key.size())==key){
375                     //resource id
376                     string rid=(subit->first).substr((subit->first).find("<"
377                                                         ,0)+1,
378                                                         (subit->first).find(">"
379                                                         ,0)-1);
380
381                     //Set up for tokenizing
382                     strvec tokenized;
383                     string::size_type pos = vit->find_first_of(';');
384                     string tokenizeme= *vit;
385                     if(pos != string::npos)
386                         tokenizeme=vit->substr(0,pos);
387
388                     //tokenizing
389                     tokenize(tokenizeme,tokenized);
390
391                     //tokenize iterator
392                     strvec::iterator tokit;
393                     /* Find the time of the change */
394                     tokit=find(tokenized.begin(),tokenized.end(),"time");
395                     if(tokit==tokenized.end())
396                         throw string("The_key_time_was_not_found_in_the_line:_"+
397                                     +tokenizeme);
398                     /* The location tokit+1 must contain a unsigned int that
399                     denotes
400                     the time of the change. */
401
402                     unsigned int changeTime;
403                     ss<<get_next_token(tokenized,tokit);
404                     if(!(ss>>changeTime))
405                         throw string(*tokit+"_is_not_of_type_unsigned_int_in_
406                                     line_"+tokenizeme);
407
408                     //Clear stringstream
409                     ss.clear();
410                     ss.str("");
411
412                     /* Find i and j values */
413                     // i
414                     int r_i;
415                     tokit=find(tokenized.begin(),tokenized.end(),"i");
416                     if(tokit!=tokenized.end()){
417                         //the i key was found.

```

```

411         ss<<get_next_token(tokenized,tokit);
412         if(!(ss>>r_i))
413             throw string("No valid i value was found in "+*tokit)
414         ;
415     }
416     else
417         throw string("Missing i value was in "+tokenizeme);
418     //Clear stringstream
419     ss.clear();
420     ss.str("");
421
422     // j
423     int r_j;
424     tokit=find(tokenized.begin(),tokenized.end(),"j");
425     if(tokit!=tokenized.end()){
426         //the j key was found.
427         ss<<get_next_token(tokenized,tokit);
428         if(!(ss>>r_j))
429             throw string("No valid i value was found in "+*tokit)
430         ;
431     }
432     else
433         throw string("Missing j value was in "+tokenizeme);
434     //Clear stringstream
435     ss.clear();
436     ss.str("");
437
438     //Making change struct and add it to change queue
439     change* c=new change;
440     (*c).id=rid;
441     (*c).time=changeTime;
442     (*c).i=r_i;
443     (*c).j=r_j;
444
445     changes.push(c);
446     changetimes.insert((*c).time);
447
448     //Add pointer to listofchanges (used for summary)
449     listofchanges.push_back(c);
450 }
451
452 void Configuration::set_additions(const string& section){
453     if(section!="[Add]")
454         throw string("Trying to set additions from section: "+section)
455     ;
456
457     stringstream ss;
458     submap::iterator sit;
459     strvec::iterator it;
460     for(sit=cfile[section].begin();sit!=cfile[section].end();sit++){
461         //for each subsection
462         for(it=sit->second.begin();it!=sit->second.end();it++){

```

```

462     /* Each *it should be a line containing information of a
463         addition.
464         This line will be tokenized and then addition structs will
465             be created
466             and inserted into the addition queue. */
465     strvec tokenized;
466     string::size_type pos = it->find_first_of(';');
467     string tokenizeme = *it;
468     if(pos != string::npos)
469         tokenizeme=it->substr(0,pos);
470
471     //tokenizing
472     tokenize(tokenizeme,tokenized);
473
474     strvec::iterator tokit;
475     /* Some keys must be present in the line (in the string
476         vector
477         containing the tokenized line). The first mandatory key is
478             "time".
479         This key denotes when the addition will happen. */
478     tokit=find(tokenized.begin(),tokenized.end(),"time");
479     if(tokit==tokenized.end())
480         throw string("The_key_time_was_not_found_in_the_line:"+
481             tokenizeme);
482     /* The location tokit+1 must contain a unsigned int that
483         denotes
484         the time of the addition. */
485
486     unsigned int addTime;
487     ss<<get_next_token(tokenized,tokit);
488     if(!(ss>>addTime))
489         throw string("Expected_"+*tokit+"_to_be_of_type_unsigned_
490             int_when_processing_line_"+tokenizeme);
491
492     //Clear stringstream
493     ss.clear();
494     ss.str("");
495     /* A mandatory value that also must be there is a valid
496         resource id.
497         This is required simple to remind the user to keep track
498             of who joins
499             the simulation (for saving purposes). Here we will iterate
500             through
501             the strvec and see if a string begins with 'u' or 's'.*/
502     bool foundrid=false;
503     string rid;
504     for(tokit=tokenized.begin();tokit!=tokenized.end() && !
505         foundrid;tokit++)
506         if(valid_id(*tokit)){
507             rid=*tokit;
508             foundrid=true;
509         }
510
511     //Making addition struct and adding to addqueue and adding
512     time to addtimes

```

