

Fine-Grained Bandwidth Adaptivity in Networks-on-Chip Using Bidirectional Channels

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University of Toronto

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Motivation

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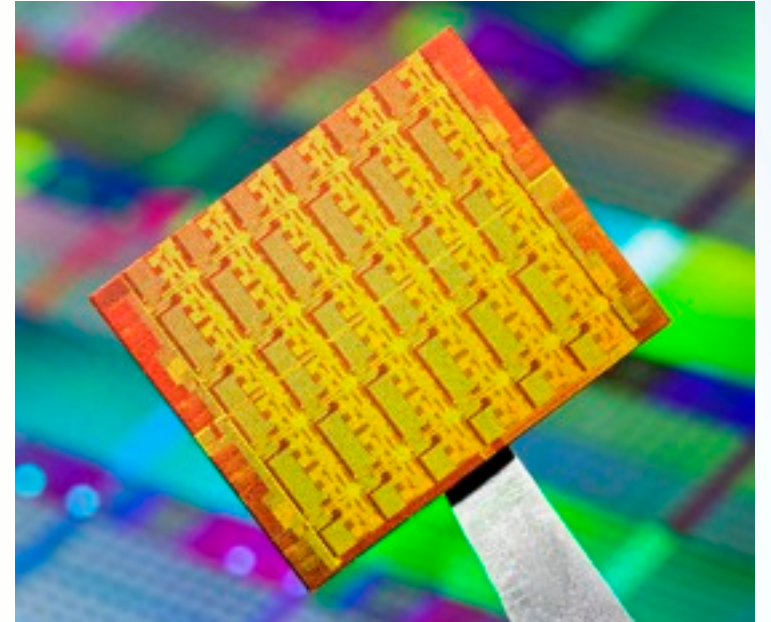
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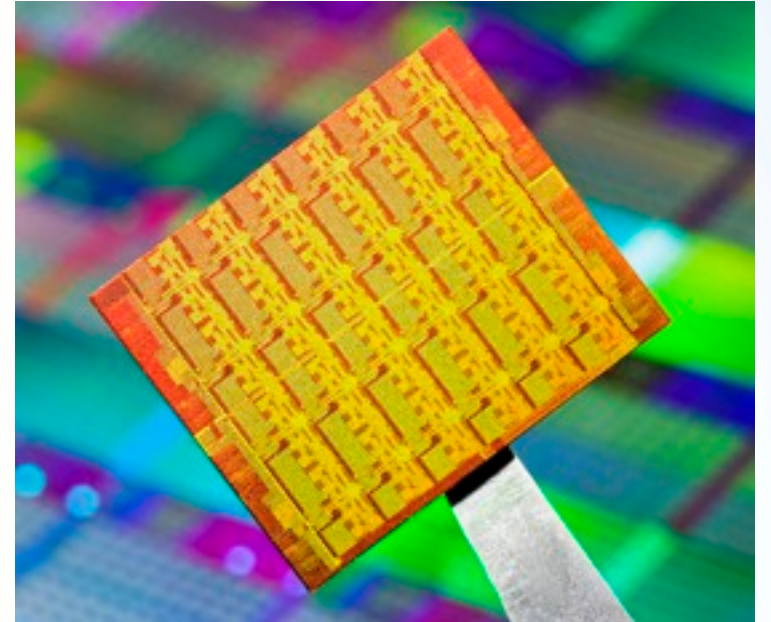
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- NoCs are crucial for scaling CMPs



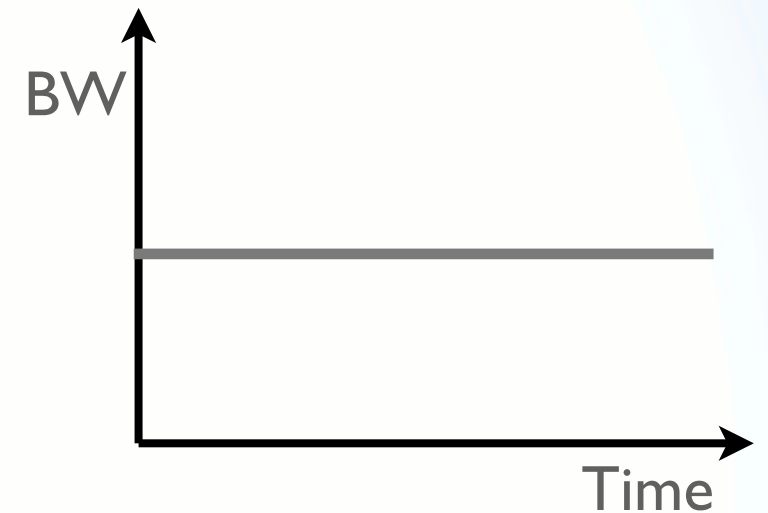
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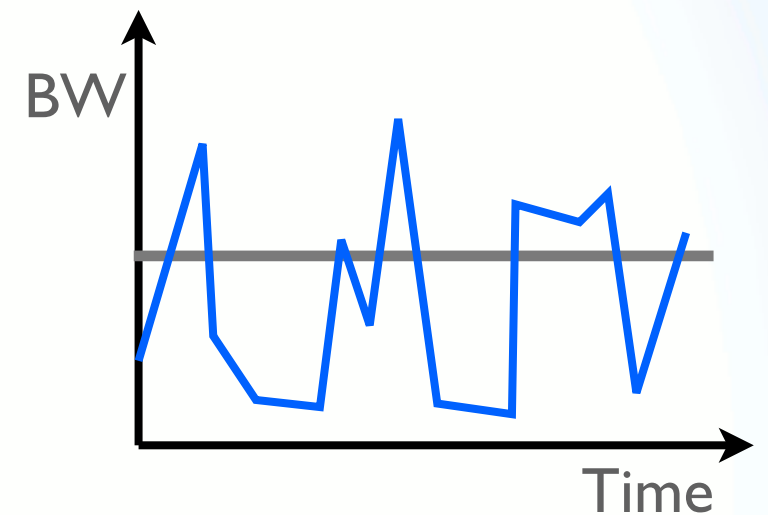
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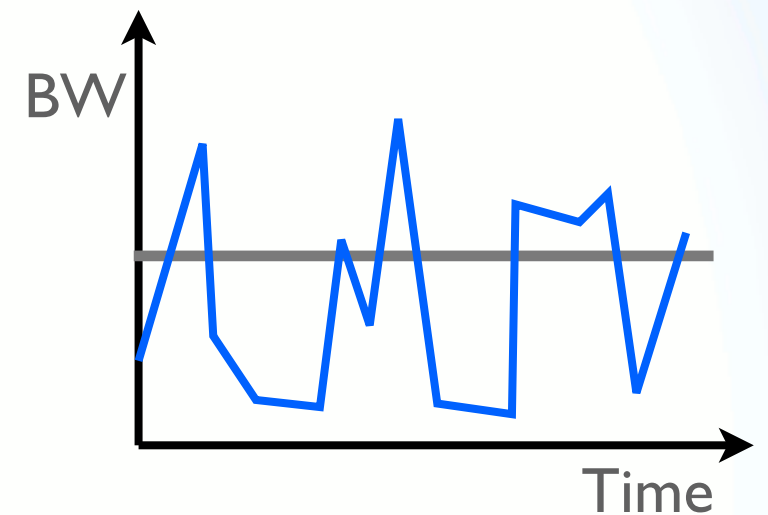
Motivation

- NoCs are crucial for scaling CMPs
- Problem:
 - NoC bandwidth resources are static
 - Bandwidth requirements are highly dynamic



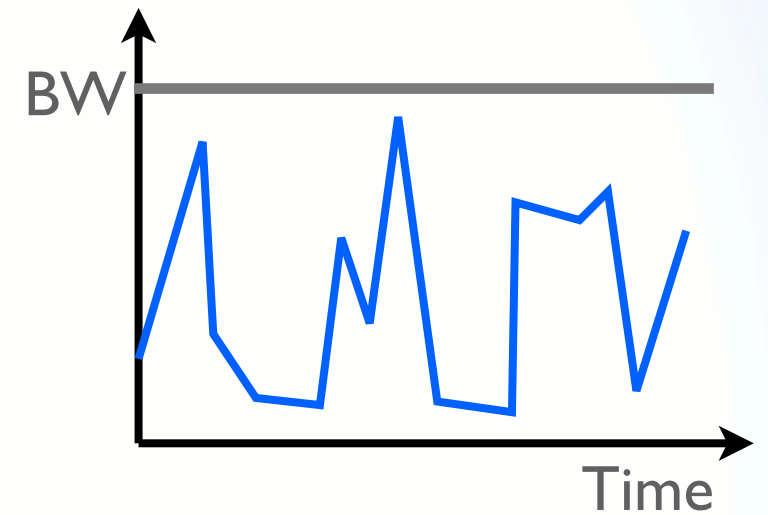
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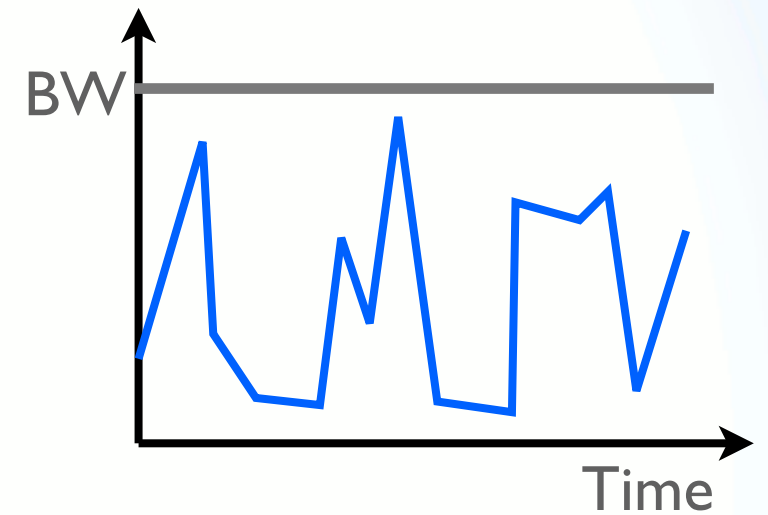
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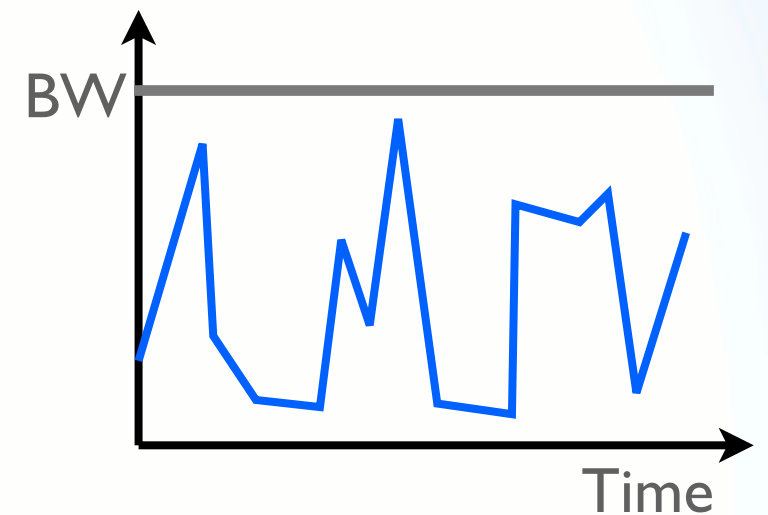


Average channel utilization: < 5%



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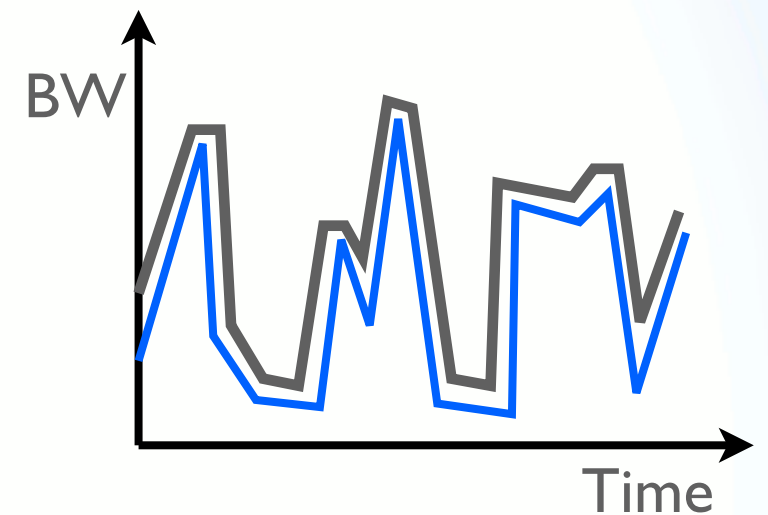


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 - Adapt link BW to demands

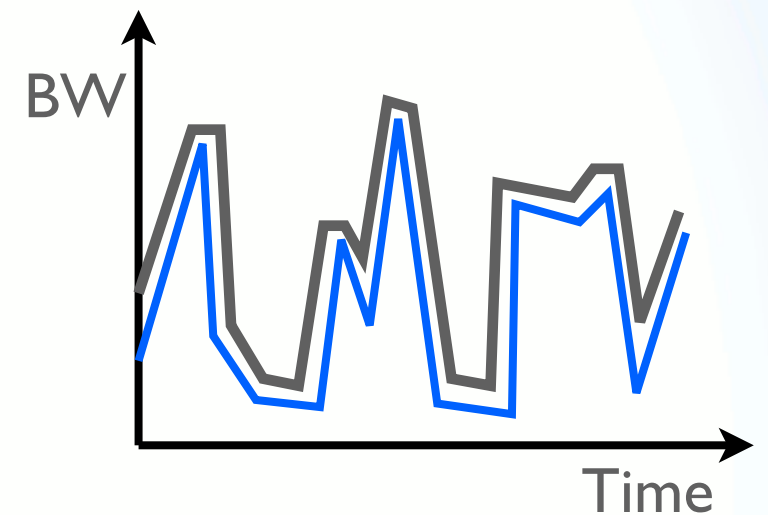


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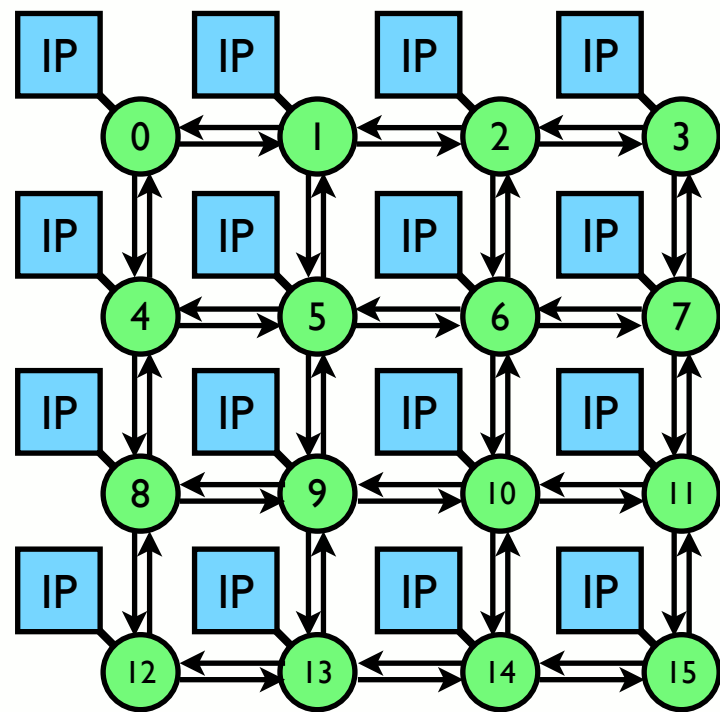
Save up to 75% of BW resources



Motivation – Static NoC

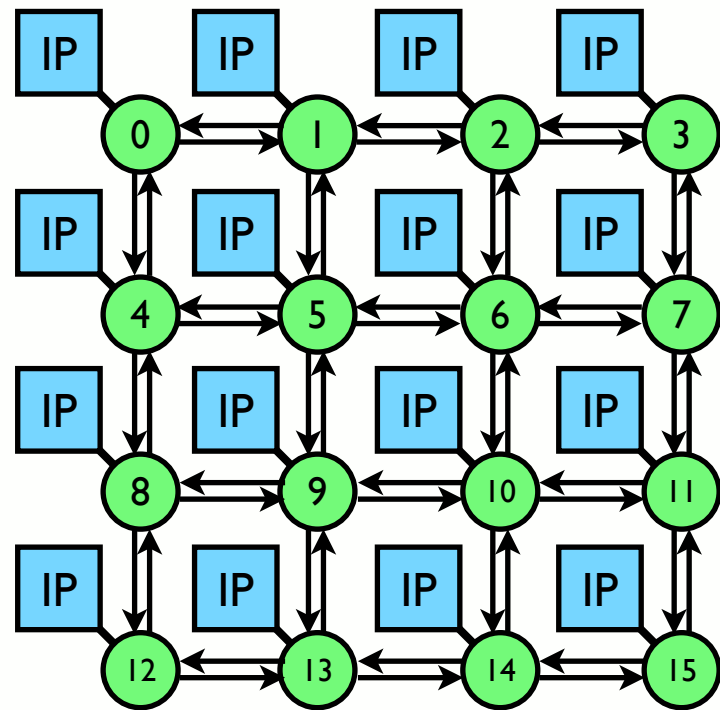


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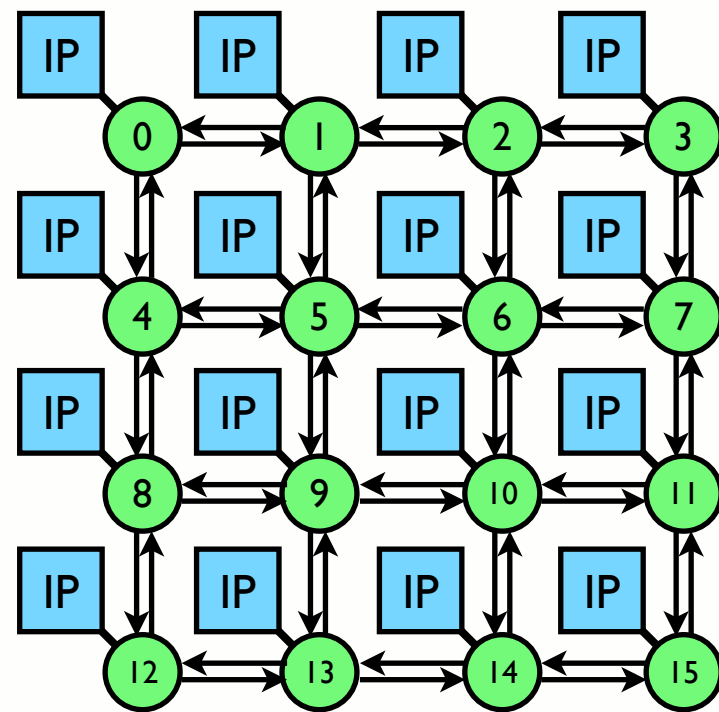
- Static Topology