```

505     addition* a=new addition;
506     (*a).id=rid;
507     (*a).time=addTime;
508     additions.push(a);
509     addtimes.insert((*a).time);
510
511     /* Now the time of addition has been found and the id of the
512        resource. Now its time to check for optional values.
513        It is however recommended (purely for consistency of the
514        configuration file) to specify optional values in the
515        [Resource] section. The method here is just for shorthand.
516
517        The only optional values allowed are i & j. Meaning that
518        only the behavioural values can be specified.
519
520        REMEMBER : Any option values specified in the [Add]
521                   section
522        will override values from the [Resource] section. */
523
524     //Check for i value
525     bool specified_i=false;
526     int r_i;
527     tokit=find(tokenized.begin(),tokenized.end(),"i");
528     if(tokit!=tokenized.end()){
529         //the i key was found.
530         ss<<get_next_token(tokenized,tokit);
531         if(!(ss>>r_i))
532             throw string("No_valid_i_value_was_found_when_in_"+*tokit
533                );
534         specified_i=true;
535     }
536     //Clear stringstream
537     ss.clear();
538     ss.str("");
539
540     //Check for j value
541     bool specified_j=false;
542     int r_j=-1;
543     tokit=find(tokenized.begin(),tokenized.end(),"j");
544     if(tokit!=tokenized.end()){
545         //the i key was found.
546         ss<<get_next_token(tokenized,tokit);
547         if(!(ss>>r_j))
548             throw string("No_valid_i_value_was_found_when_in_"+*tokit
549                );
550         specified_j=true;
551     }
552     //Clear stringstream
553     ss.clear();
554     ss.str("");
555
556     /* Update the resources map if i and j was specified.
557        Here a warning will be printed if it overwrites
558        previously defined values.
559        */

```



```

557     if(specified_i || specified_j){
558         //i or j was specified.
559         if(resources.find(rid)==resources.end()){
560             /*No values have previously been specified
561             Now we can just make a new properties struct
562             and insert. Base is still rDefaults */
563             properties p =simOptions.rDefaults;
564             if(specified_i)
565                 p.i=r_i;
566             if(specified_j)
567                 p.j=r_j;
568
569             resources[rid]=p;
570         }
571         else{
572             /* properties struct already exist for resource rid.
573             We need to change it. (overwrite i and j)
574             Will print a warning if different from default values
575             */
576             if(r_i!=simOptions.rDefaults.i || r_j!=simOptions.
                    rDefaults.j)
577                 if(resources[rid].i!=simOptions.rDefaults.i ||
578                     resources[rid].j!=simOptions.rDefaults.j)
579                     cout<<"Warning: _Overwriting _i _and/or _j _values _for _"<<
                        rid<<endl;
580             if(specified_i)
581                 resources[rid].i=r_i;
582             if(specified_j)
583                 resources[rid].j=r_j;
584         }
585     }
586 }
587 //Add addition* to listofadds (when printing the summary)
588 listofadds.push_back(a);
589
590 }
591 }
592
593 //cout<<"set_additions ([Add]) success"<<endl;
594 }
595
596 void Configuration::set_saveoptions(){
597     /* This method sets the saveoptions. It will
598     look for 3 sections in the cfile. The sections
599     are [CtSave], [DtSave] and [RtSave], [PrSave].
600
601     IMPORTANT: If a section is missing the default action
602     is to save nothing.
603     */
604     bool ct=false,dt=false,rt=false, pr=false;
605     //Find default saving options
606     if(valid_key("[CtSave]","default")){
607         string save=get_parameter<string>("[CtSave]","default");
608         if(save=="yes" || save=="true")
609             ct=true;

```

```

610 }
611 if(valid_key("[DtSave]","default")){
612     string save=get_parameter<string>("[DtSave]","default");
613     if(save=="yes" || save=="true")
614         dt=true;
615 }
616 if(valid_key("[RtSave]","default")){
617     string save=get_parameter<string>("[RtSave]","default");
618     if(save=="yes" || save=="true")
619         rt=true;
620 }
621 if(valid_key("[PrSave]","default")){
622     string save=get_parameter<string>("[PrSave]","default");
623     if(save=="yes" || save=="true")
624         pr=true;
625 }
626 if(ct||dt||rt||pr)
627     cout<<"Something is saved by default"<<endl;
628 //Create saveOptions object
629 sOptions=new SaveOption(ct,dt,rt);
630 //scan for [CtSave],[DtSave] and [RtSave] sections
631 set_savetype("[CtSave]","ct");
632 set_savetype("[DtSave]","dt");
633 set_savetype("[RtSave]","rt");
634 }
635
636 void Configuration::set_savetype(const string& section,const string
        & type){
637     /* Helper function to set_saveoptions */
638     if(type!="ct" && type!="dt" && type!="rt" && type!="pr")
639         throw string("using_set_savetype_with_invalid_type:_" +type);
640
641     string subsection="none";
642     string nokey="default";
643     if(cfile.find(section)!=cfile.end()){
644         //There is a [savetype] section.
645         for(strvec::iterator it=cfile[section][subsection].begin();
646             it!=cfile[section][subsection].end();it++)
647             if(trim(*it).substr(0,nokey.size())!=nokey){
648                 /* *it might contain a string of a key (a resource id)
649                    and a list of resource id's. These resource id's are
650                    the resources that the resource with id = key will
651                    save combined trust info on.
652                 */
653                 string::size_type pos = it->find_first_of(';');
654                 string tokenizeme= *it;
655                 if(pos != string::npos)
656                     tokenizeme=it->substr(0,pos);
657
658                 //tokenizing
659                 strvec tokenized;
660                 tokenize(tokenizeme,tokenized);
661
662                 //id is the resource that will save information
663                 string id;

```

```

664         if(valid_id(tokenized[0]))
665             id=tokenized[0];
666         else
667             throw string(id+"_appears_to_be_an_invalid_id_key_(["
668                 CtSave]_)"");
669
670         //Check that every token appears as a valid resource id
671         for(strvec::iterator tokit=tokenized.begin();tokit!=
672             tokenized.end();tokit++){
673             if(*tokit=="full"){
674                 //Saves everything
675                 if(type=="ct")
676                     sOptions->add_full_save(id,type);
677                 else if(type=="dt")
678                     sOptions->add_full_save(id,type);
679                 else if(type=="rt")
680                     sOptions->add_full_save(id,type);
681                 else if(type=="pr")
682                     sOptions->add_full_save(id,type);
683                 else
684                     throw string("Unknown_type:_" +type);
685             }
686             else if(!valid_id(*tokit))
687                 throw string (*tokit+"_appears_invalid_in_"+id+"'_s_"
688                     ctsave_list");
689             else if(id!=*tokit){
690                 //Add *tokit at saveoptions.
691                 //cout<<id<<" saves "<<*tokit<<" of type"<<type<<endl;
692                 if(type=="ct")
693                     sOptions->save_ct(id,*tokit);
694                 else if(type=="dt")
695                     sOptions->save_dt(id,*tokit);
696                 else if(type=="rt")
697                     sOptions->save_rt(id,*tokit);
698                 else if(type=="pr")
699                     sOptions->save_rt(id,*tokit);
700                 else
701                     throw string("Unknown_type:_" +type);
702             }
703         }
704     }
705 }
706
707 void Configuration::set_mandatory_options(){
708     /* get and set the mandatory options.
709     These options are the bare minimum needed to
710     run a simulation.
711     */
712     try{
713         simOptions.users=get_parameter<unsigned int>("[ Config]","users"
714             );

```

```

714     simOptions.sites=get_parameter<unsigned int>("[ Config]","sites"
715     );
716     simOptions.time=get_parameter<unsigned int>("[ Config]","time");
717     simOptions.mu=get_parameter<double>("[ Config]","mu");
718     simOptions.normalmu=get_parameter<double>("[ Config]","normalmu"
719     );
720     simOptions.normalsigma=get_parameter<double>("[ Config]","
721     normalsigma");
722     simOptions.path=get_parameter<string>("[ Config]","path");
723     //If there is no seed specified at commandline seed
724     if(!commandlineSeed){
725         //Check if there is a seed specified in cmap. If not
726         //current time will be used as seed.
727         if(valid_key("[ Config]","seed"))
728             simOptions.seed=get_parameter<unsigned long int>("[ Config]"
729             ,"seed");
730         else
731             simOptions.seed=time(0);
732     }
733     else if(valid_key("[ Config]","seed")){
734         cout<<"Warning: _the_seed_"<<simOptions.seed
735         <<"_was_specified_at_the_commandline,_but_the_config_file"
736         <<"_suggests_the_seed_"
737         <<get_parameter<unsigned long int>("[ Config]","seed")<<
738         endl;
739     }
740     //Default behaviour for the resources
741     set_properties(simOptions.rDefaults);
742 }
743 //cout<<"set_mandatory_options success"<<endl;
744 }
745 catch(string e){
746     string exc;
747     exc="When_trying_to_set_the_mandatory_options";
748     exc+="_the_following_error_occured\n"+e;
749     throw string(exc);
750 }
751 }
752 void Configuration::set_heuristic(){
753     customHeuris=false;
754 }
755 if(valid_key("[ Heuristics]","experience")){
756     heuristic["experience"]=get_parameter<string>("[ Heuristics]","
757     experience");
758     cout<<"Info:_Custom_experience_heuristics_specified"<<endl;
759     customHeuris=true;
760 }
761 if(valid_key("[ Heuristics]","oldvalue")){
762     heuristic["oldvalue"]=get_parameter<string>("[ Heuristics]","
763     oldvalue");
764     cout<<"Info:_Custom_oldvalue_heuristics_specified"<<endl;
765     customHeuris=true;
766 }

```