Motivation – Static NoC

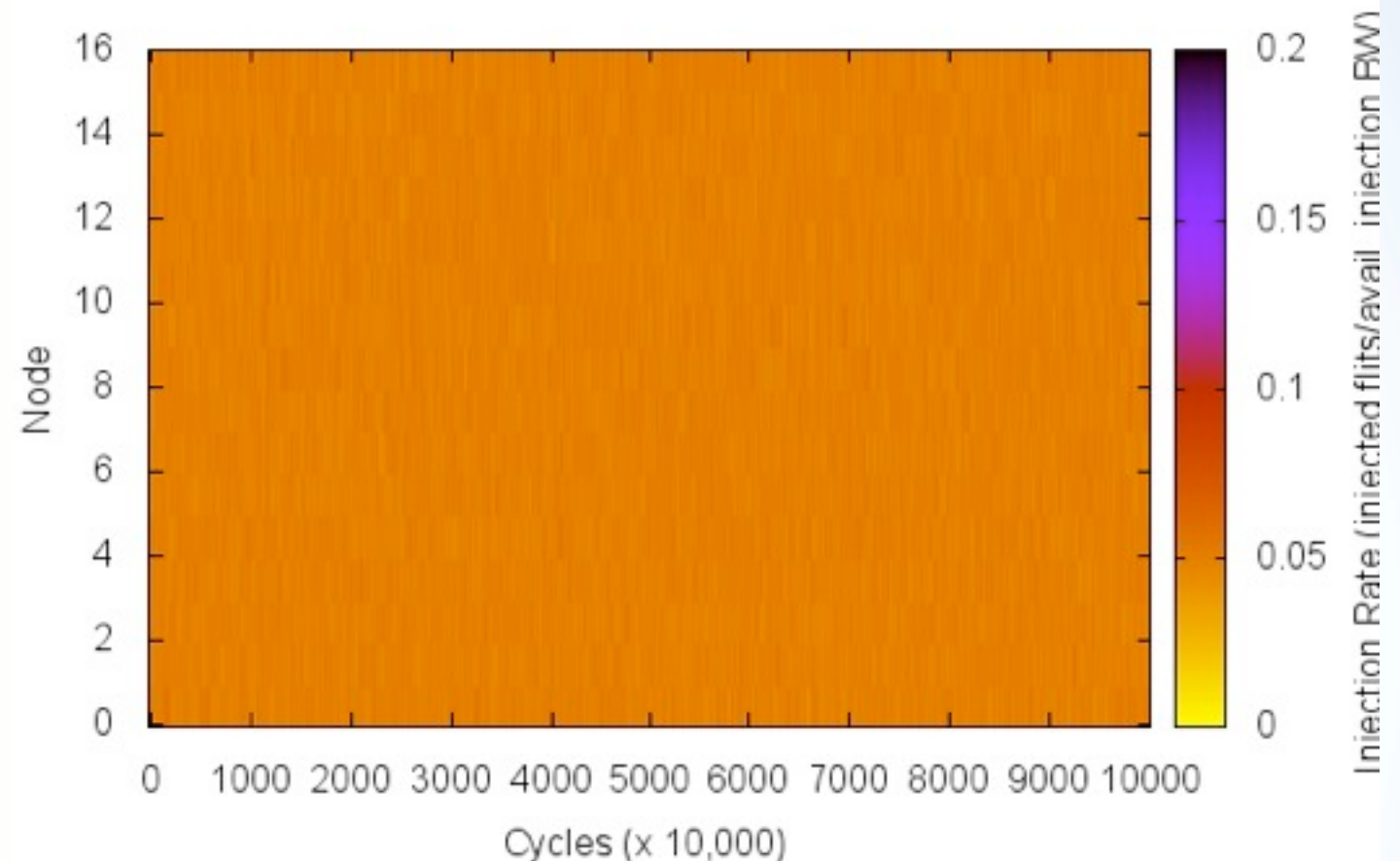


- Static Topology
- Static Bandwidth

Motivation – Static NoC



Uniform Random

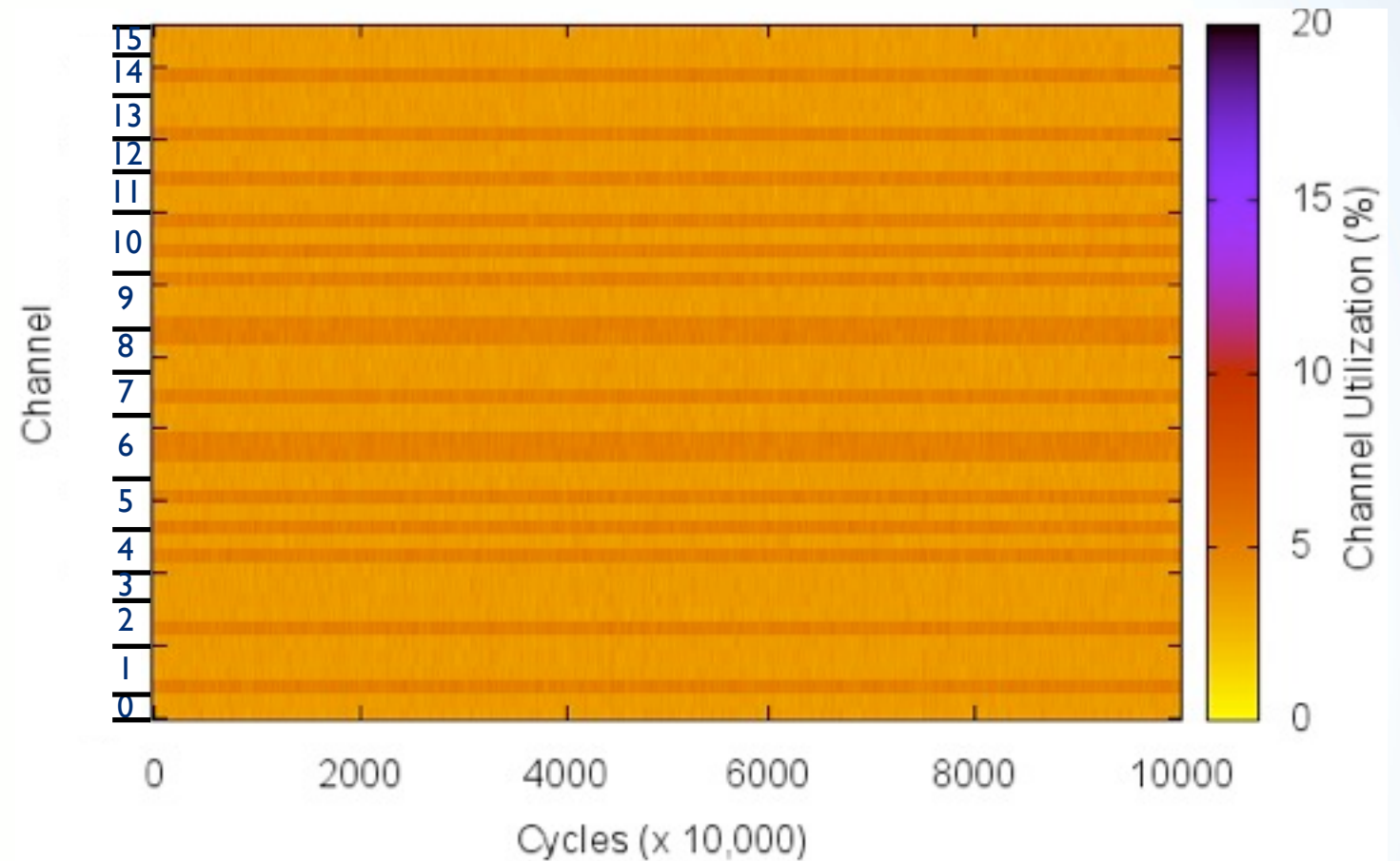
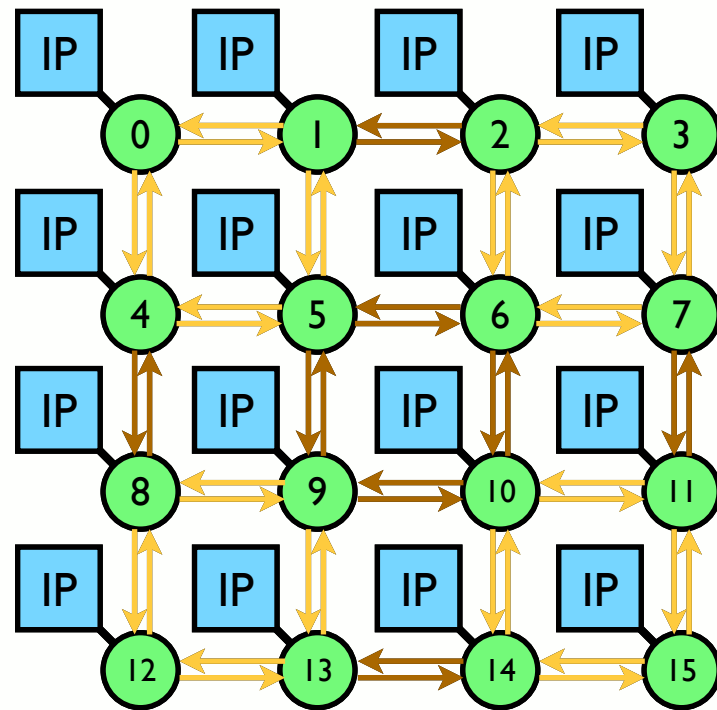


- Static Topology
- Static Bandwidth
- Static workloads for evaluation



Motivation – Static NoC

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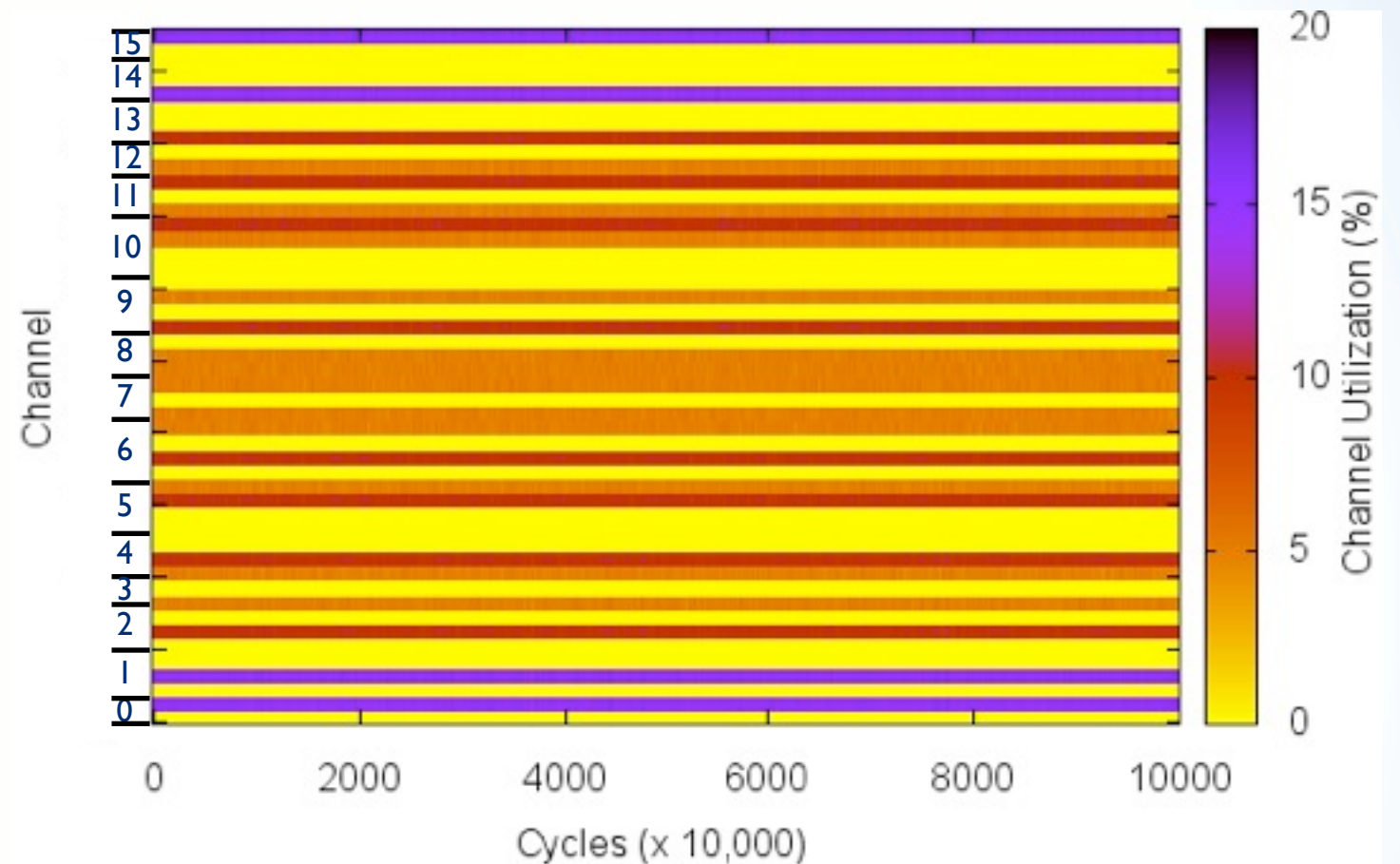
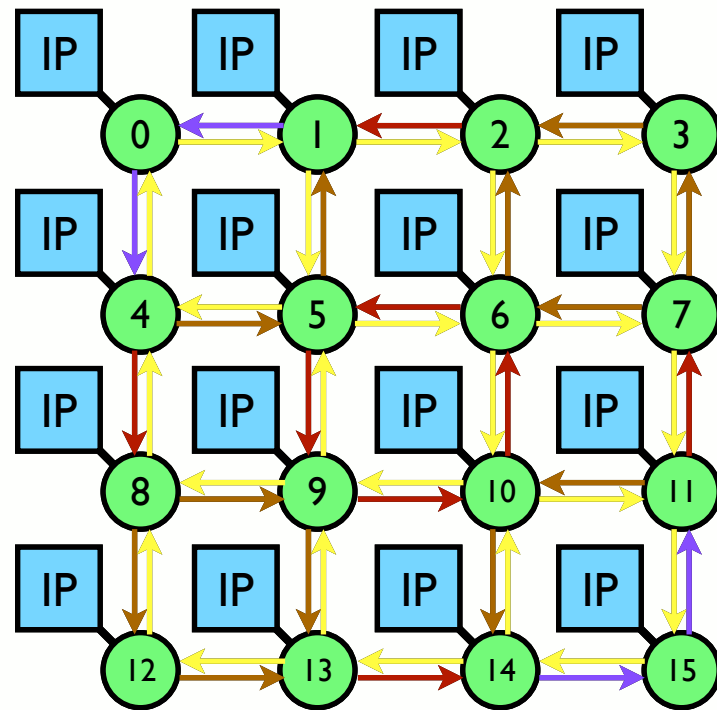


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Motivation – Static NoC

Transpose

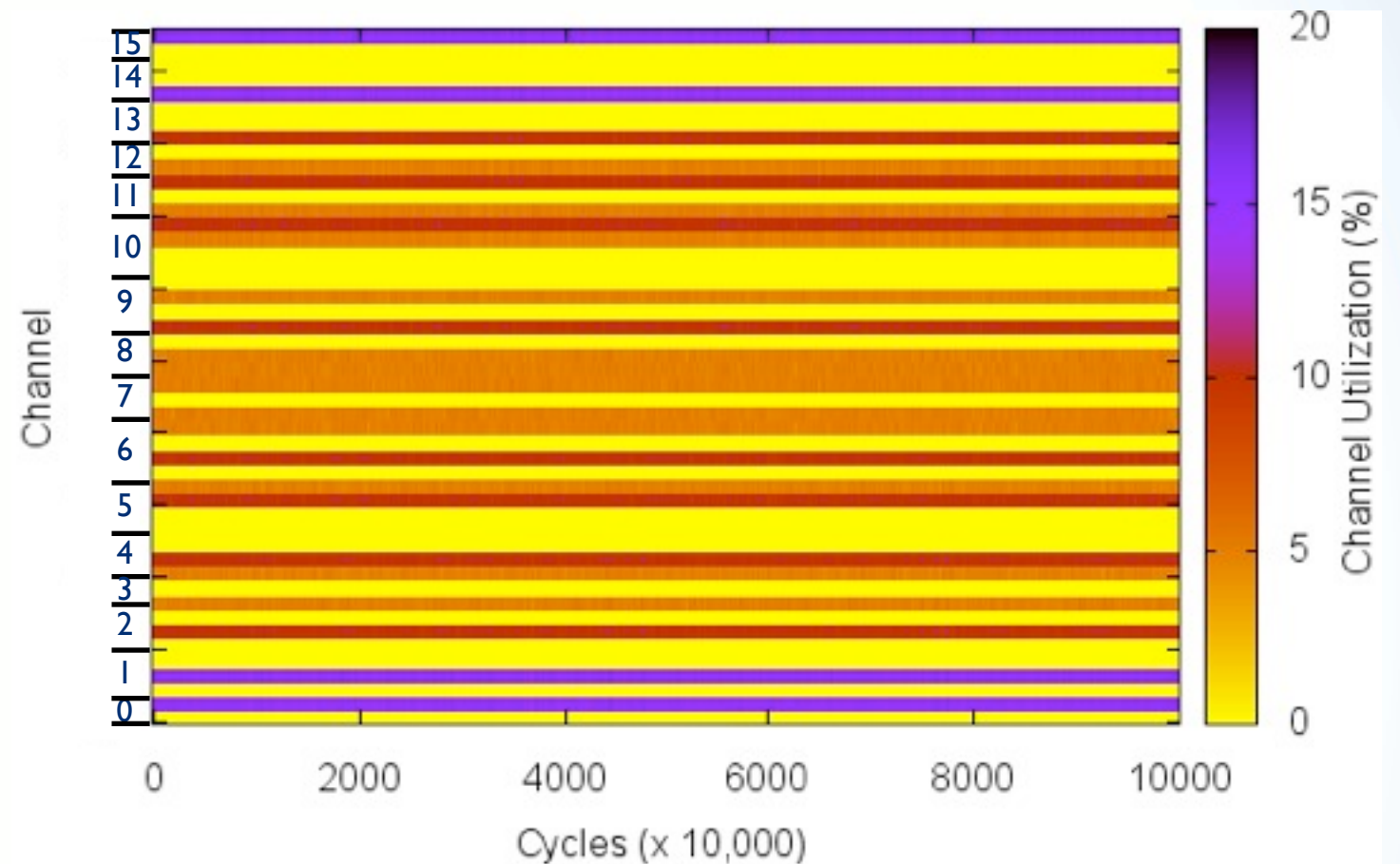
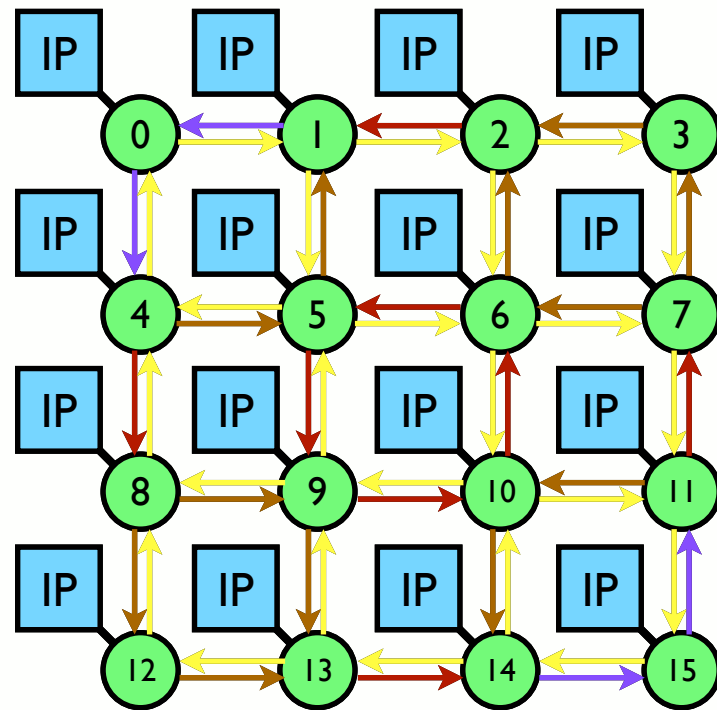


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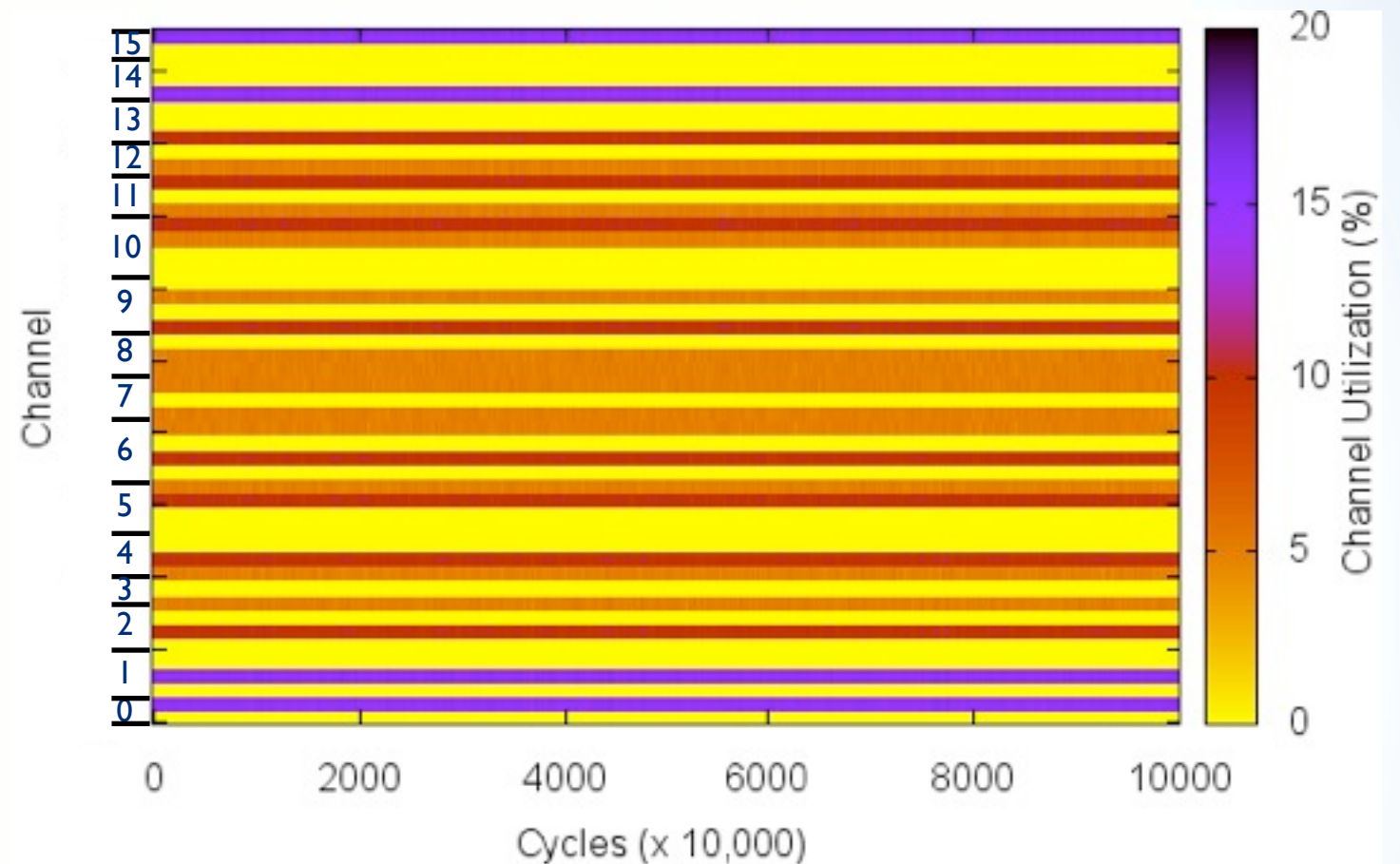
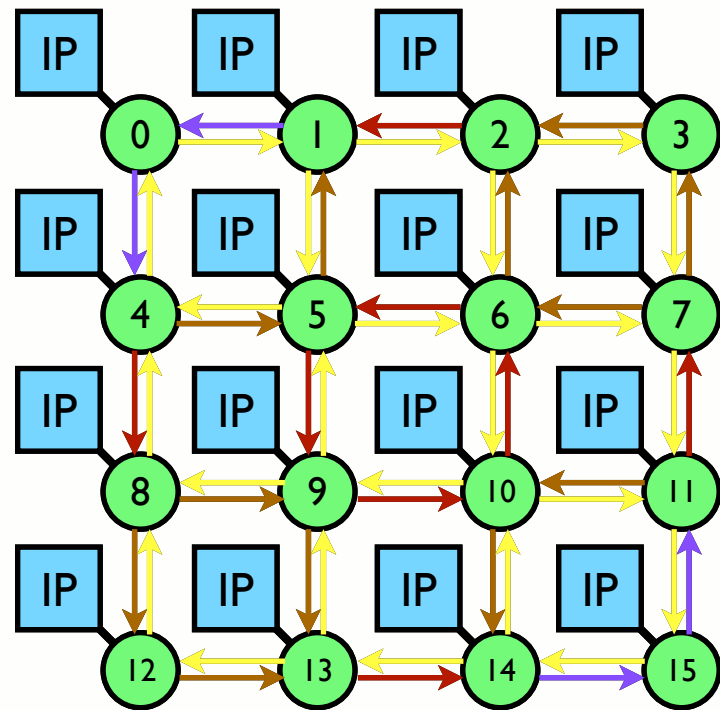