```

761     if(valid_key("[Heuristics]", "newvalue")){
762         heuristic["newvalue"]=get_parameter<string>("[Heuristics]", "
           newvalue");
763         cout<<" Info: _Custom_newvalue_heuristics_specified"<<endl;
764         customHeuris=true;
765     }
766 }
767 }
768
769
770 void Configuration::tokenize(const string& str,
771                             vector<string>& tokens,
772                             const string& delimiters){
773     // Skip delimiters at beginning.
774     string::size_type lastPos = str.find_first_not_of(delimiters, 0);
775     // Find first "non-delimiter".
776     string::size_type pos      = str.find_first_of(delimiters, lastPos
777 );
778
779     while (string::npos != pos || string::npos != lastPos){
780         // Found a token, add it to the vector.
781         tokens.push_back(str.substr(lastPos, pos - lastPos));
782         // Skip delimiters. Note the "not_of"
783         lastPos = str.find_first_not_of(delimiters, pos);
784         // Find next "non-delimiter"
785         pos = str.find_first_of(delimiters, lastPos);
786     }
787 }
788
789 string Configuration::trim(string& s, const string& drop)
790 {
791     string r=s.erase(s.find_last_not_of(drop)+1);
792     return r.erase(0,r.find_first_not_of(drop));
793 }
794
795 bool Configuration::compare_properties(const properties& p1, const
796 properties& p2){
797     if(p1.i!=p2.i)
798         return false;
799     if(p1.j!=p2.j)
800         return false;
801     if(p1.threshold!=p2.threshold)
802         return false;
803     if(p1.trust_eval!=p2.trust_eval)
804         return false;
805     if(p1.rep_eval!=p2.rep_eval)
806         return false;
807     if(p1.smoothing!=p2.smoothing)
808         return false;
809     if(p1.sigma!=p2.sigma)
810         return false;
811     if(p1.stranger!=p2.stranger)
812         return false;
813     if(p1.trust_eval=="discrete")
814         if(p1.discrete_ratio!=p2.discrete_ratio)
815             return false;

```

```

813     if(p1.rep_eval=="beta")
814         if(p1.forgetting!=p2.forgetting)
815             return false;
816         else if(p1.base_weight!=p2.base_weight)
817             return false;
818     if(p1.rep_eval=="discrete")
819         if(p1.rep_ratio!=p2.rep_ratio)
820             return false;
821     if(p1.groups!=p2.groups)
822         return false;
823
824     //Properties seem equal
825     return true;
826 }
827
828 bool Configuration::compare_economy(const economy& e1,const economy
& e2){
829     if(e1.markup!=e2.markup)
830         return false;
831     if(e1.depreciation!=e2.depreciation)
832         return false;
833     if(e1.discount!=e2.discount)
834         return false;
835     if(e1.hardwarecost!=e2.hardwarecost)
836         return false;
837     if(e1.protectionlevel!=e2.protectionlevel)
838         return false;
839
840     //Properties seem equal
841     return true;
842 }
843
844
845 strvec Configuration::find_and_tokenize(const string& section,
846                                         const string& key,
847                                         const string& subsection,
848                                         const string& delim){
849     //String vector for the tokenized result
850     strvec result;
851     //Iterators
852     strvec::iterator it;
853     //Flag indicated if key is found
854     bool foundKey= false;
855
856     //Check if section is defined in cfile map
857     if(cfile.find(section)==cfile.end())
858         throw string(section + "_section_was_not_found_in_config_file")
;
859
860     //Check if subsection is defined in cfile map
861     if(cfile[section].find(subsection)==cfile[section].end())
862         throw string(subsection + "_subsection_was_not_found_in_config_
file");
863
864     //Find key in the string vector at cfile[section][subsection]

```

```

865     it=cfile[section][subsection].begin();
866     while(!foundKey && it!=cfile[section][subsection].end()){
867         // *it must start with key
868         if(trim(*it).substr(0,key.size())==key){
869             /* Found keyword. Tokenizing the line. Only the line
870                upto the the first comment, denoted by ; is parsed. */
871             string::size_type pos = it->find_first_of(';');
872             string tokenizeme= *it;
873             if(pos != string::npos)
874                 tokenizeme=it->substr(0,pos);
875
876             //tokenizing
877             tokenize(tokenizeme,result,delim);
878             //Indicating that the key was found
879             foundKey=true;
880         }
881         it++;
882     }
883     if(!foundKey)
884         throw string(key+" _was_not_found_in_section_"+section);
885
886     //Return tokenized string as a vector of strings
887     return result;
888 }
889
890 properties* Configuration::get_properties(const string& id){
891     if(resources.find(id)==resources.end())
892         return &simOptions.rDefaults;
893     else
894         return &resources[id];
895 }
896
897 economy* Configuration::get_economy(const string& id){
898     //const string& id){
899     if(economyresources.find(id)!=economyresources.end())
900         return &economyresources[id];
901     else
902         return &ep;
903 }
904
905
906
907 void Configuration::set_bank(){
908     double startmoney=0.0;
909
910
911     for(submap::iterator it=cfile["[Bank]"].begin(); it!=cfile["
912         [Bank]"].end(); it++)
913     {
914         if(it->first!="none")
915             {
916                 string bid = (it->first).substr((it->first).find("<
917                     ",0)+1,

```

```

916                                     (it->first).find(">
917                                     ",0)-1);
918     startmoney=get_parameter<double>("[Bank]", "
919     initialmoney",(it->first));
920     nationaltreasure.Set_Accounts(bid, startmoney, false
921     );
922     }
923 }
924 }
925 startmoney=get_parameter<double>("[Bank]", "initialmoney");
926
927     nationaltreasure.Set_init_money(startmoney);
928     cout << "BANK.OPRETTET" << endl;
929
930
931 }
932 }
933
934 void Configuration::putsalaryonaccount(string id, double amount)
935 {
936     nationaltreasure.PutSalaryOnAccount(id, amount);
937 }
938
939 bool Configuration::hascredit(string userid){
940     if(nationaltreasure.Has_Credit(userid))
941     return true;
942     else
943     return false;
944 }
945
946 saveOptions* Configuration::get_save_options(const string& id1,
947     const string& id2){
948     return sOptions->get_save_options(id1, id2);
949 }
950
951 strset* Configuration::get_rep_group(const string& rId){
952     //Check valid rId
953     if(rId[0]!='u' && rId[0]!='s')
954     throw string("Tried_to_get_a_reputation_group_for_"+rId);
955
956     strset* res= new strset();
957
958     stringstream ss;
959
960     //get rId's number
961     int myNr;
962     ss<<rId.substr(1,rId.length()-1);

```



```

966     if (!(ss >> myNr))
967         throw string("Invalid_resource_id_" + rId);
968
969     // Clear stringstream
970     ss.clear();
971     ss.str("");
972
973     // number of reputation group members
974     int numberofrefs=4;
975     if (valid_key("[Config]", "numberofrefs"))
976         numberofrefs=get_parameter<int>("[Config]", "numberofrefs");
977
978     // find max amount of resources
979     int max;
980     if (rId[0] == 'u')
981         max=simOptions.users-1;
982     else
983         max=simOptions.sites-1;
984
985     // counter
986     int c=myNr;
987     vector<int> nrs;
988     if (c==numberofrefs)
989         while(nrs.size()<numberofrefs && nrs.size()<=max){
990             c--;
991             if (c!=myNr)
992                 nrs.push_back(c);
993         }
994     else if (c>numberofrefs)
995         while(nrs.size()<numberofrefs && nrs.size()<=max){
996             c--;
997             if (c!=myNr)
998                 nrs.push_back(c);
999         }
1000     else{
1001         for(int i=0; i<=numberofrefs && i<=max ; i++){
1002             if (i!=myNr)
1003                 nrs.push_back(i);
1004         }
1005
1006         // create id's and insert
1007         for(int i=0; i<nrs.size(); i++){
1008             ss<<rId[0]<<nrs[i];
1009             res->insert(ss.str());
1010
1011             ss.clear();
1012             ss.str("");
1013         }
1014         // print
1015         //cout<<"For "<<rId<<" - ";
1016         //copy(res->begin(), res->end(), ostream_iterator<string>(cout, "
1017         //cout<<endl;
1018         return res;
1019

```