- Static Topology
- Static Bandwidth
- Static workloads for evaluation
- Specified at design time for worst case scenario



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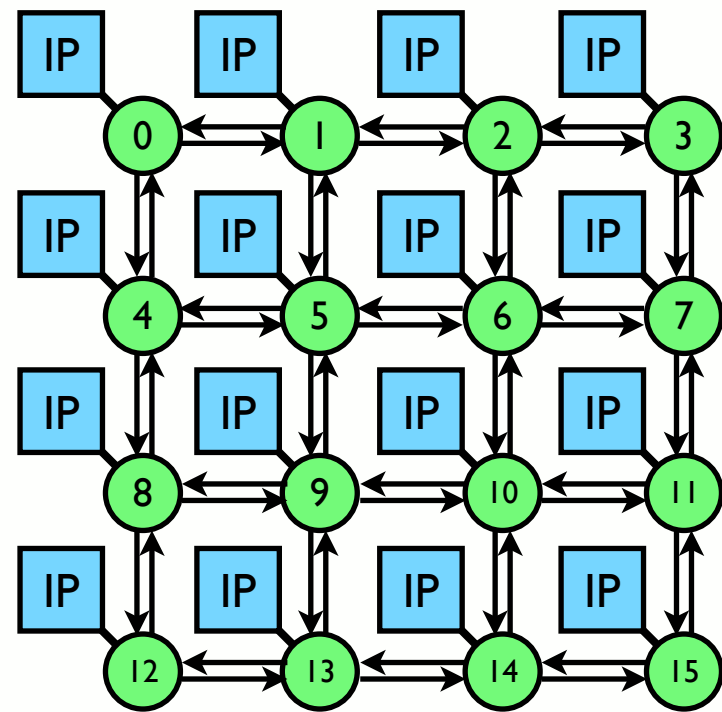
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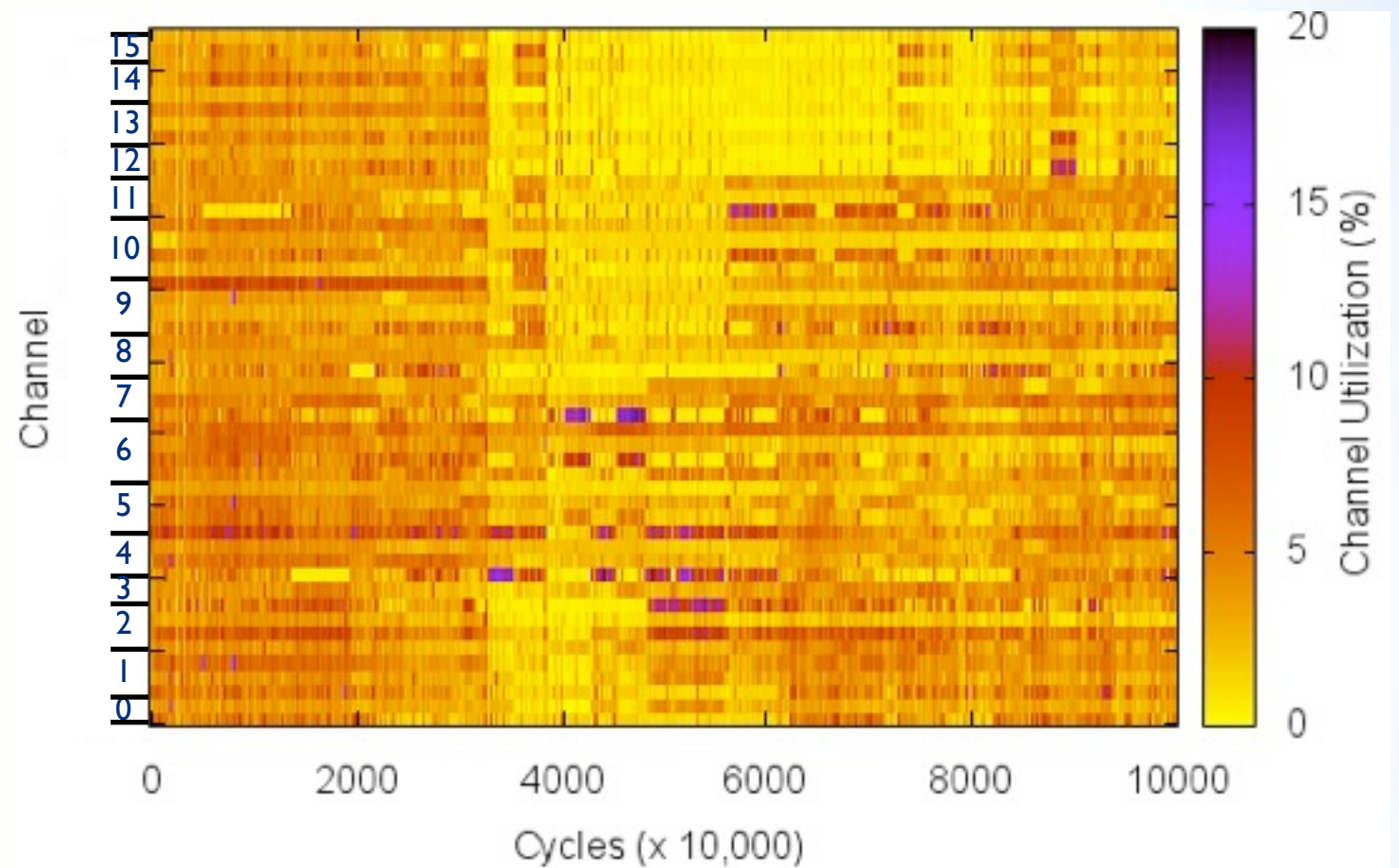
- Static Topology
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- Static workloads for evaluation
- Specified at design time for worst case scenario
- Static NoCs can handle temporally- and spatially-stable traffic well



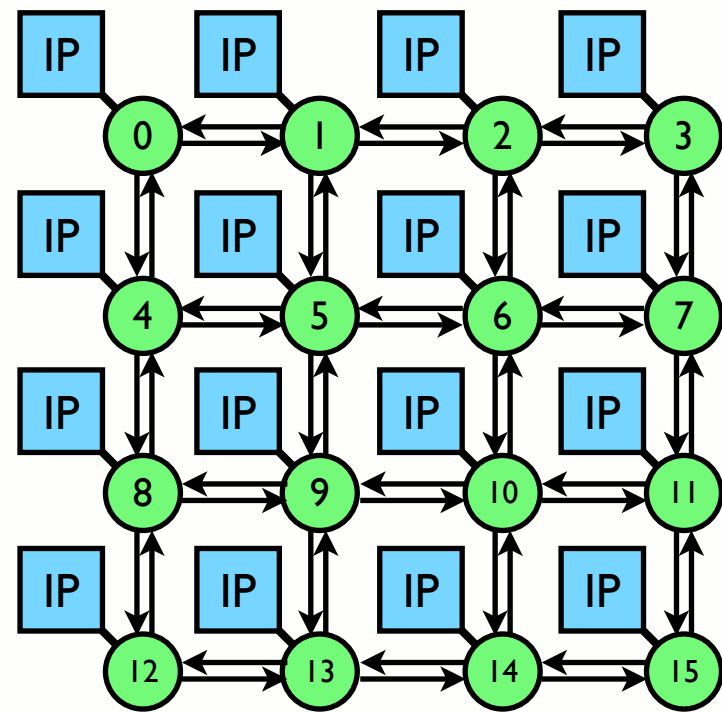
Motivation – Real NoC Traffic



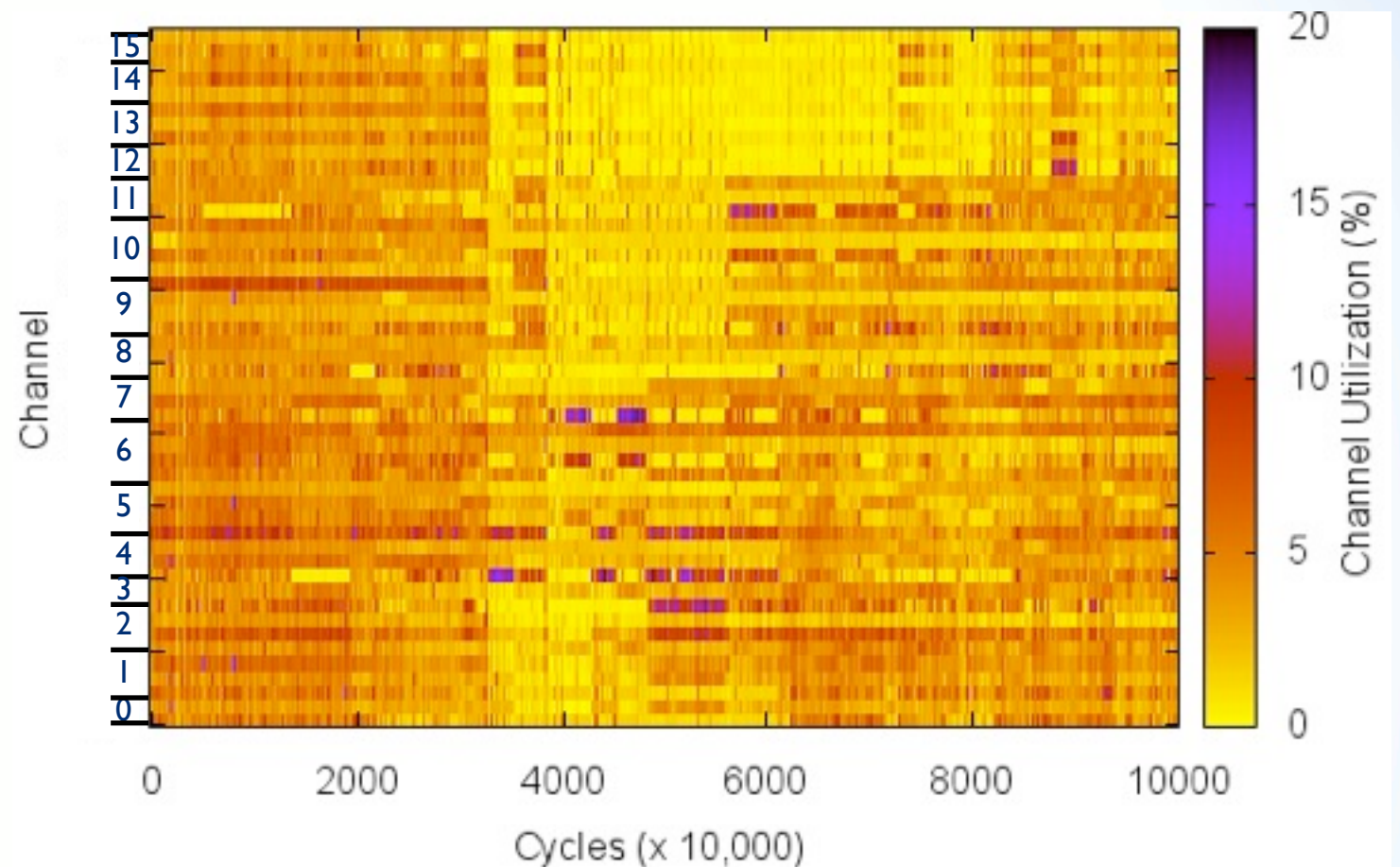
Blackscholes



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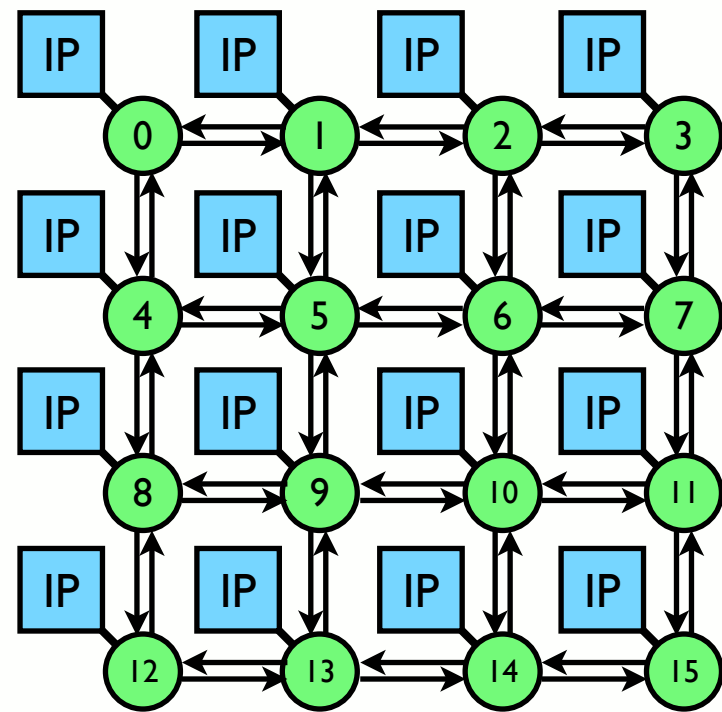


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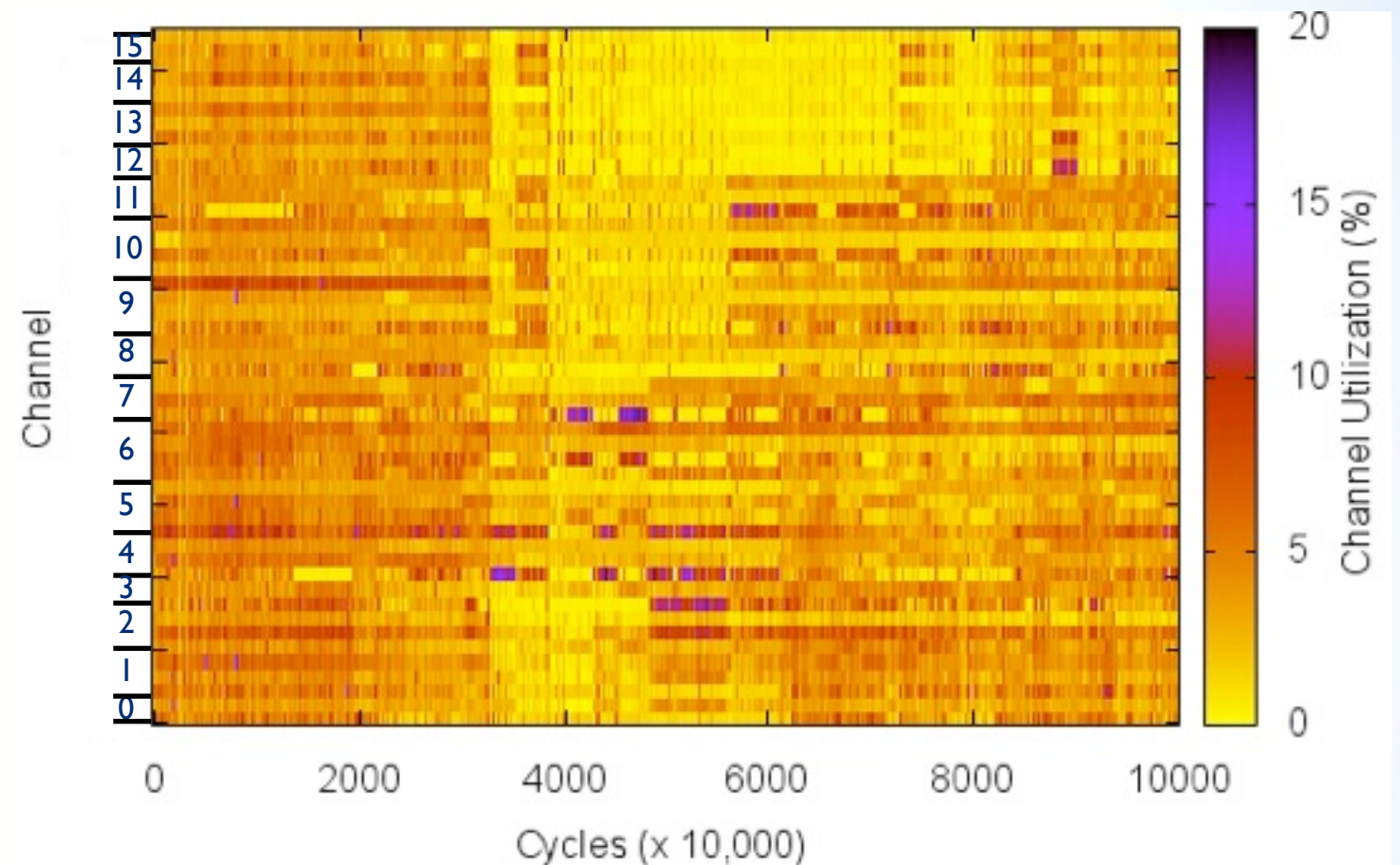


- Highly dynamic workloads

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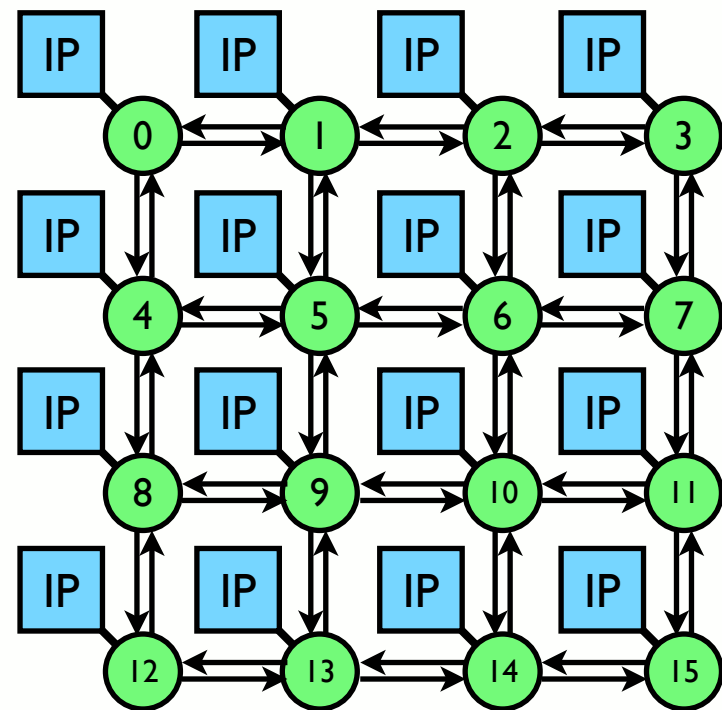


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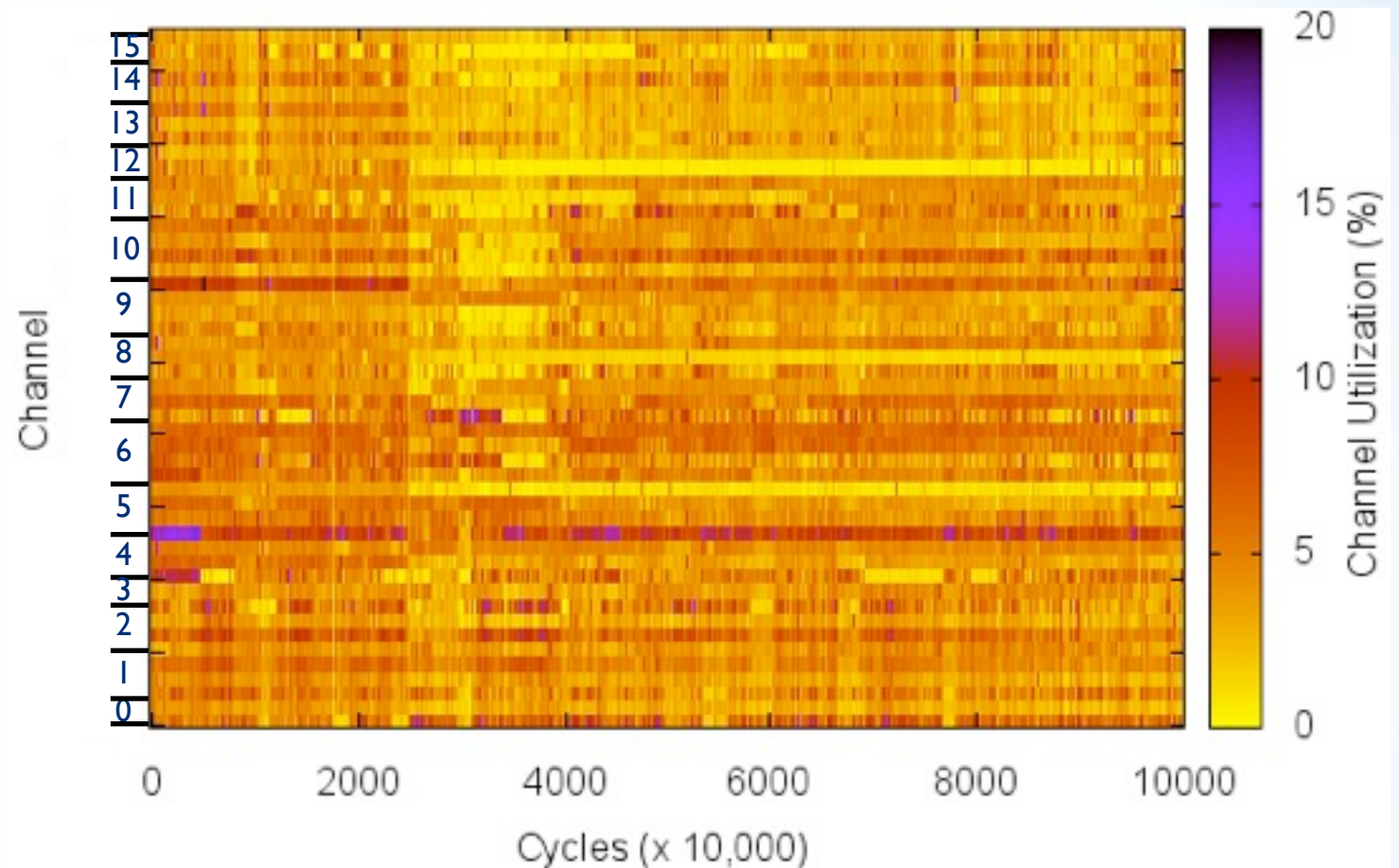