```

1020 }
1021
1022 int Configuration::compare_groups(const string& id1,const string&
    id2){
1023     //if(1)
1024     //return 1;
1025
1026     if(resources.find(id1)==resources.end() && resources.find(id2)==
        resources.end())
1027         return simOptions.rDefaults.groups.size();
1028     else{
1029         intset* g1;
1030         intset* g2;
1031         if(resources.find(id1)!=resources.end())
1032             g1=&resources[id1].groups;
1033         else
1034             g1=&simOptions.rDefaults.groups;
1035         if(resources.find(id2)!=resources.end())
1036             g2=&resources[id2].groups;
1037         else
1038             g2=&simOptions.rDefaults.groups;
1039         //make the set intersection for find the number of share groups
1040         intset r;
1041         set_intersection(g1->begin(),g1->end(),
1042             g2->begin(),g2->end(),
1043             inserter(r,r.begin()));
1044
1045         //Comparing
1046         return r.size();
1047     }
1048 }
1049 double Configuration::parse_exp(string e){
1050     string exp=e;
1051     //The expression has to end with a newline
1052     exp+="\n";
1053     //make char*
1054     char* expression=new char[exp.size()+1];
1055     //copy chars to expression
1056     strcpy(expression,exp.c_str());
1057     double result= parse_this(expression);
1058     //delete space allocate to expression
1059     delete[] expression;
1060     //Check if the parsing went ok.
1061     if(error_flag()!=0)
1062         throw string("Error_parsing_expression:" +e);
1063     return result;
1064 }
1065 bool Configuration::more_additions(unsigned int ctime){
1066     if(additions.empty())
1067         return false;
1068     if(additions.top()->time!=ctime)
1069         return false;
1070     else
1071         return true;
1072 }

```

```

1073
1074 addition* Configuration::get_next_add(){
1075     if(additions.empty())
1076         throw string("Asked_for_one_more_addition_but_none_existed_in_
            queue");
1077
1078     addition* a;
1079     a=additions.top();
1080     additions.pop();
1081     //increment users or sites counter
1082     if((*a).id[0]=='u')
1083         simOptions.users++;
1084     else if((*a).id[0]=='s')
1085         simOptions.sites++;
1086
1087     return a;
1088 }
1089
1090 bool Configuration::more_changes(unsigned int ctime){
1091     if(changes.empty())
1092         return false;
1093     if(changes.top()->time!=ctime)
1094         return false;
1095     else
1096         return true;
1097 }
1098
1099 change* Configuration::get_next_change(){
1100     if(changes.empty())
1101         throw string("Asked_for_one_more_addition_but_none_existed_in_
            queue");
1102
1103     change* c;
1104     c=changes.top();
1105     changes.pop();
1106     return c;
1107 }
1108
1109
1110 /* ----- Test functions ----- */
1111
1112
1113 void Configuration::print_config(){
1114     cmap::iterator it;
1115     submap::iterator subit;
1116     strvec::iterator vit;
1117     for(it=cfile.begin();it!=cfile.end();it++)
1118         for(subit=(it->second).begin();subit!=(it->second).end();subit
            ++){
1119             cout<<"printing_section:"<<it->first
1120                 <<"_and_subsection_"<<subit->first<<endl;
1121             for(vit=(subit->second).begin();vit!=(subit->second).end();
                vit++)
1122                 cout<<*vit<<endl;
1123         }

```

```

1124 }
1125
1126 void Configuration::print_properties(const properties* p){
1127
1128     cout<<"printing_properties"<<endl;
1129     cout<<"i"<<p->i<<endl;
1130     cout<<"j"<<p->j<<endl;
1131     cout<<"groups=";
1132     copy(p->groups.begin(), p->groups.end(), ostream_iterator<int>(
        cout, "_"));
1133     cout<<endl;
1134     cout<<"threshold"<<p->threshold<<endl;
1135     cout<<"trust_eval"<<p->trust_eval<<endl;
1136     cout<<"rep_eval"<<p->rep_eval<<endl;
1137     cout<<"smoothing"<<p->smoothing<<endl;
1138     cout<<"sigma"<<p->sigma<<endl;
1139     cout<<"stranger"<<p->stranger<<endl;
1140     if(p->trust_eval=="discrete")
1141         cout<<"discrete_ratio"<<p->discrete_ratio<<endl;
1142     if(p->rep_eval=="beta"){
1143         cout<<"forgetting"<<p->forgetting<<endl;
1144         cout<<"base_weight"<<p->base_weight<<endl;
1145     }
1146     if(p->rep_eval=="discrete")
1147         cout<<"rep_ratio"<<p->rep_ratio<<endl;
1148 }
1149
1150 void Configuration::print_resources(){
1151     cout<<"Default_properties:_"<<endl;
1152     print_properties(&simOptions.rDefaults);
1153     for(romap::iterator it=resources.begin();it!=resources.end();it
        ++){
1154         cout<<"resource_"<<it->first<<"_has_the_properties:_"<<endl;
1155         print_properties(&it->second);
1156     }
1157 }
1158
1159 void Configuration::print_strvec(const strvec& v){
1160     cout<<"Content_of_strvec:_";
1161     copy(v.begin(), v.end(), ostream_iterator<string>(cout, ",_"));
1162     cout<<endl;
1163 }
1164
1165 void Configuration::print_difference(const string& id,const
    properties* p){
1166     if(!compare_properties(*p,simOptions.rDefaults)){
1167         cout<<id<<"_has_special_properties"<<endl;
1168         if(p->i!=simOptions.rDefaults.i)
1169             cout<<"i_is_"<<p->i<<"_and_the_default_is_"
                <<simOptions.rDefaults.i<<endl;
1170         if(p->j!=simOptions.rDefaults.j)
1171             cout<<"j_is_"<<p->j<<"_and_the_default_is_"
                <<simOptions.rDefaults.j<<endl;
1172         if(p->threshold!=simOptions.rDefaults.threshold)
1173             cout<<"threshold_is_"<<p->threshold<<"_and_the_default_is_"

```

```

1176         <<simOptions.rDefaults.threshold<<endl;
1177     if(p->trust_eval!=simOptions.rDefaults.trust_eval)
1178         cout<<" trust_eval_is_"<<p->trust_eval<<" _and_the_default_is_"
1179         <<simOptions.rDefaults.trust_eval<<endl;
1180     if(p->rep_eval!=simOptions.rDefaults.rep_eval)
1181         cout<<" rep_eval_is_"<<p->rep_eval<<" _and_the_default_is_"
1182         <<simOptions.rDefaults.rep_eval<<endl;
1183     if(p->smoothing!=simOptions.rDefaults.smoothing)
1184         cout<<" smoothing_is_"<<p->smoothing<<" _and_the_default_is_"
1185         <<simOptions.rDefaults.smoothing<<endl;
1186     if(p->sigma!=simOptions.rDefaults.sigma)
1187         cout<<" sigma_is_"<<p->sigma<<" _and_the_default_is_"
1188         <<simOptions.rDefaults.sigma<<endl;
1189     if(p->stranger!=simOptions.rDefaults.stranger)
1190         cout<<" stranger_is_"<<p->smoothing<<" _and_the_default_is_"
1191         <<simOptions.rDefaults.smoothing<<endl;
1192     if(p->trust_eval=="discrete")
1193         if(simOptions.rDefaults.trust_eval=="discrete"){
1194             if(p->discrete_ratio!=simOptions.rDefaults.discrete_ratio)
1195                 cout<<" discrete_ratio_is_"<<p->discrete_ratio<<" _and_the_
1196                     default_is_"
1197                     <<simOptions.rDefaults.discrete_ratio<<endl;
1198         }
1199     else
1200         cout<<" discrete_ratio_is_"<<p->discrete_ratio
1201         <<" _and_there_is_no_default"<<endl;
1202     if(p->rep_eval=="beta")
1203         if(simOptions.rDefaults.rep_eval=="beta"){
1204             if(p->forgetting!=simOptions.rDefaults.forgetting)
1205                 cout<<" forgetting_is_"<<p->forgetting<<" _and_the_default_
1206                     is_"
1207                     <<simOptions.rDefaults.forgetting<<endl;
1208         }
1209     else
1210         cout<<" forgetting_is_"<<p->forgetting
1211         <<" _and_there_is_no_default"<<endl;
1212     if(p->rep_eval=="beta")
1213         if(simOptions.rDefaults.rep_eval=="beta"){
1214             if(p->base_weight!=simOptions.rDefaults.base_weight)
1215                 cout<<" forgetting_is_"<<p->base_weight<<" _and_the_default_
1216                     is_"
1217                     <<simOptions.rDefaults.base_weight<<endl;
1218         }
1219     else
1220         cout<<" base_weight_is_"<<p->base_weight
1221         <<" _and_there_is_no_default"<<endl;
1222     if(p->rep_eval=="discrete")
1223         if(simOptions.rDefaults.rep_eval=="discrete"){
1224             if(p->rep_ratio!=simOptions.rDefaults.rep_ratio)
1225                 cout<<" rep_ratio_is_"<<p->rep_ratio<<" _and_the_default_is_
1226                     _"
1227                     <<simOptions.rDefaults.rep_ratio<<endl;

```

```

1227     }
1228     else
1229         cout<<"rep_ratio_is_"<<p->rep_ratio
1230         <<"_and_there_is_no_default"<<endl;
1231
1232     //compare groups
1233     if(p->groups!=simOptions.rDefaults.groups){
1234         cout<<"Groups_are:_";
1235         copy(p->groups.begin(), p->groups.end(),
1236             ostream_iterator<int>(cout, "_"));
1237         cout<<endl;
1238         cout<<"Default_groups_are:_";
1239         copy(simOptions.rDefaults.groups.begin(),
1240             simOptions.rDefaults.groups.end(),
1241             ostream_iterator<int>(cout, "_"));
1242         cout<<endl;
1243     }
1244 }
1245 }
1246
1247 void Configuration::countcredit(bool check){
1248     //cout << " det virker" << endl;
1249     nationaltreasure.count_creditrejected(check);
1250 }
1251
1252 void Configuration::deductmoney(string sid,string uid,double amount
1253     )
1254 {
1255     nationaltreasure.Transfer_money_user_to_site(sid,uid,amount);
1256 }
1257 void Configuration::setaccounts(string uid, double amount, bool
1258     initial)
1259 {
1260     nationaltreasure.Set_Accounts(uid, amount,initial);
1261     //cout << "account " << uid << " sat" << endl;
1262 }
1263 }
1264
1265 void Configuration::print_accounts()
1266 {
1267     nationaltreasure.Print_accounts();
1268 }
1269 void Configuration::print_summary(){
1270     /* This function prints the result of the read config file
1271     It is highly recommended that you print this information
1272     and save as a log file. It prints alot of information
1273     but it is usefull when debugging and/or investigating
1274     strange simulation results.
1275     */
1276     cout<<"-----Simulation_summary-----"<<endl;
1277     //print the main simulation setup
1278     cout<<"##_Main_simulation_setup_##"<<endl;
1279     mandatory* m=get_mandatory();

```