- Highly dynamic workloads
- Large temporal and spatial BW variance

Motivation – Real NoC Traffic



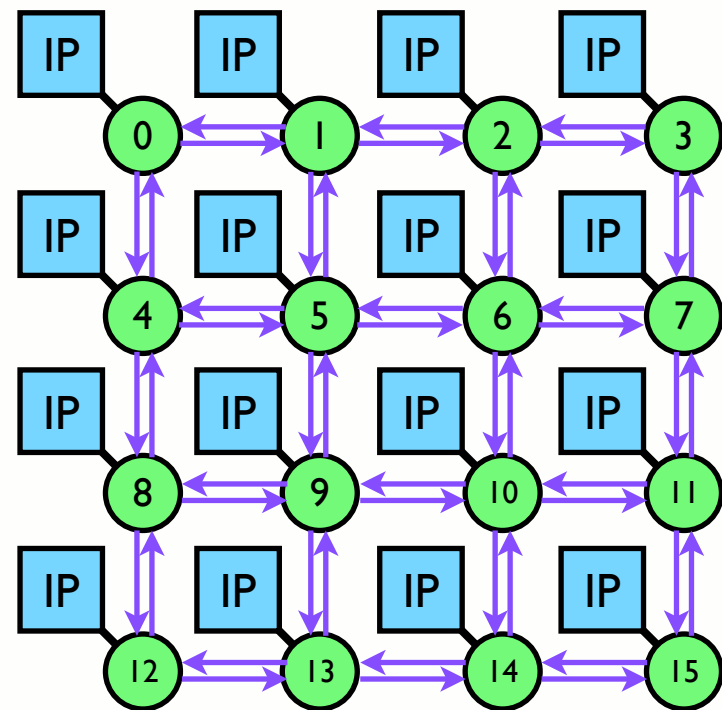
Streamcluster



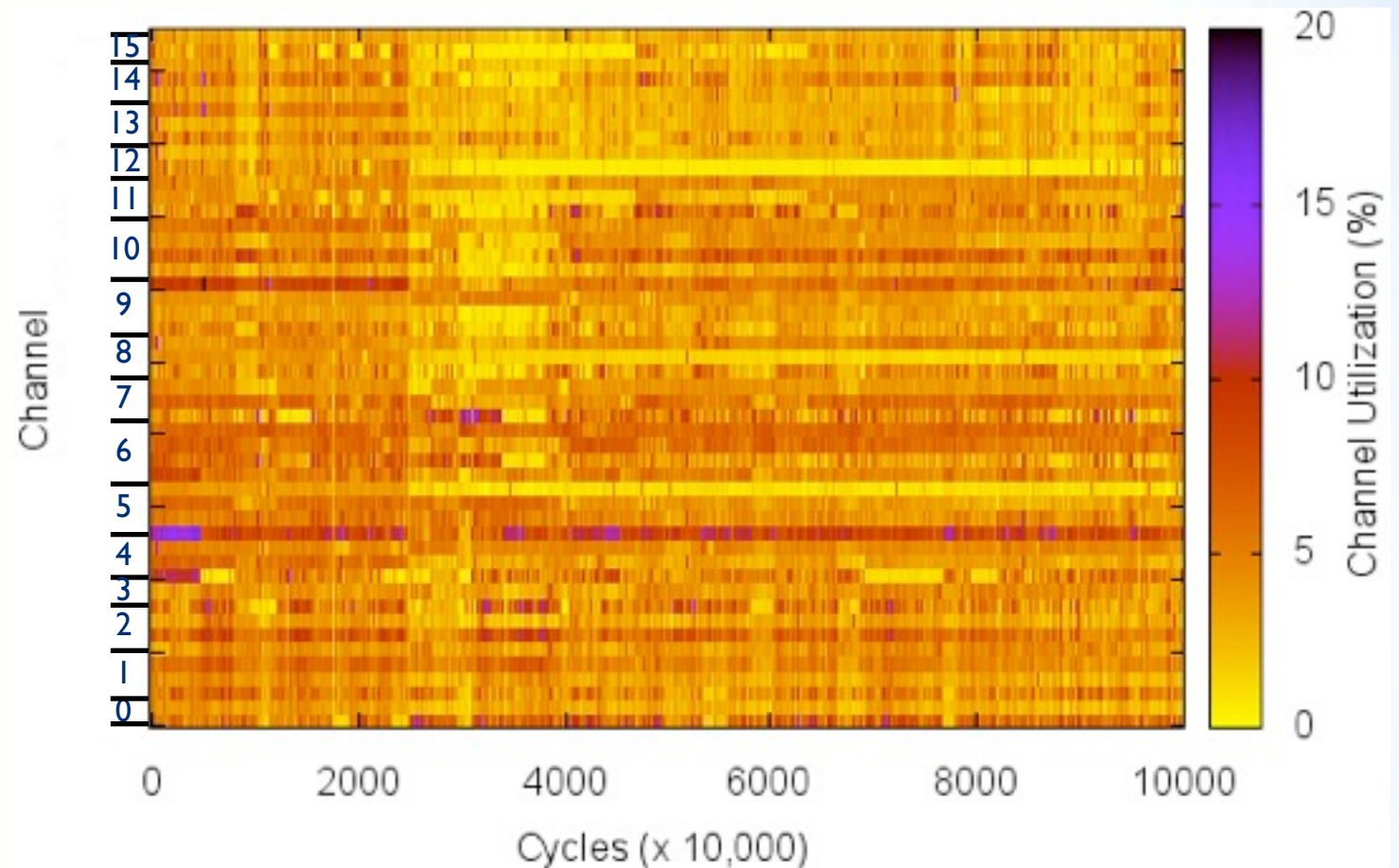
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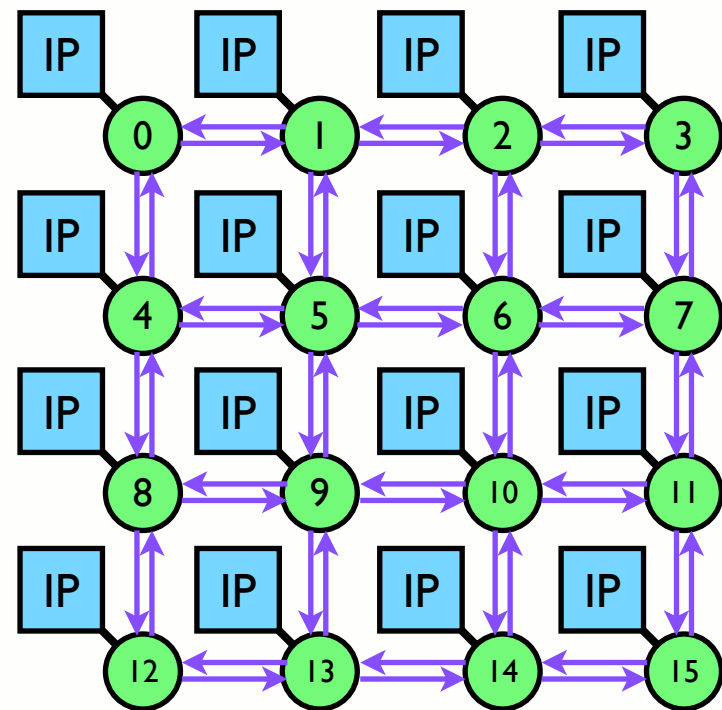


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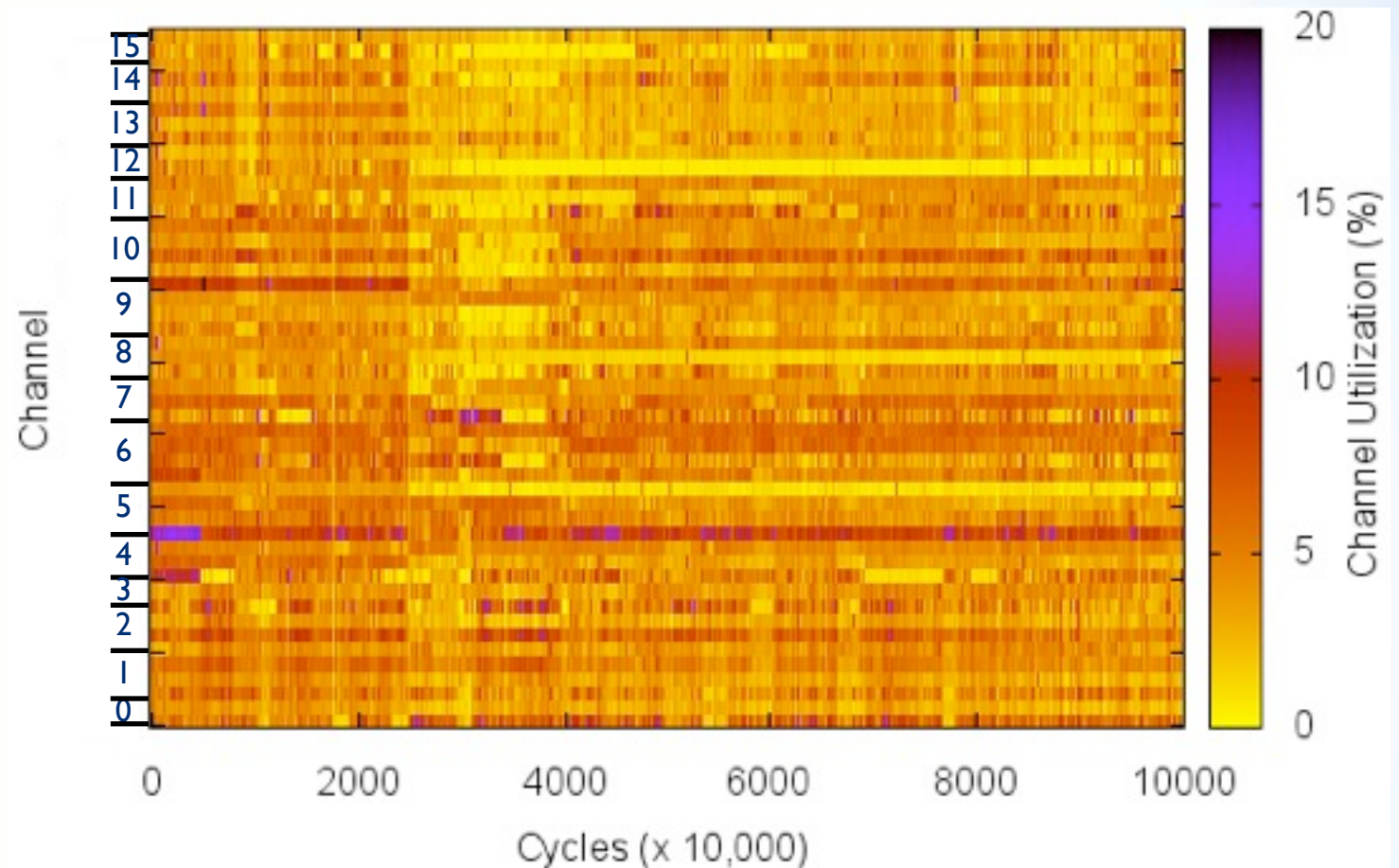
- Significant area and power overhead with traditional NoC implementation



Motivation – Real NoC Traffic



Streamcluster



- Highly dynamic workloads
- Large temporal and spatial BW variance

- Significant area and power overhead with traditional NoC implementation
- Channels are underutilized most of the time



Channel Utilization

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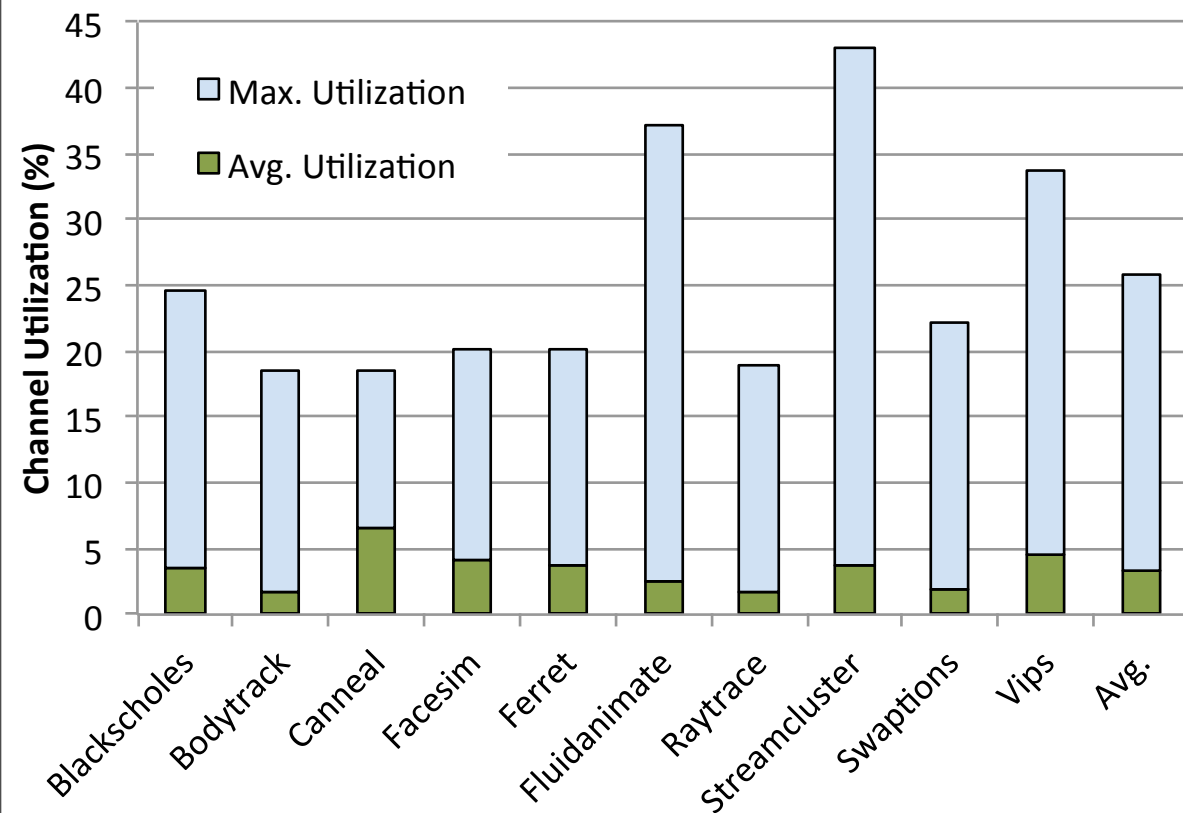
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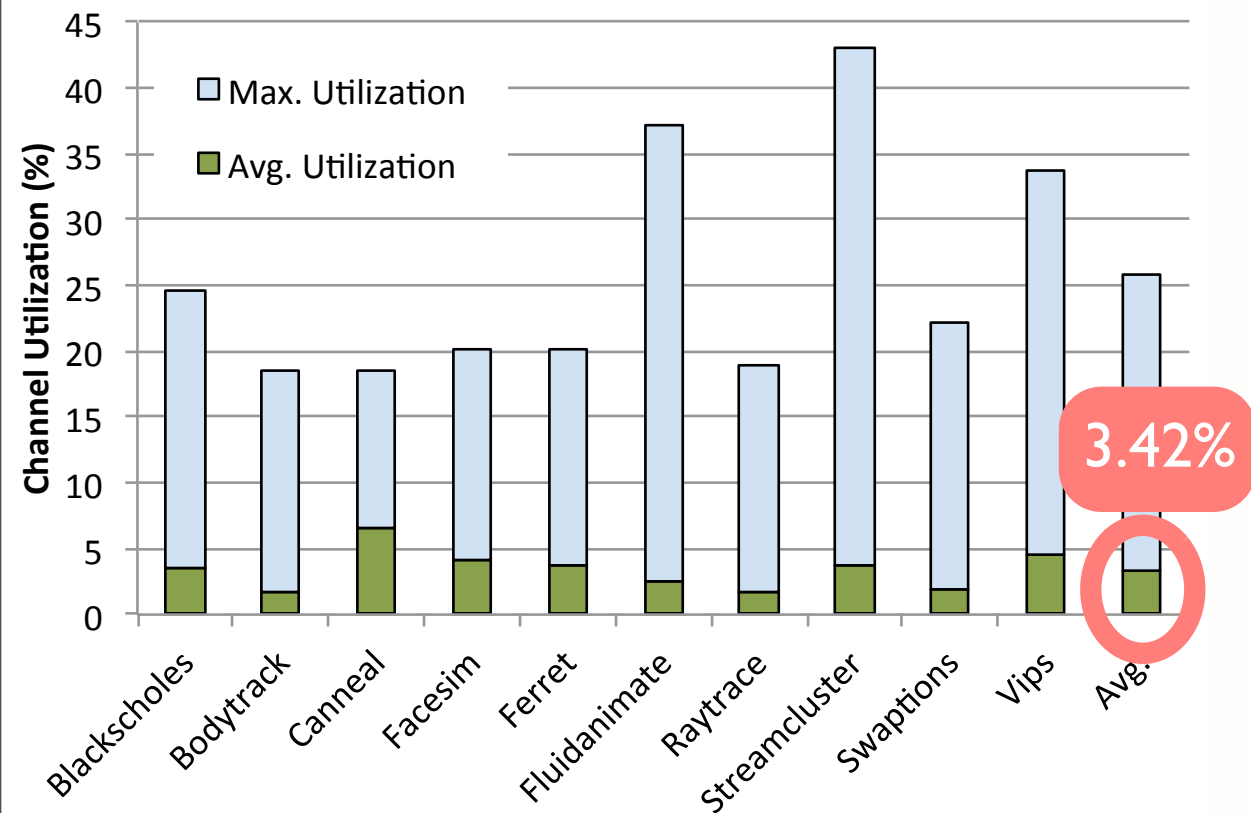