```

1280 cout<<"The simulation will start out with having"<<m->users
1281 <<" users and"<<m->sites<<" sites."<<endl;
1282 cout<<"The simulation time is set to"<<m->time<<
1283 " using a mu value of"<<m->mu<<endl;
1284 cout<<" for the poisson distribution (approximately"
1285 <<m->time/m<<mu<<" job events)"<<endl;
1286 cout<<"The results will be saved to a subdirectory in"<<m->path
<<endl;
1287 cout<<"The random number seed is"<<m->seed<<endl;
1288 cout<<"—"<<endl;
1289 cout<<"The default properties for resources are:"<<endl;
1290 cout<<"Combination for direct and reputation trust using the"
1291 <<m->rDefaults.trust_eval<<" approach"<<endl;
1292 cout<<"and reputation calculation using the"<<m->rDefaults.
    rep_eval
<<" approach."<<endl;
1293
1294 if(m->rDefaults.trust_eval=="discrete")
1295     cout<<"The discrete ratio between direct and reputation trust
        is"
1296     <<m->rDefaults.discrete_ratio<<endl;
1297 if(m->rDefaults.rep_eval=="beta"){
1298     cout<<"The forgetting factor for the beta reputation is"
1299     <<m->rDefaults.forgetting<<endl;
1300     cout<<"The base weight for the beta reputation is"
1301     <<m->rDefaults.base_weight<<endl;
1302 }
1303 if(m->rDefaults.rep_eval=="discrete")
1304     cout<<"The rep ratio between old and new reputation trust is"
1305     <<m->rDefaults.rep_ratio<<endl;
1306
1307 cout<<"Behaviour [ i , j ]=["<<m->rDefaults.i<<" , "<<m->rDefaults.j<<
    ]"<<endl;
1308 cout<<"Groups=";
1309 copy(m->rDefaults.groups.begin(),
1310      m->rDefaults.groups.end(),
1311      ostream_iterator<int>(cout, " "));
1312 cout<<"(each number is a group)."<<endl;
1313 cout<<"The trust threshold is"<<m->rDefaults.threshold<<endl;
1314 cout<<"The direct weight (when calculating direct trust) is"
1315     <<m->rDefaults.direct_weight<<endl
1316     <<" using a smoothing factor of"<<m->rDefaults.smoothing<<
    endl
1317     <<" and sigma value"<<m->rDefaults.sigma<<endl;
1318 cout<<"The default trust value that is assigned to strangers is"
1319     <<m->rDefaults.stranger<<endl;
1320 cout<<"—"<<endl;
1321
1322 cout<<"##_Main_simulation_setup_end_##"<<endl<<endl;
1323
1324 //list changes details (listofchanges)
1325 cout<<"##_list_of_changes_##"<<endl;
1326 changelist::iterator cit;
1327 for(cit=listofchanges.begin();cit!=listofchanges.end();cit++){
1328     properties* p=get_properties((*cit)->id);

```

```
1329     cout<<(*cit)->id<<"_will_at_time="<<(*cit)->time<<"_change_[i,j
        ]_from_"
1330         <<p->i<<"_"<<p->j<<"_]_to_"
1331         <<(*cit)->i<<"_"<<(*cit)->j<<""]<<endl;
1332     print_difference((*cit)->id,p);
1333 }
1334 cout<<endl;
1335
1336 //print additions details (listofadds)
1337 cout<<"##_list_of_additions.##_"<<endl;
1338 addlist::iterator ait;
1339 for(ait=listofadds.begin();ait!=listofadds.end();ait++){
1340     cout<<(*ait)->id<<"_will_be_added_at_time="<<(*ait)->time<<endl
        ;
1341     properties* p=get_properties((*ait)->id);
1342     print_difference((*ait)->id,p);
1343 }
1344 cout<<endl;
1345
1346 //print saving options
1347 sOptions->print_save_info();
1348 };
```



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