Friday, 11 May, 12

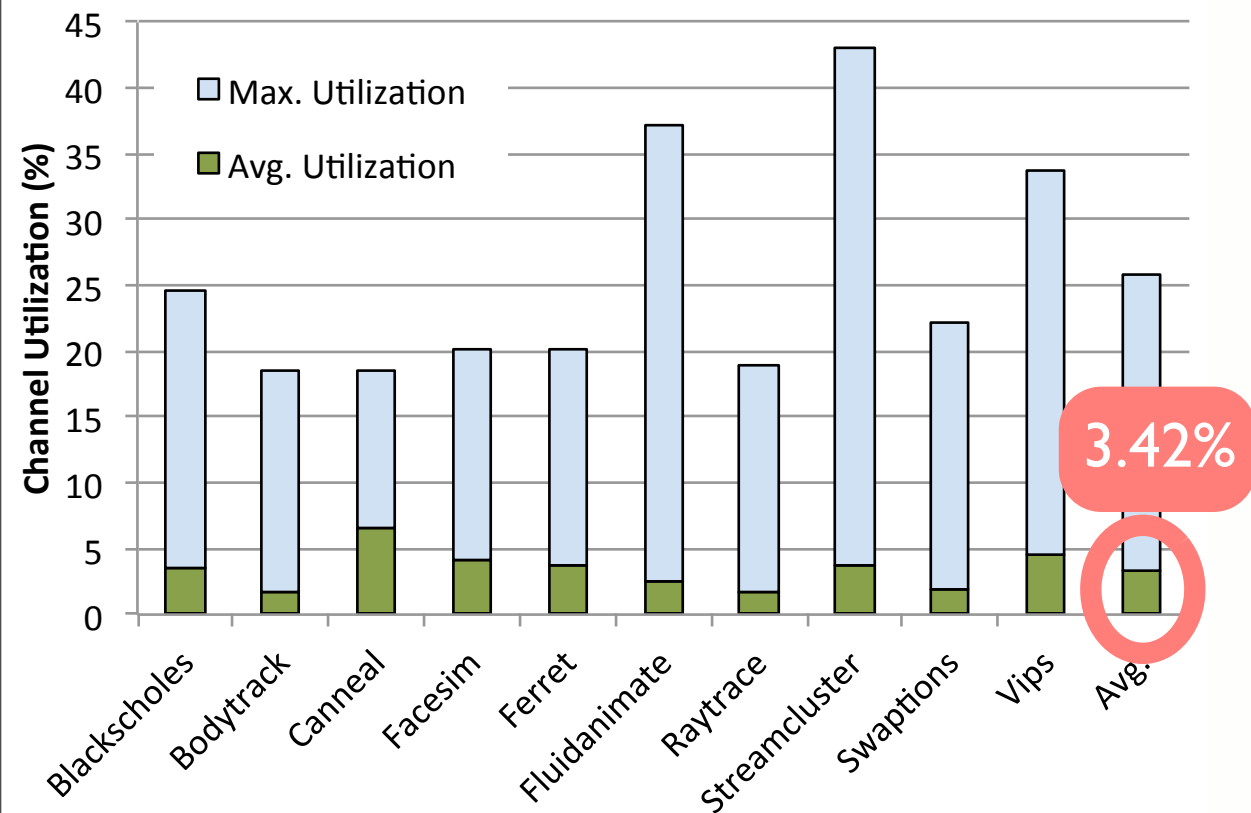
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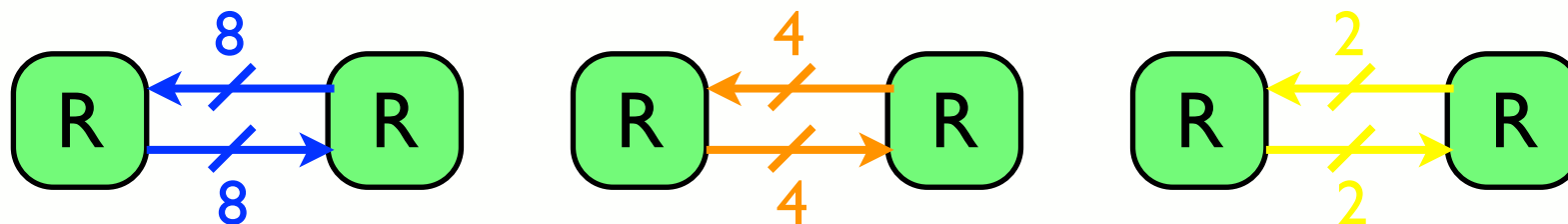
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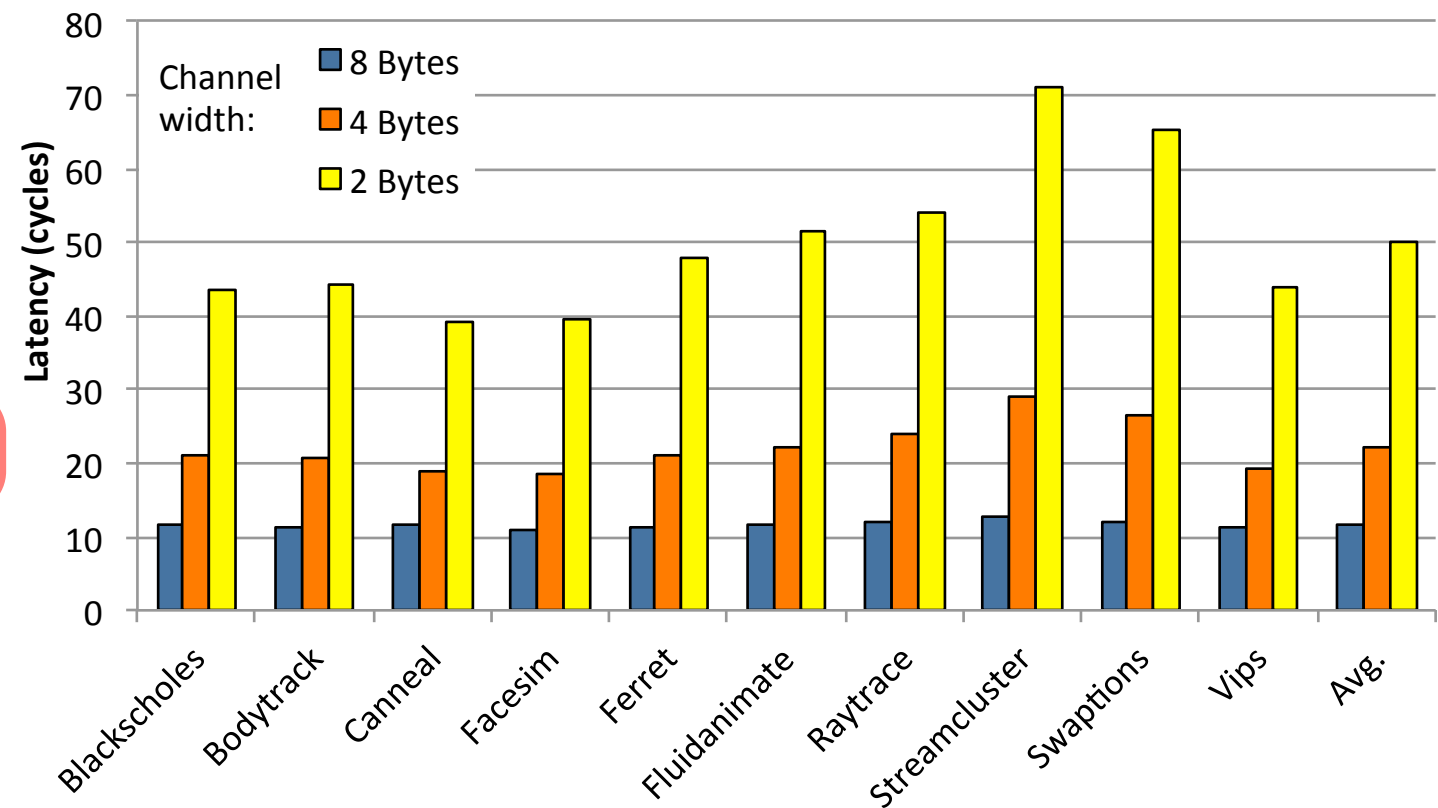
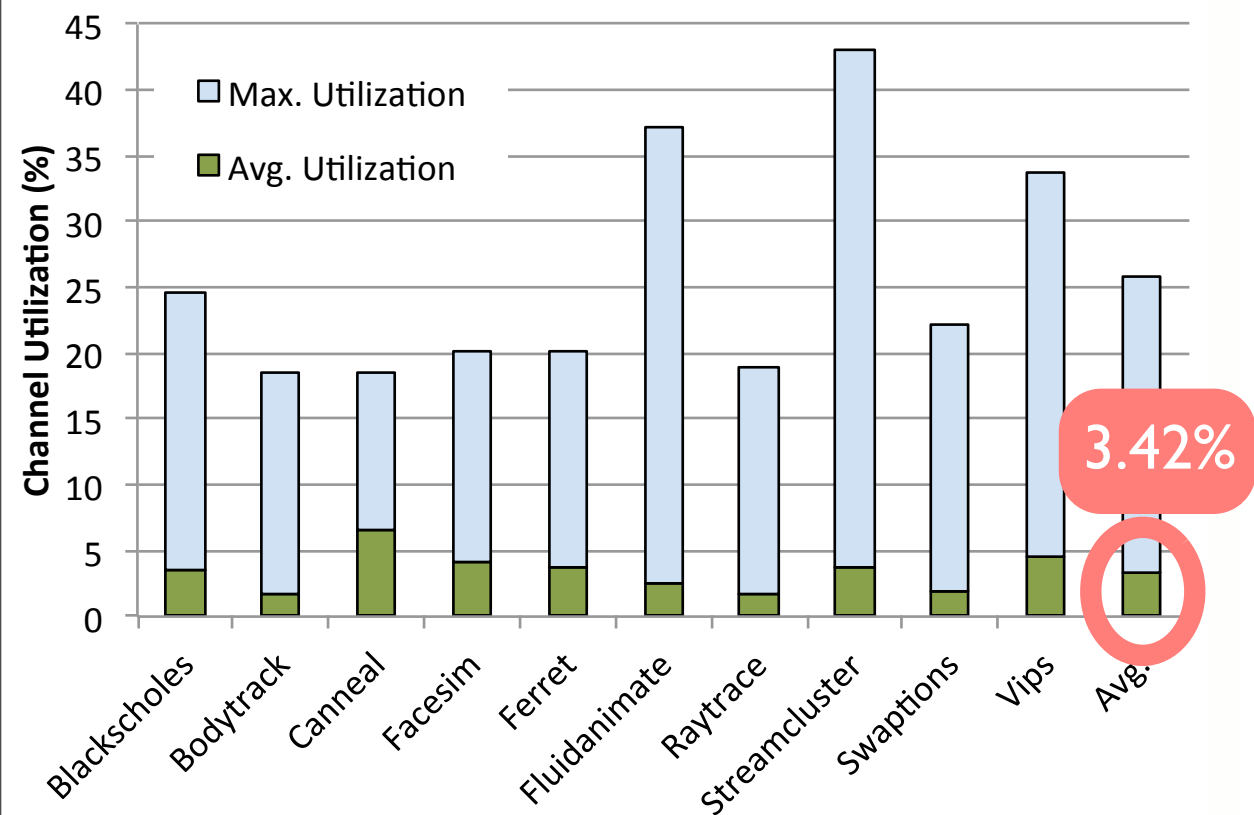
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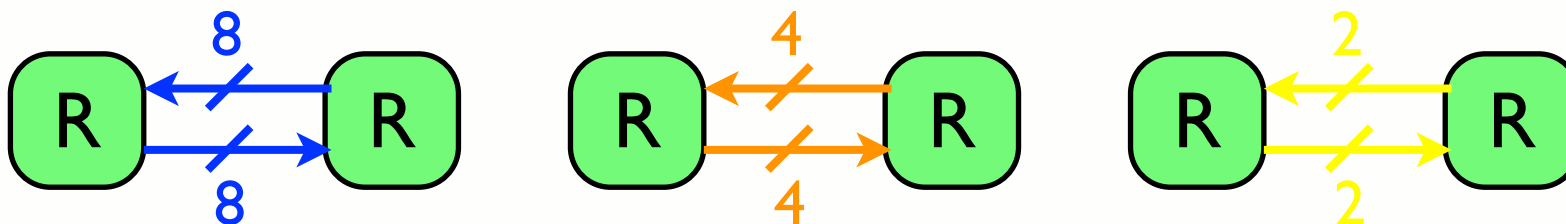
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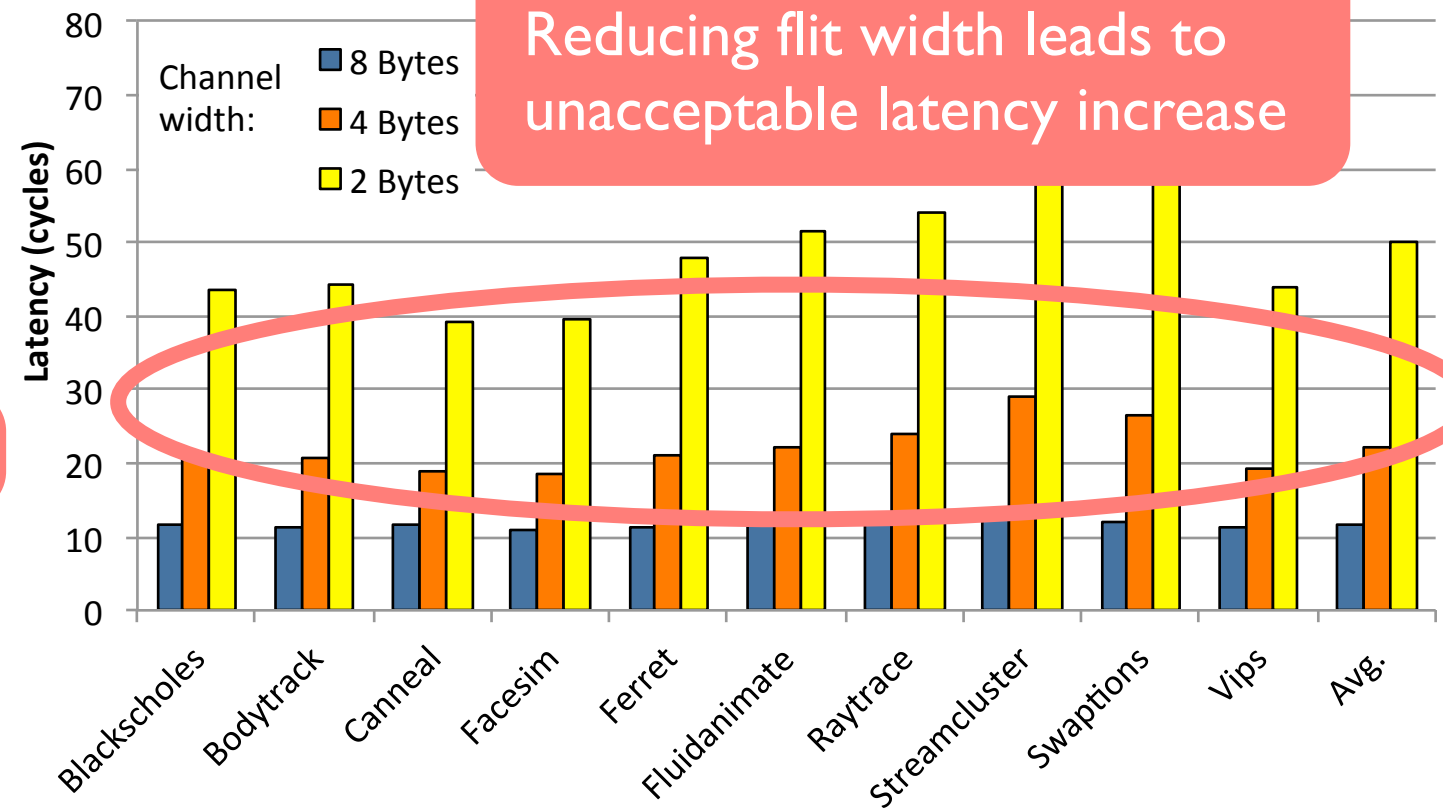
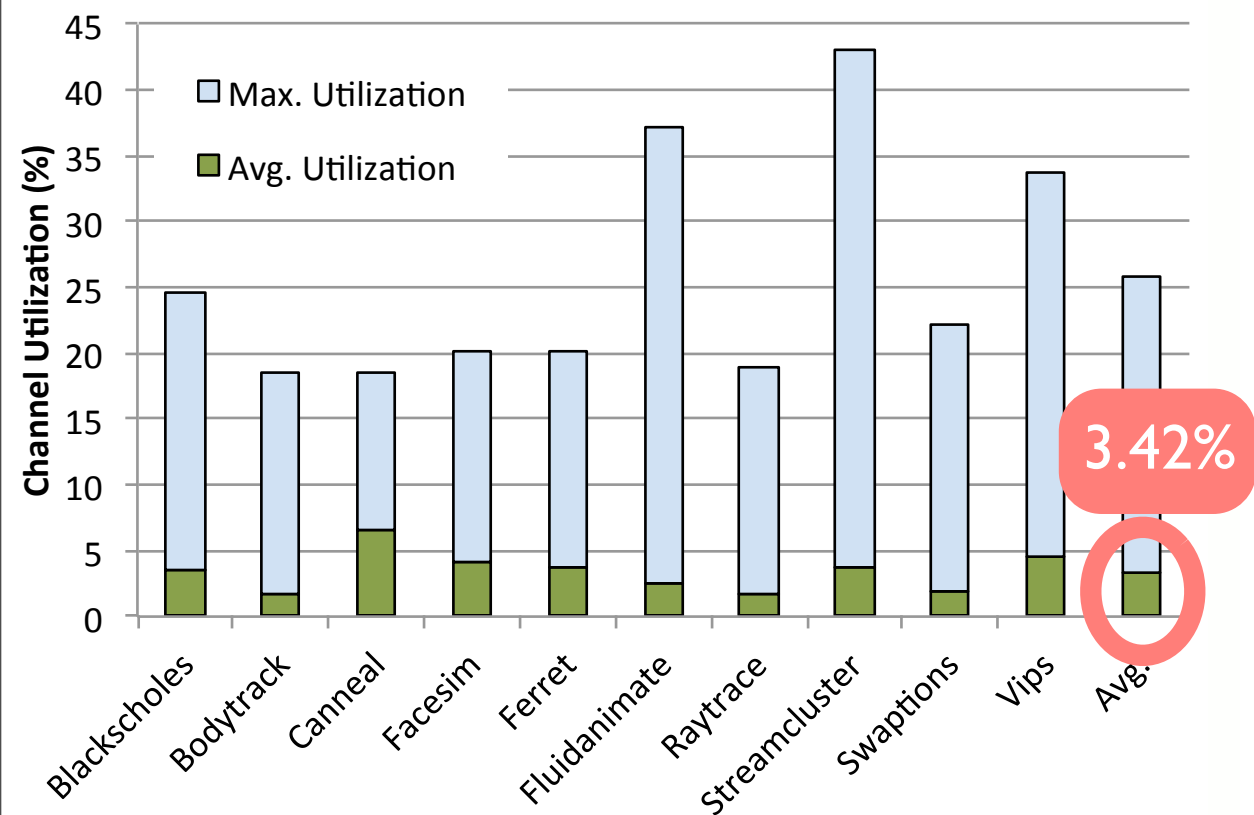
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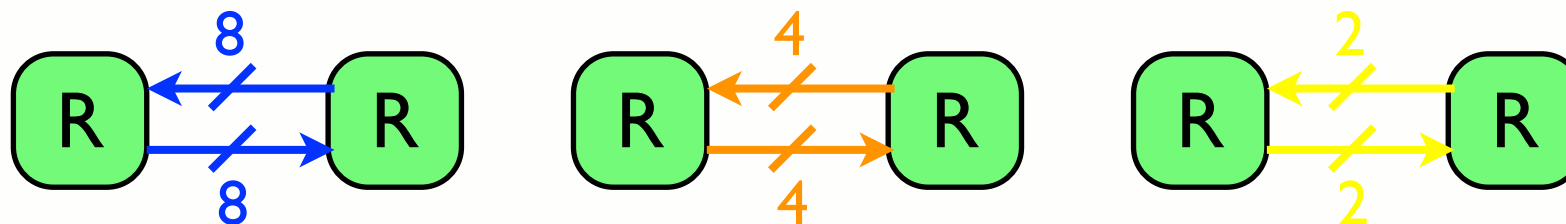
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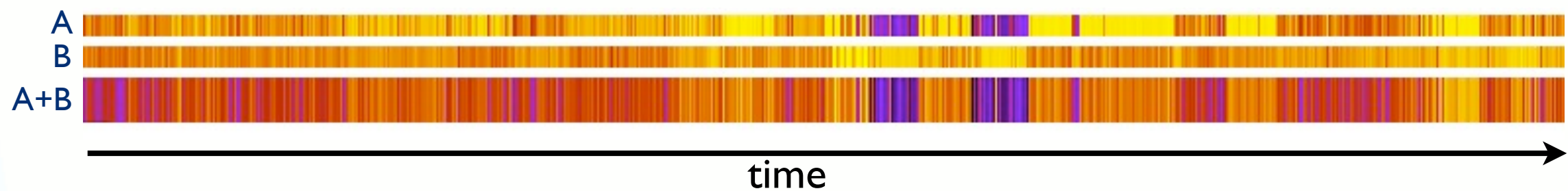
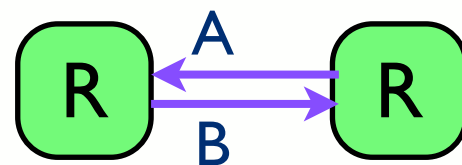


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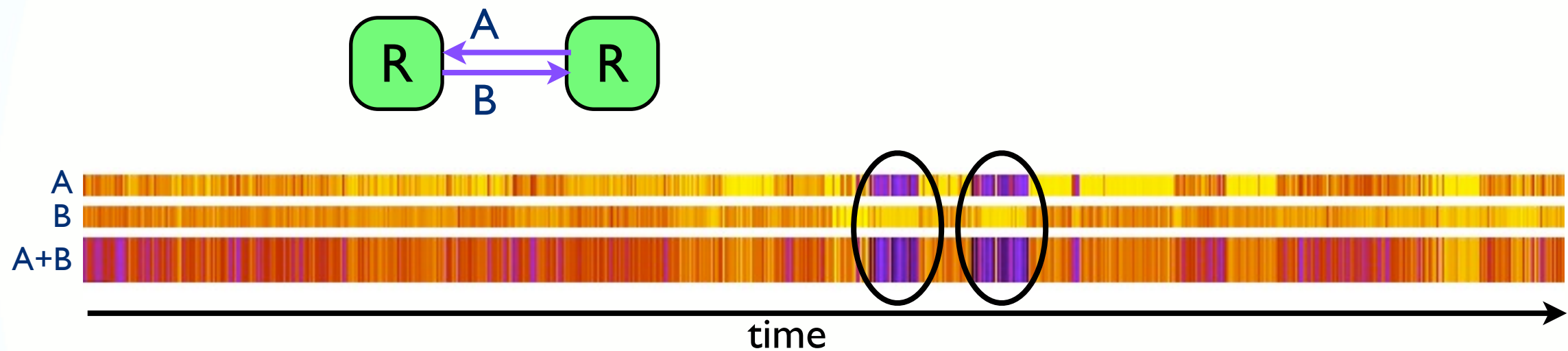
Bidirectional Channels

- Bidirectional channels to share channel resources:



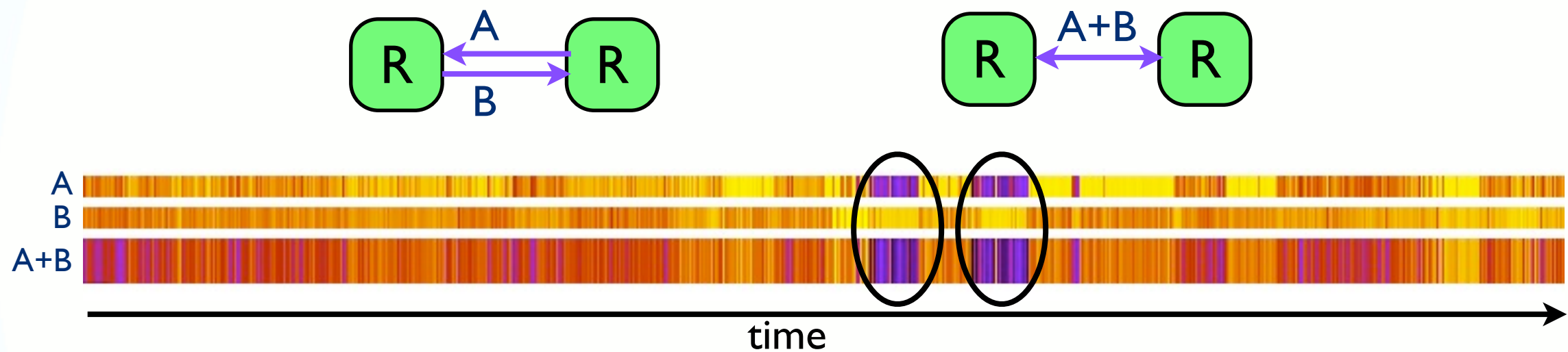
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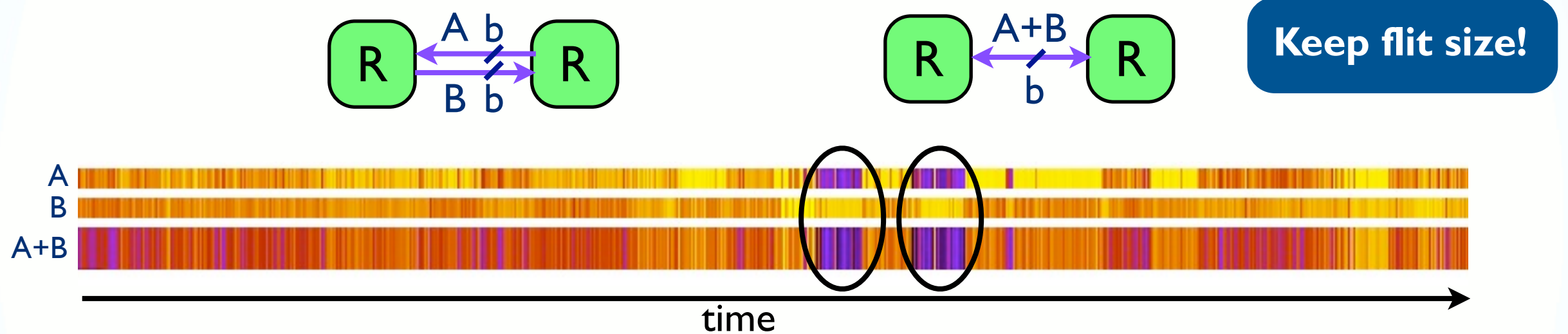
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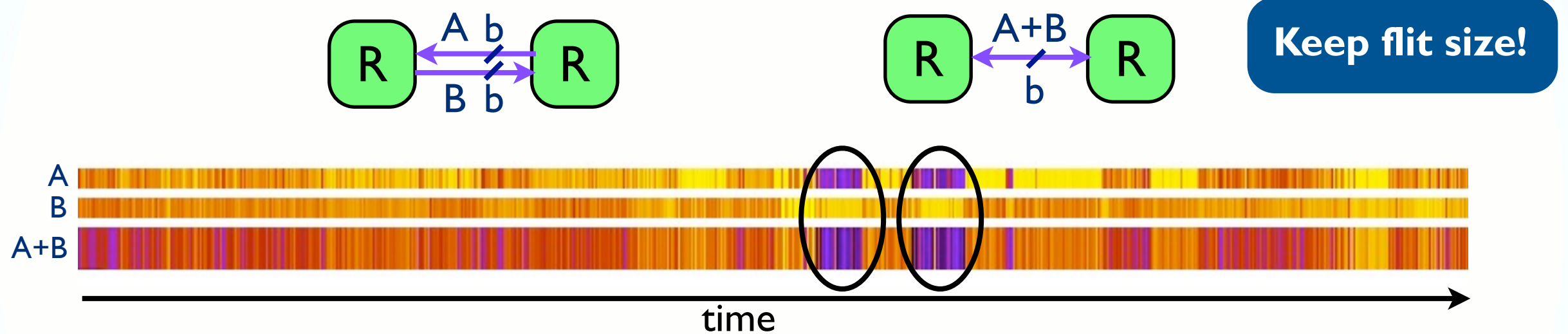
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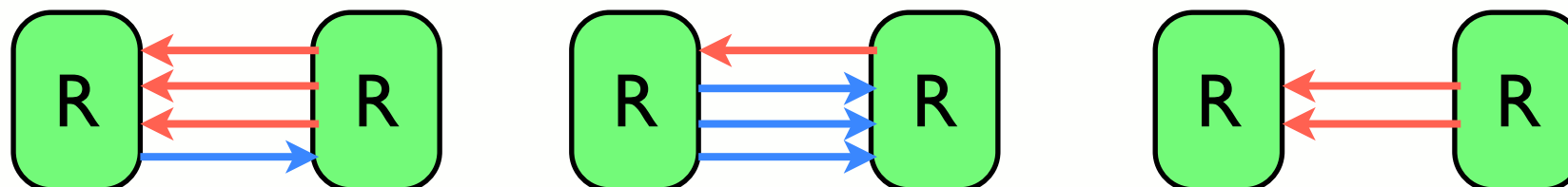


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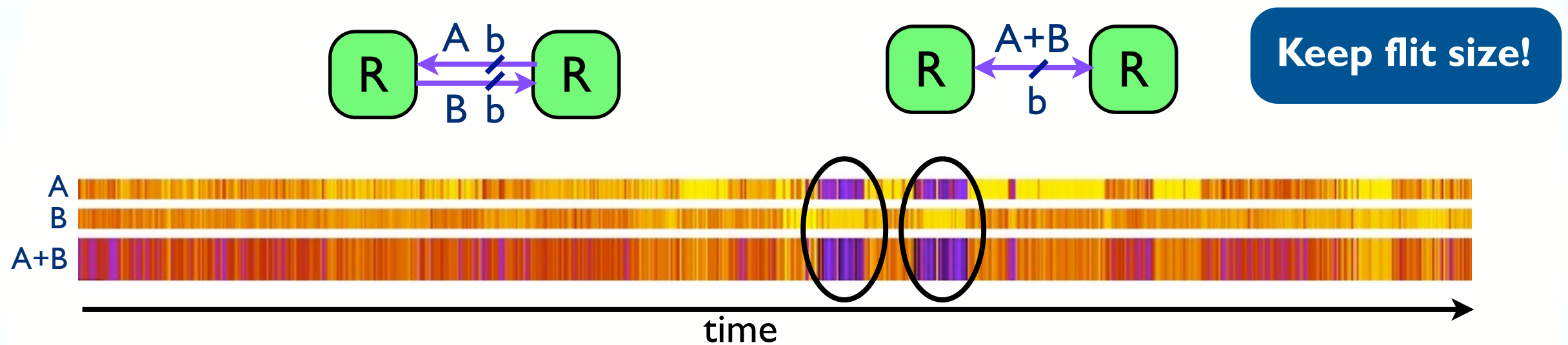


- Adding flexibility with fine-grained BW adaptivity

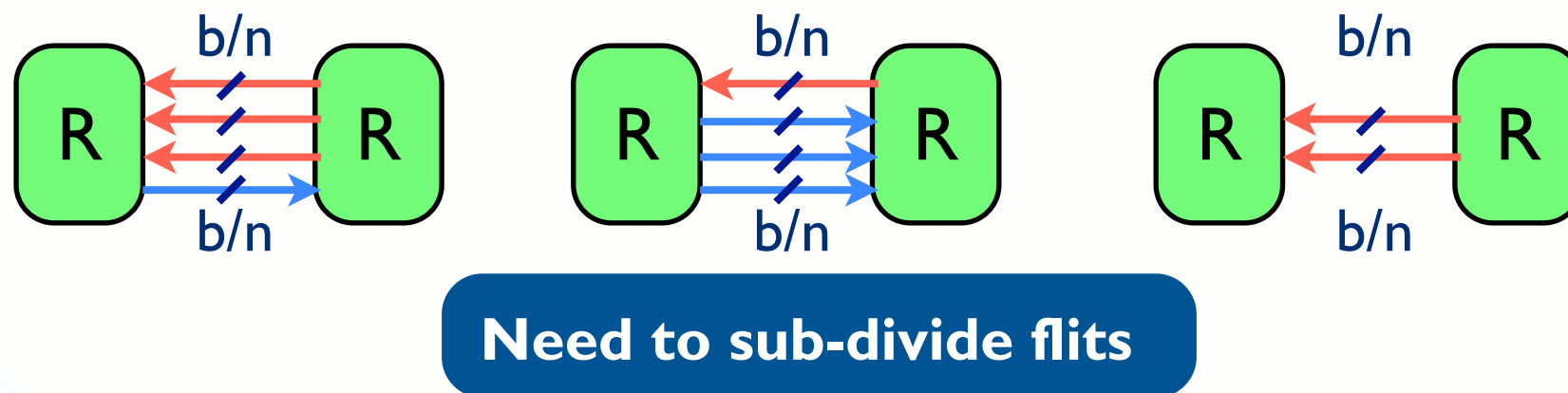


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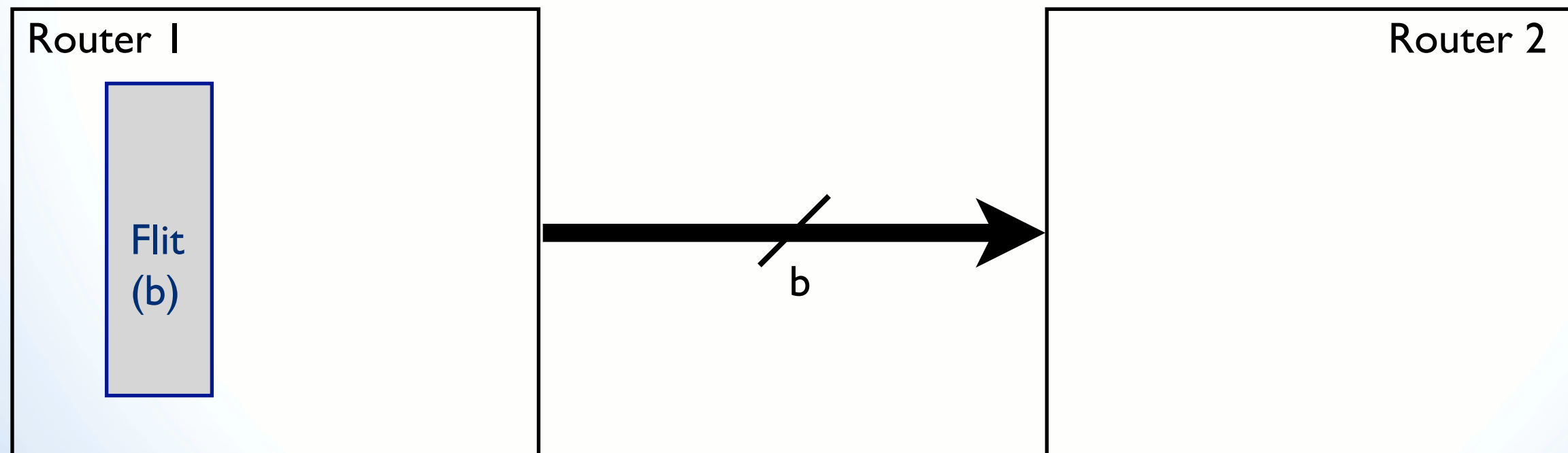


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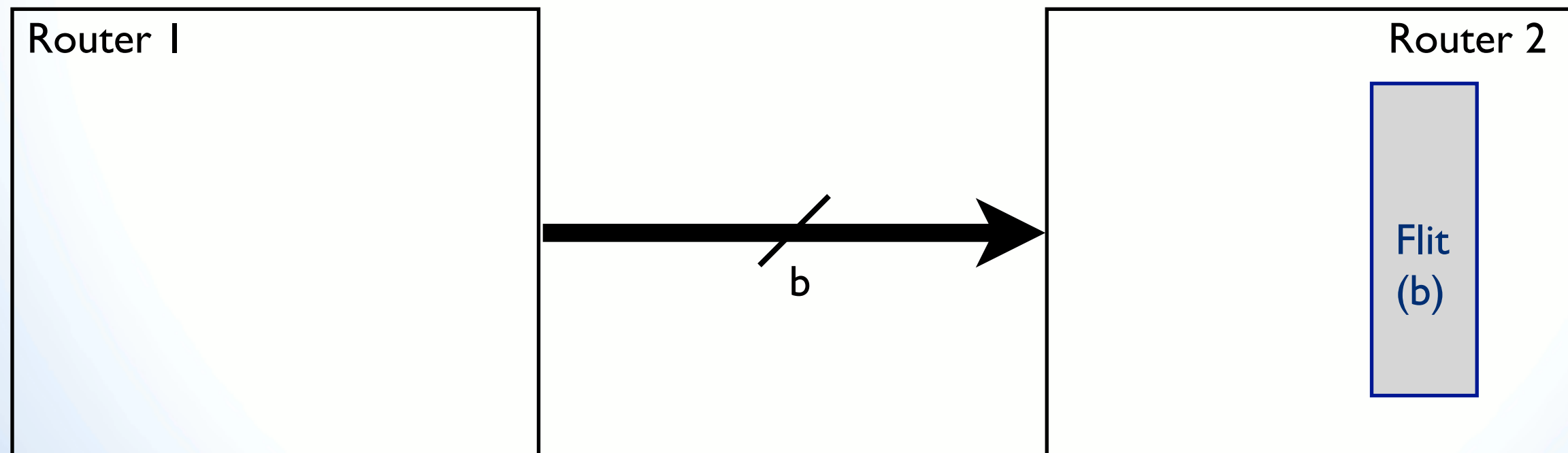
Decoupling Flit Width From Channel Width

- Conventionally in NoC, flit width is coupled to channel width



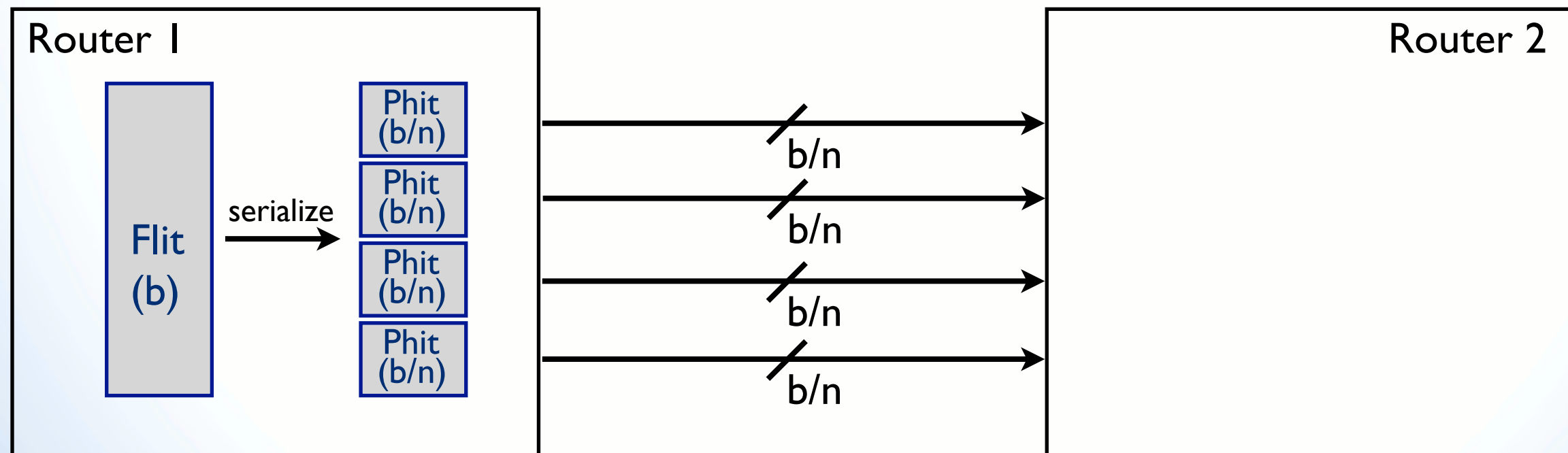
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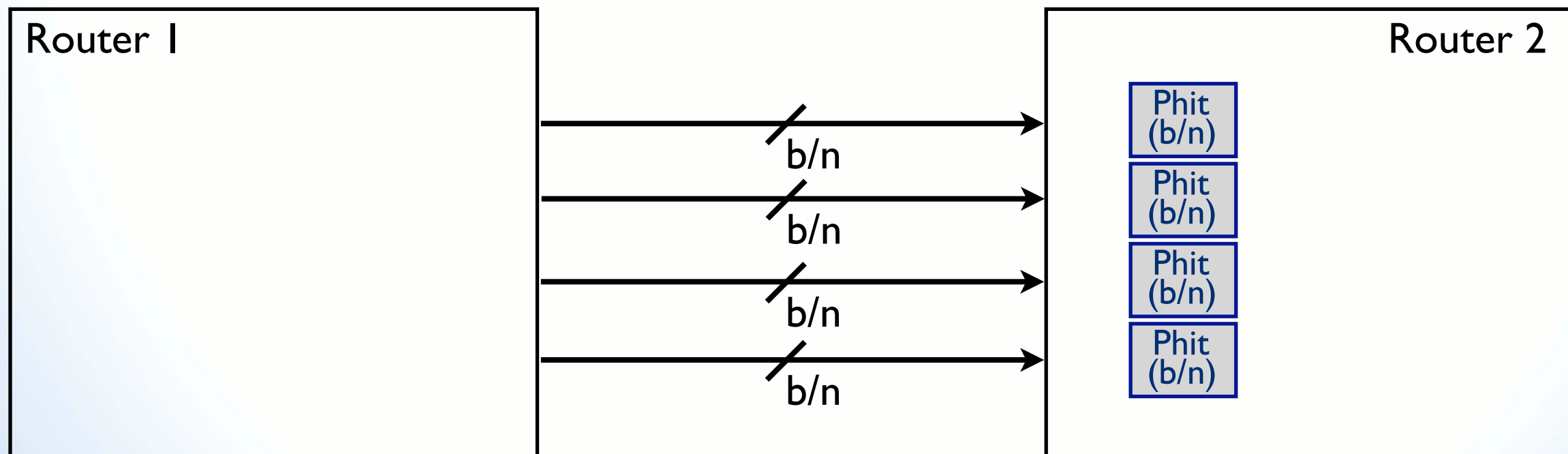
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- Break flits (flow control units) into phits (physical transfer units) to decouple channel width from flit width



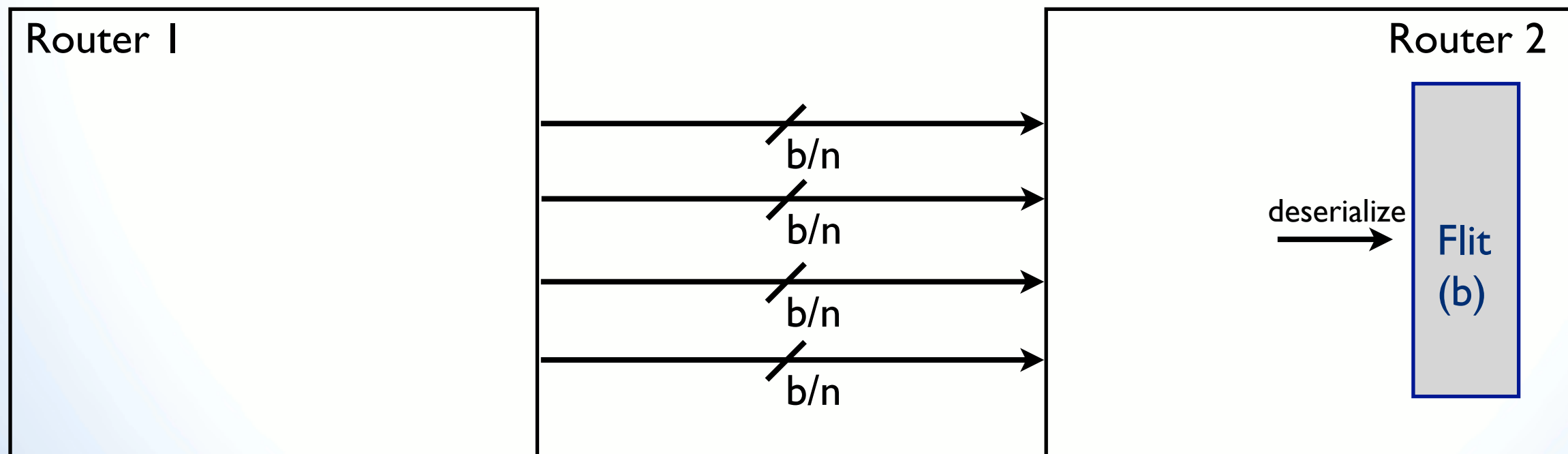
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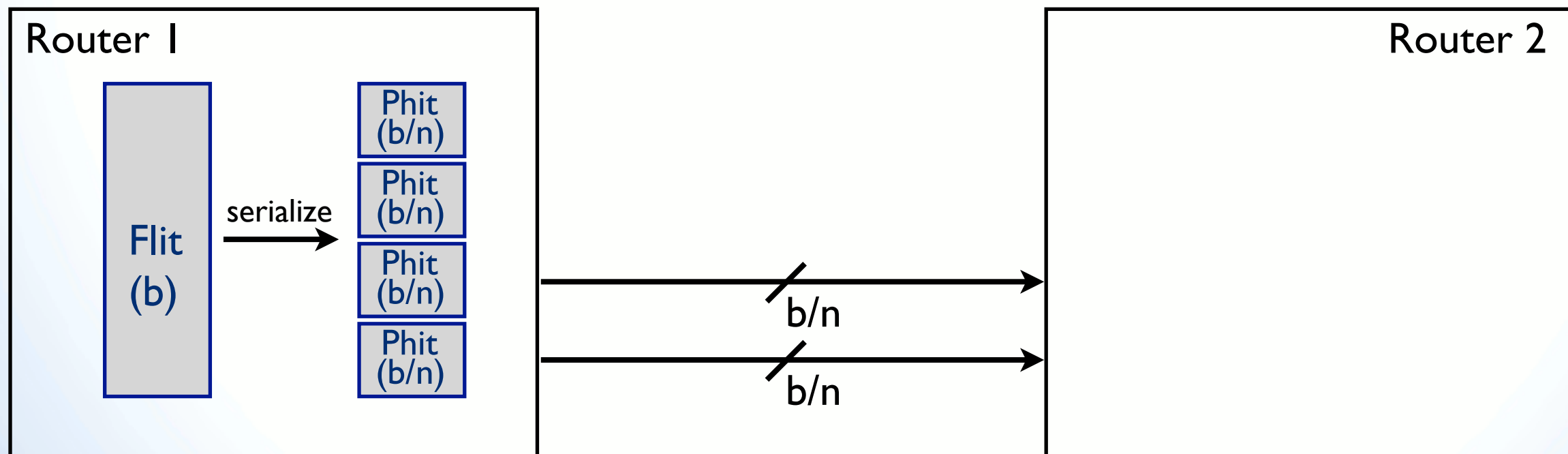
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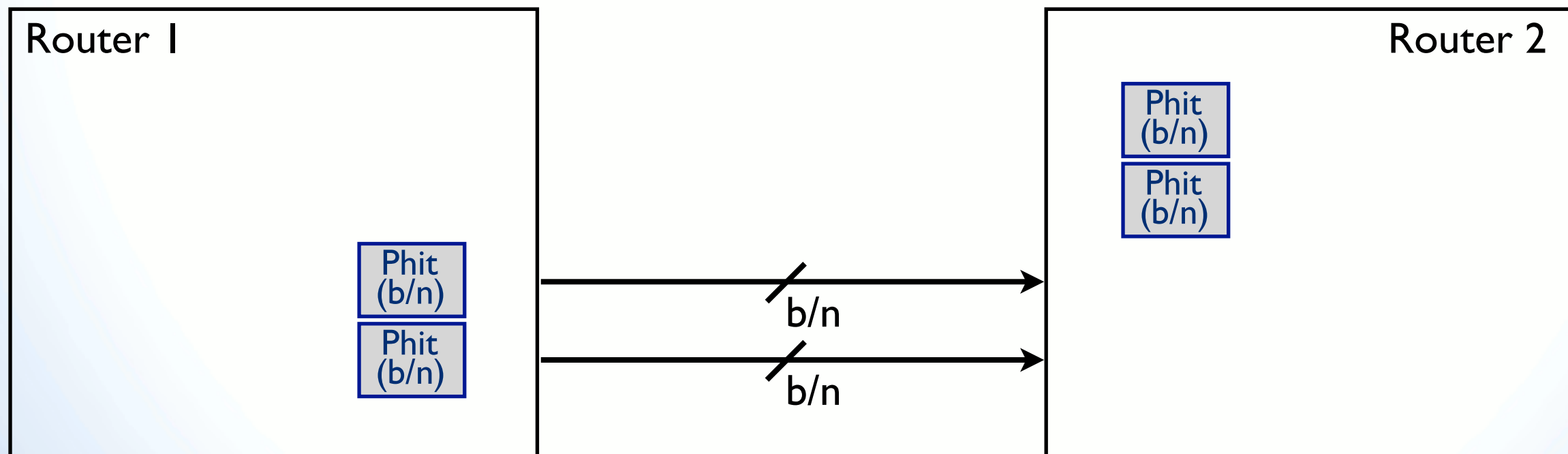
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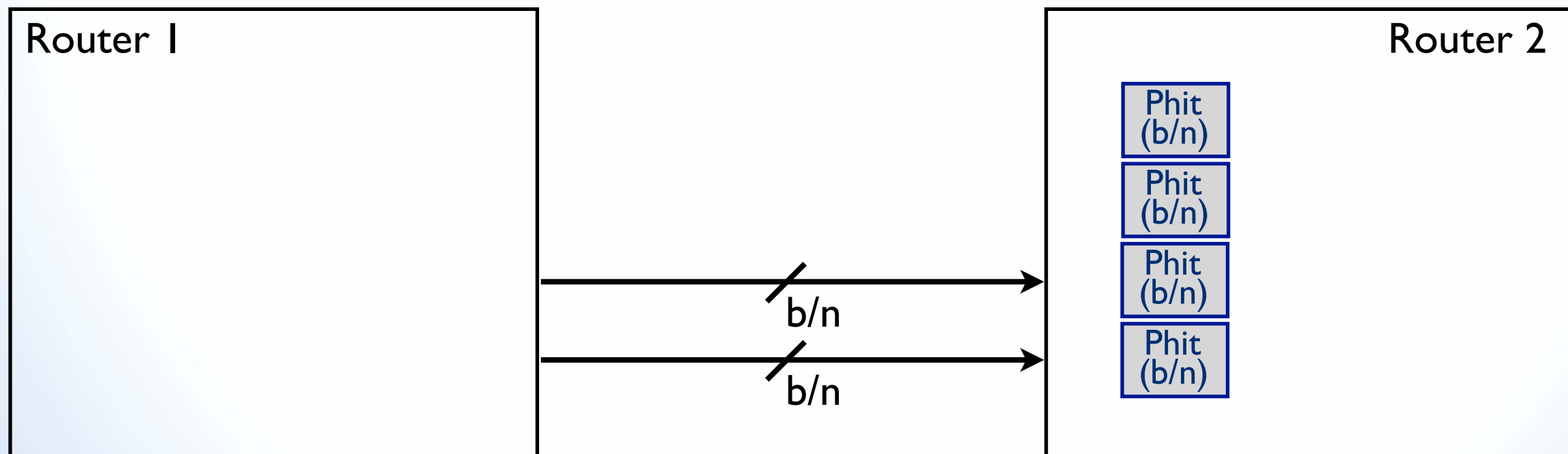
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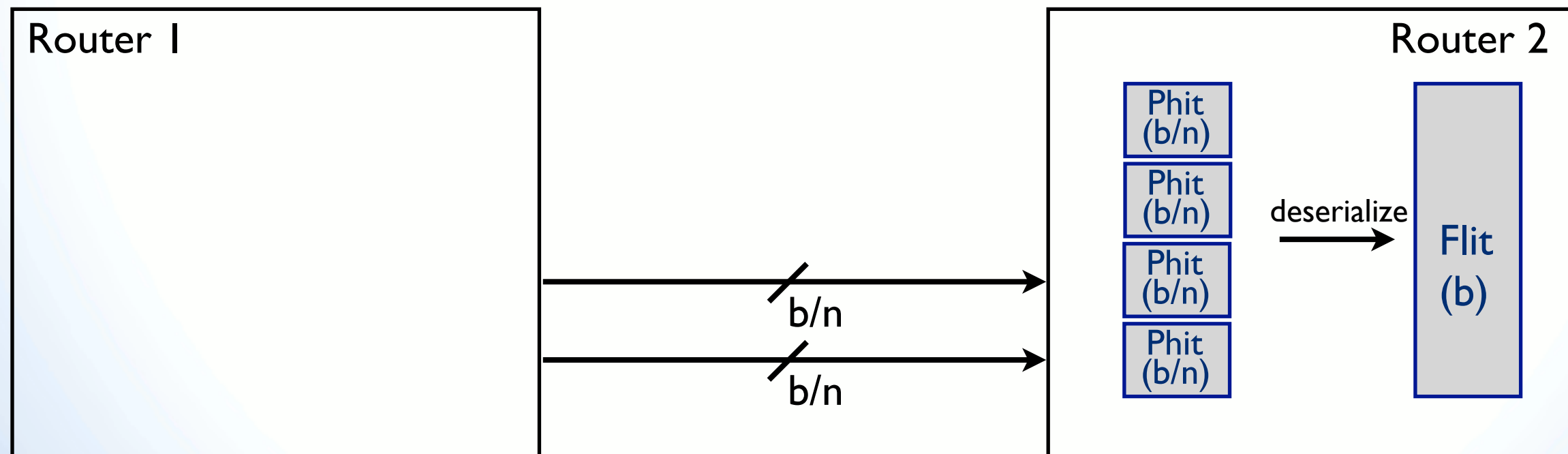
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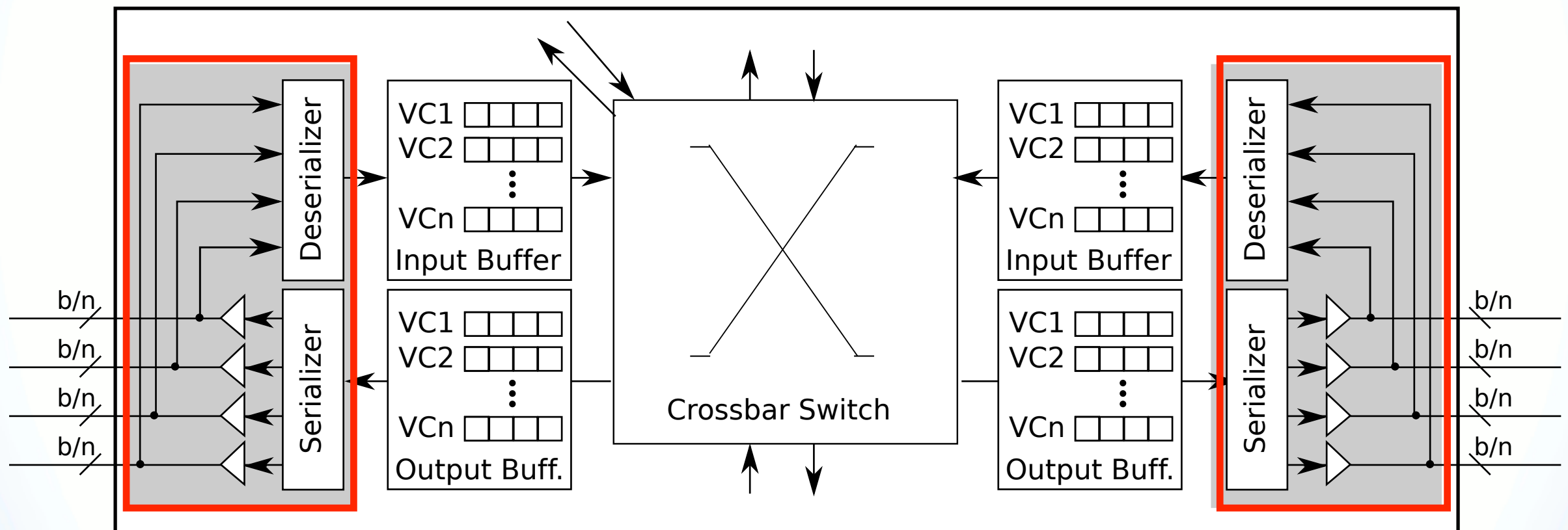
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Microarchitecture

- Bandwidth-Addaptive Router (BAR): Only minimal modifications to standard VC router necessary

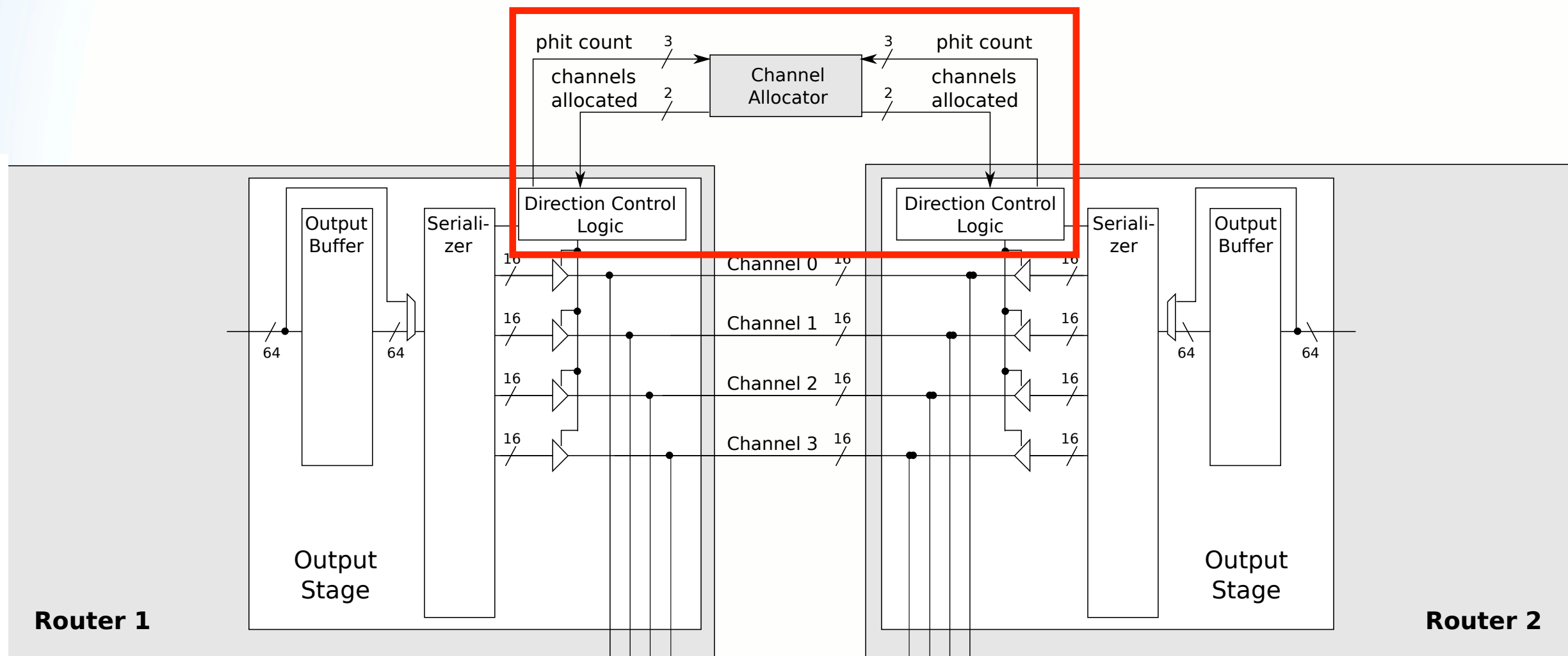


- Intra- & inter-router flow control is still flit-based



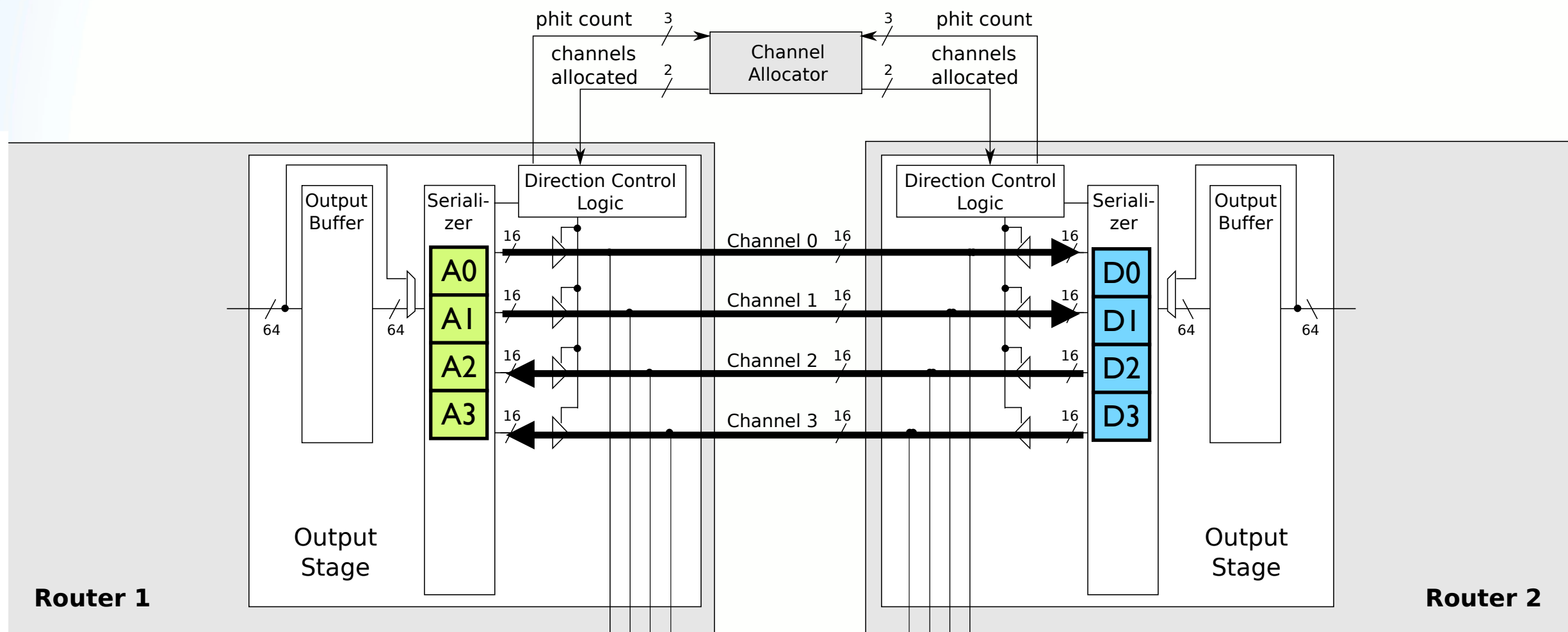
Bandwidth Allocation

- Pressure-based allocation of channels to directions



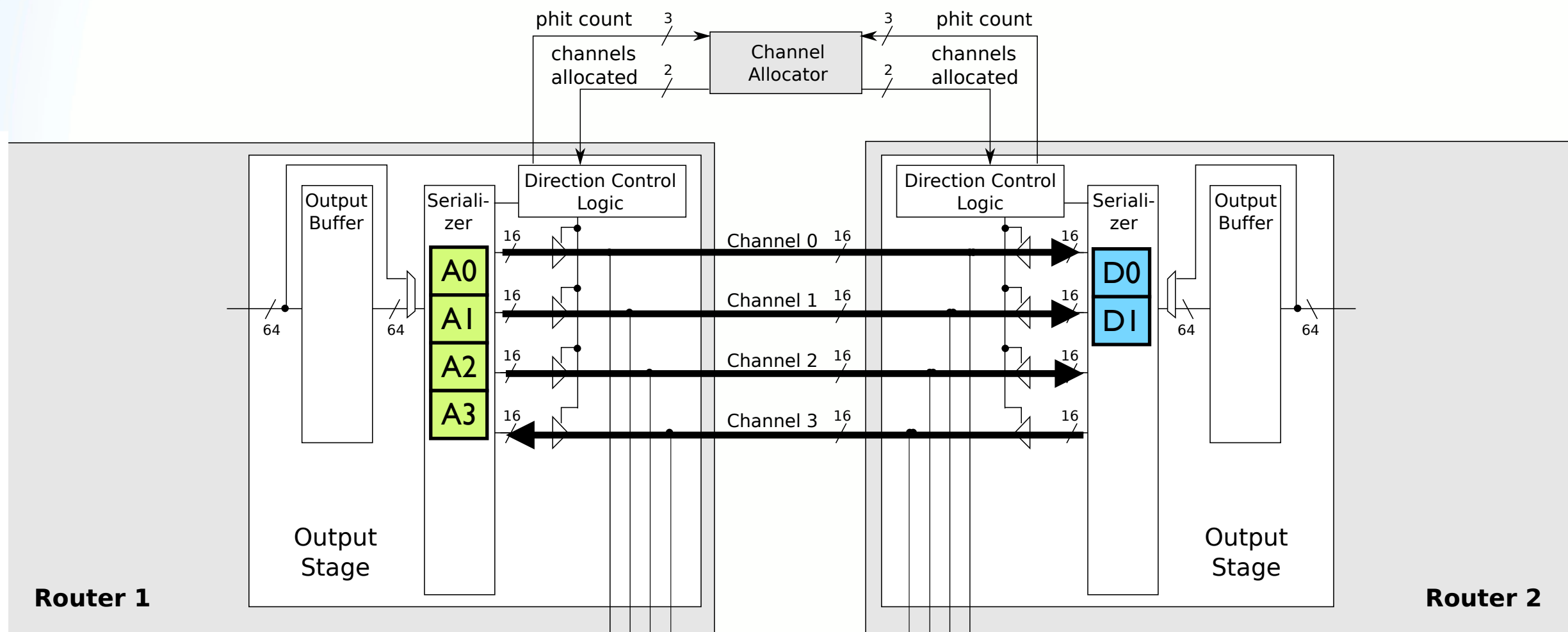
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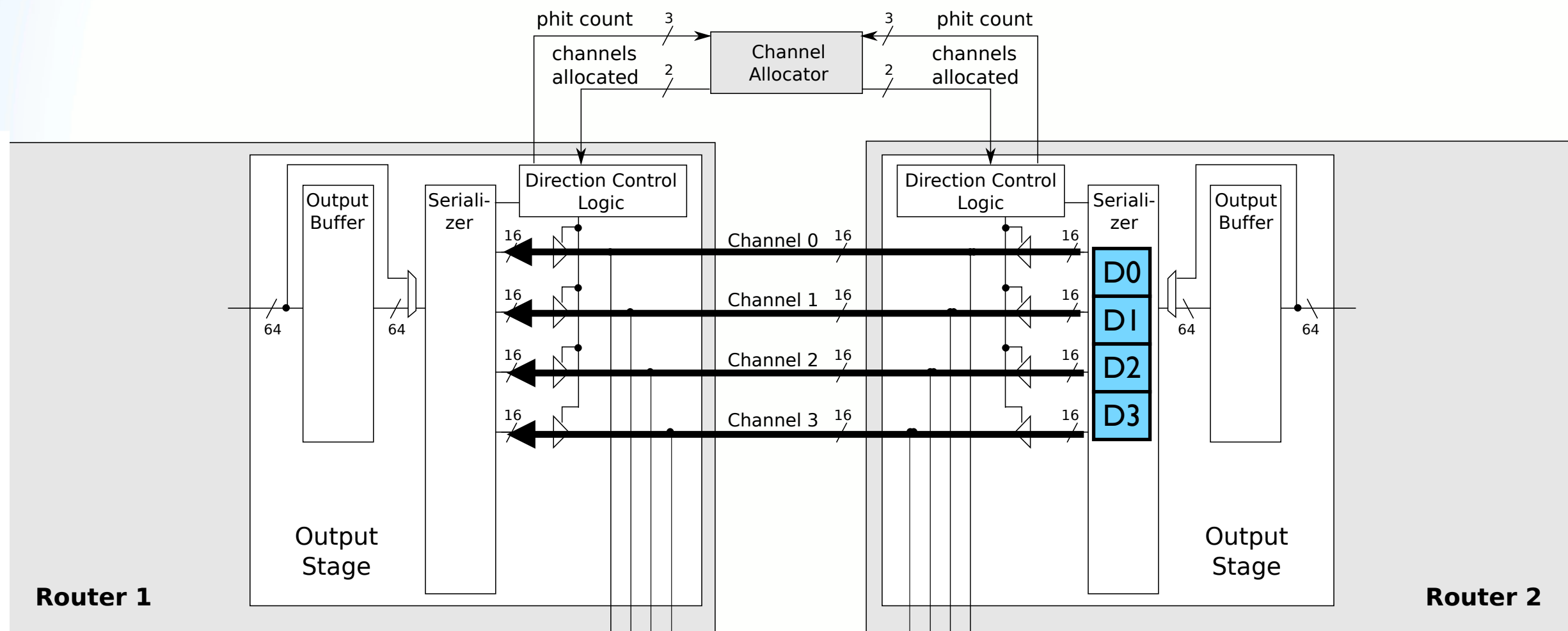
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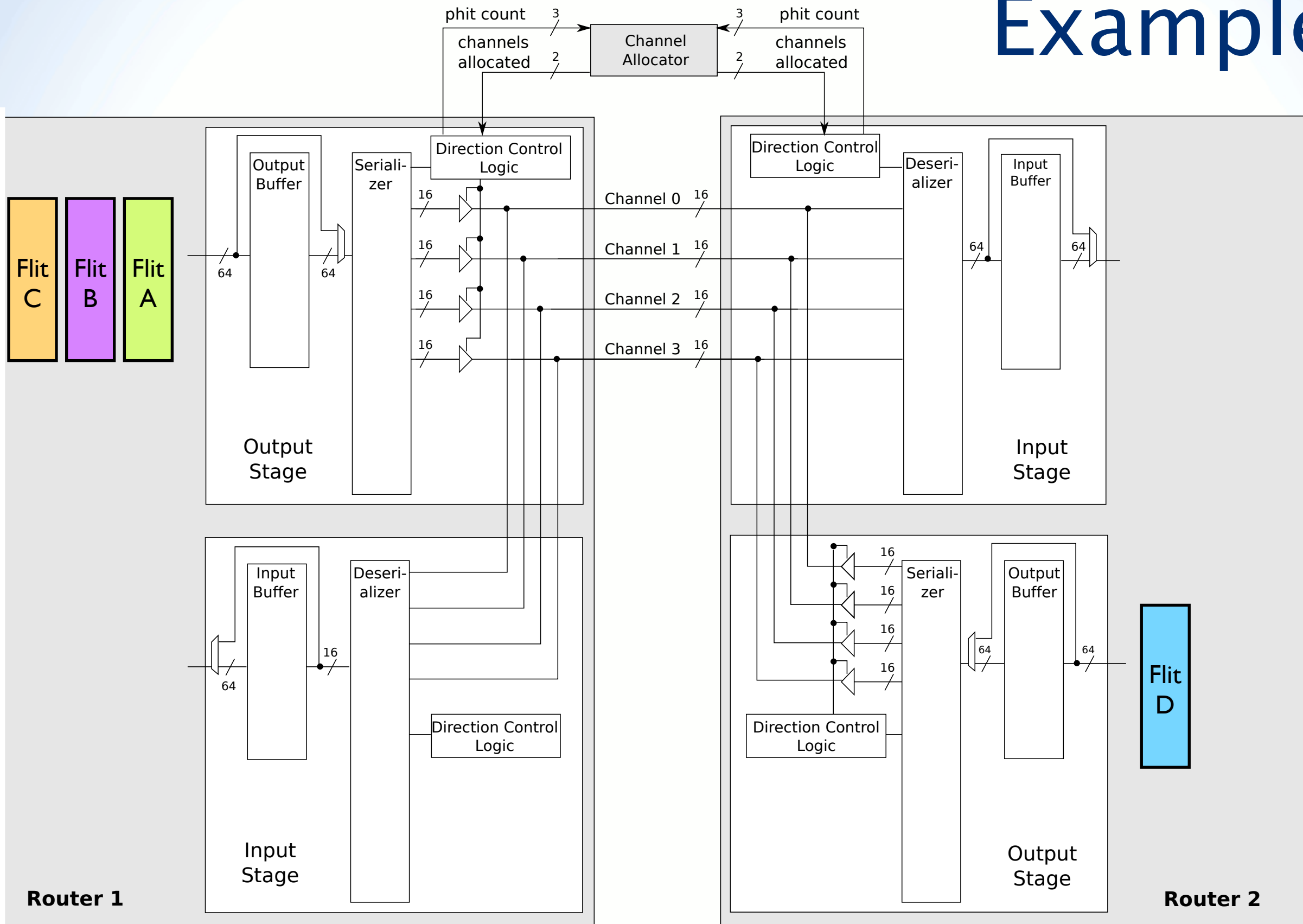


Bandwidth Allocation

- Pressure-based allocation of channels to directions



Example



Router 1

Router 2

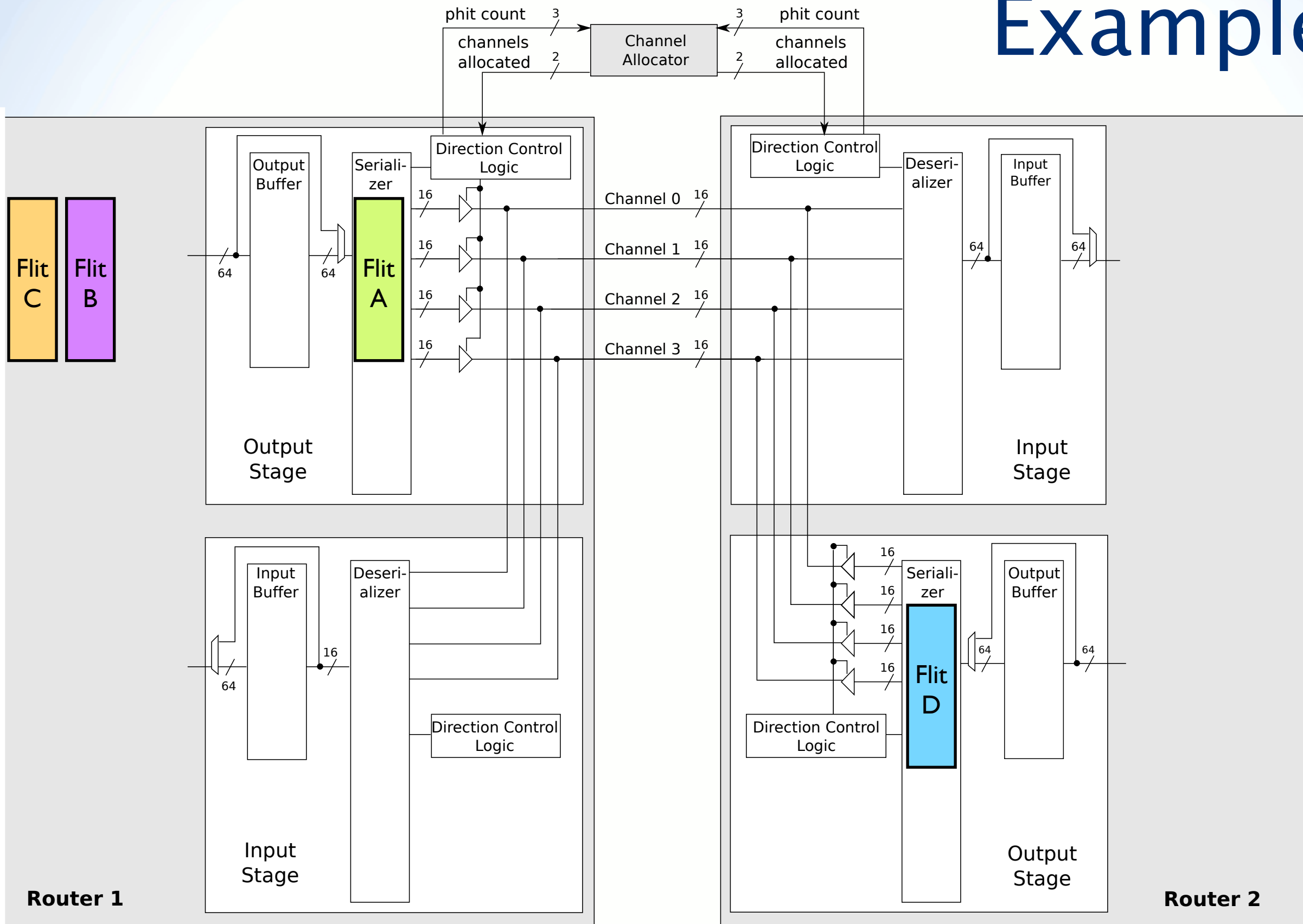
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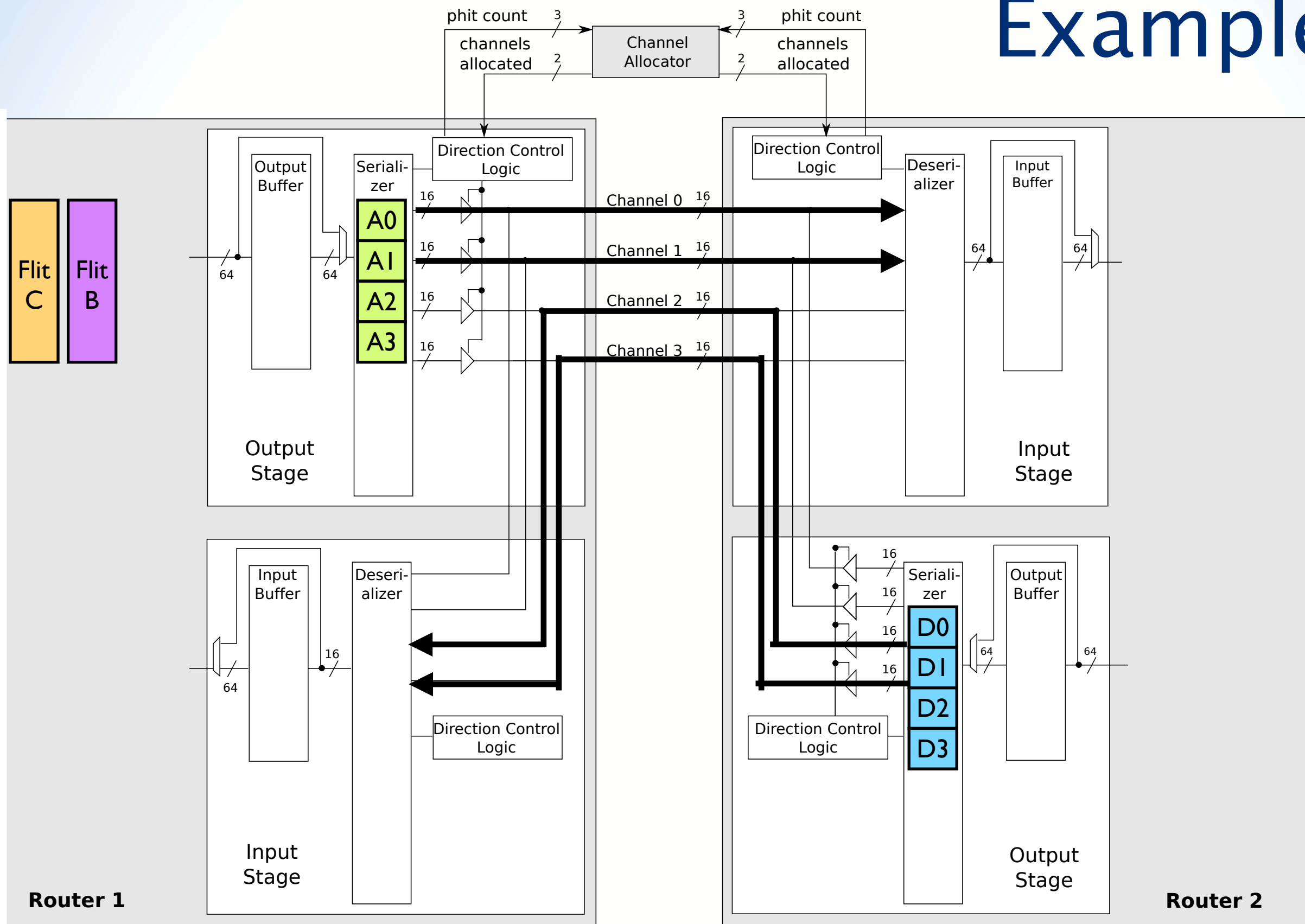
University of Toronto



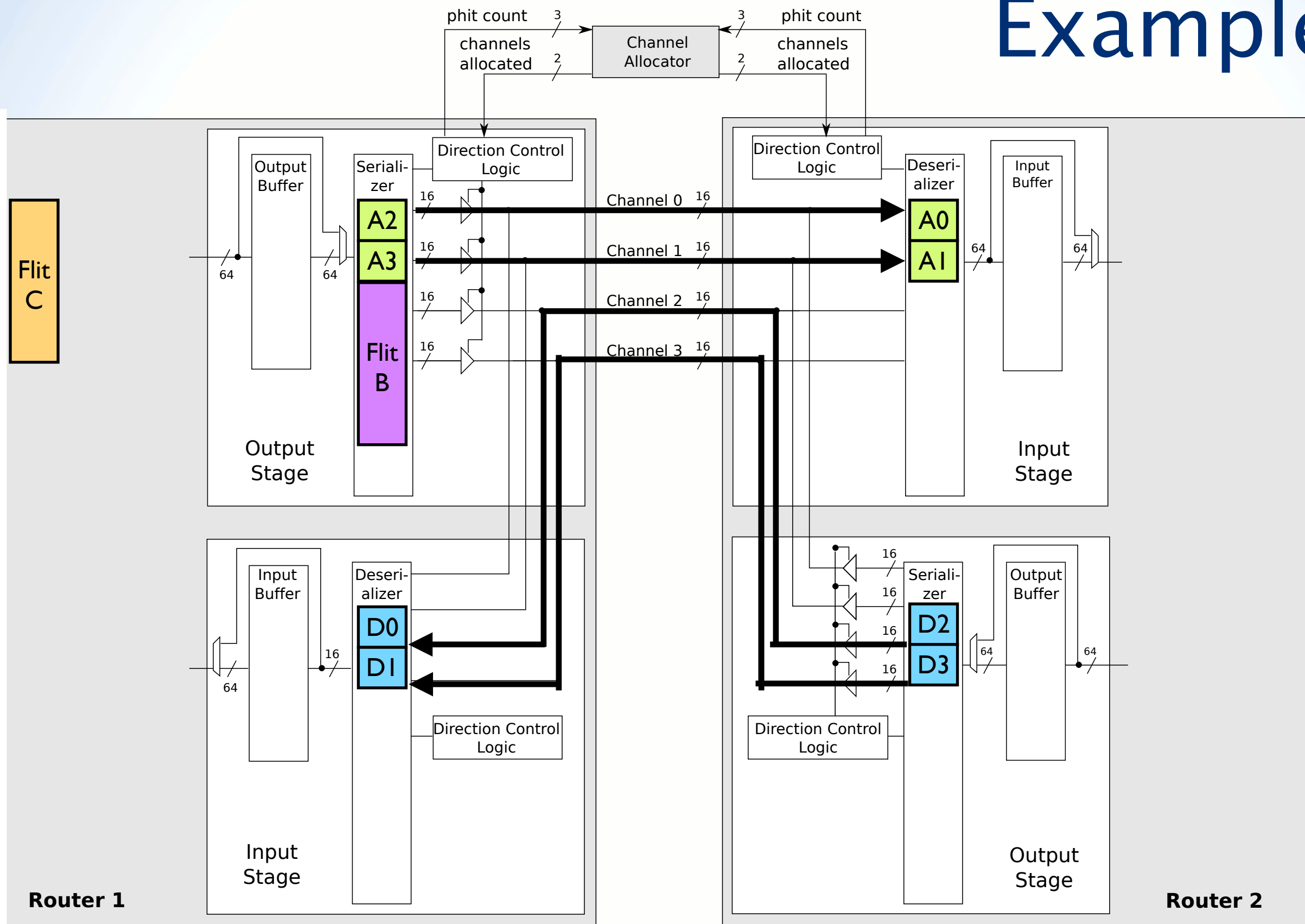
Example



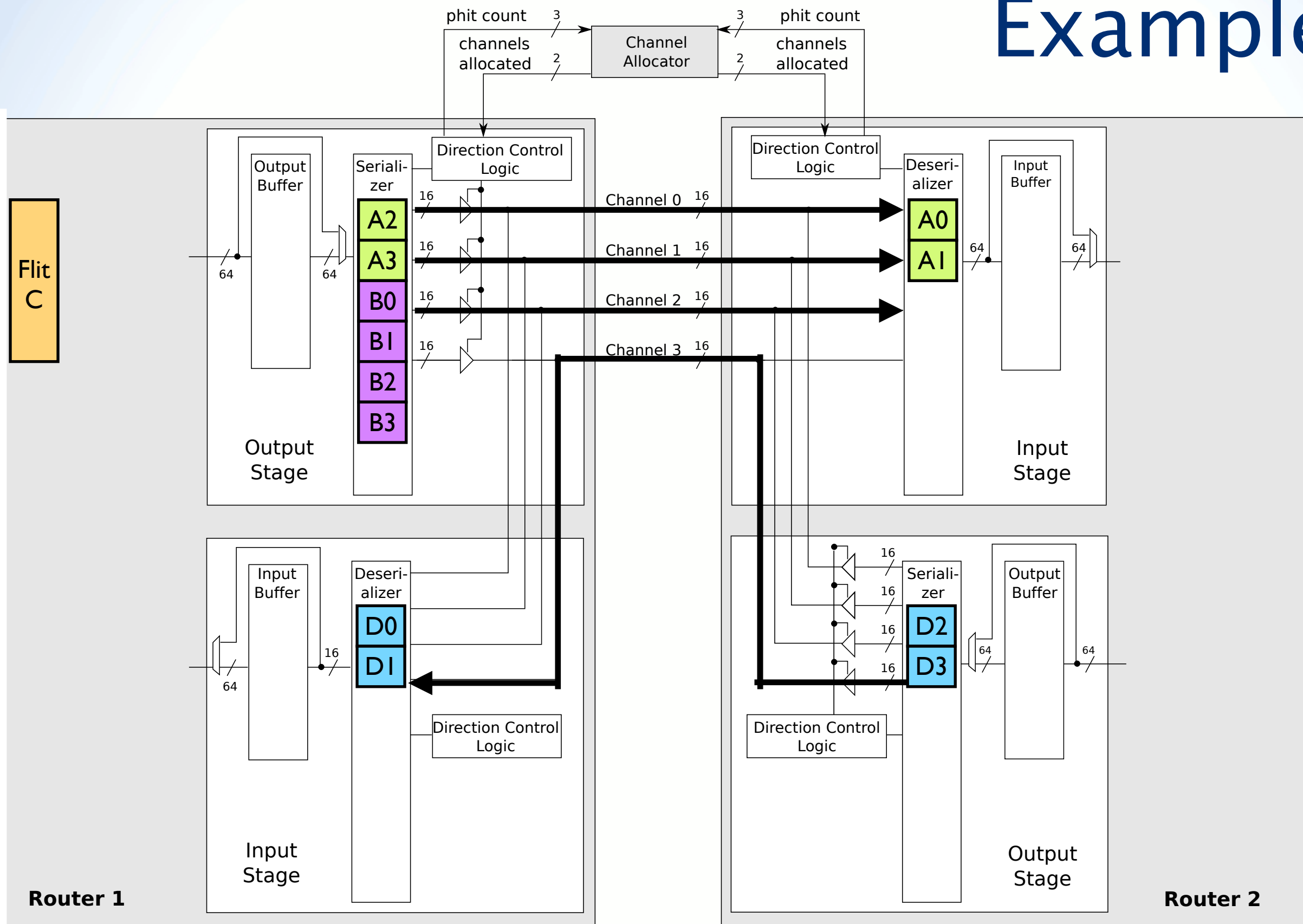
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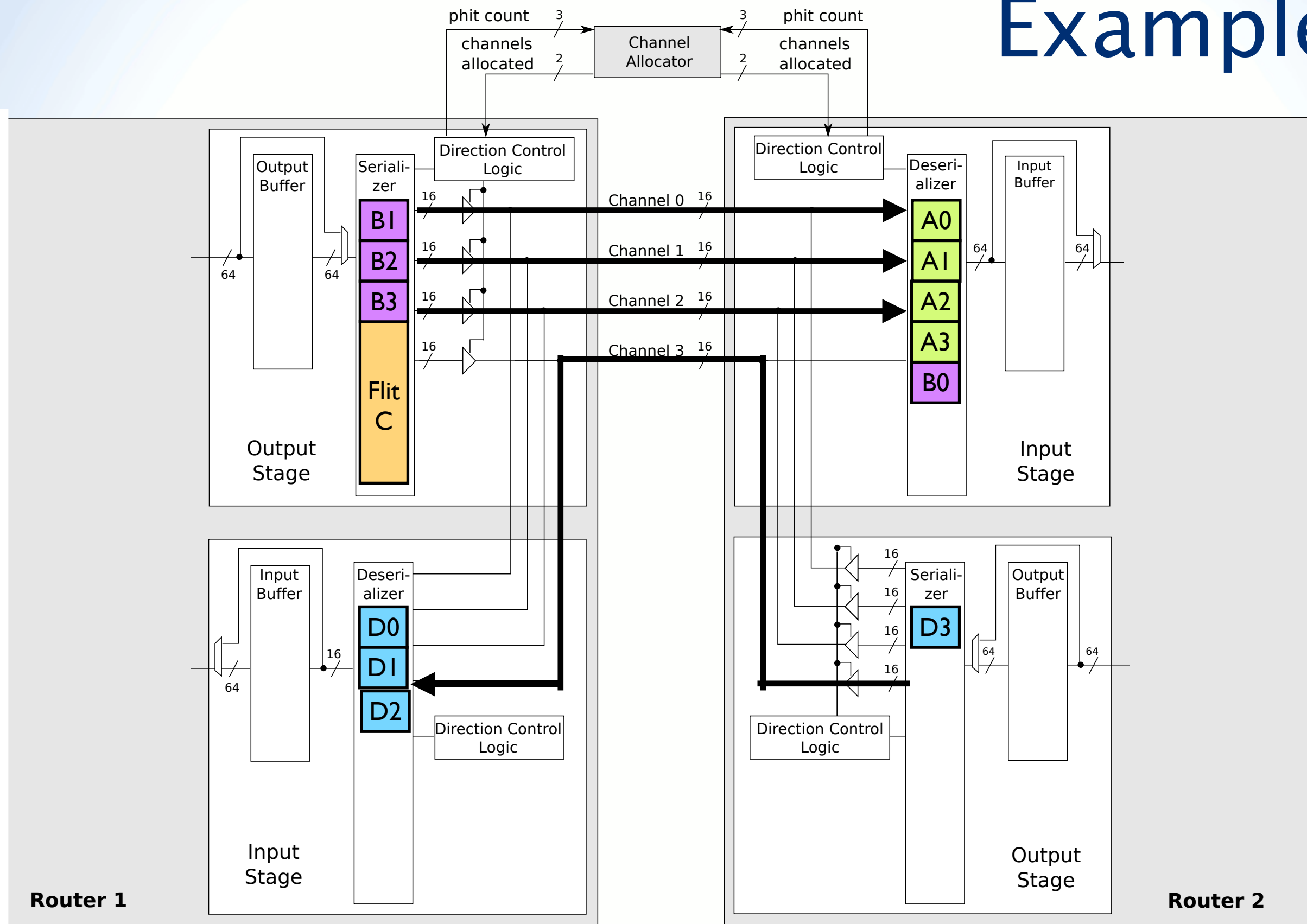
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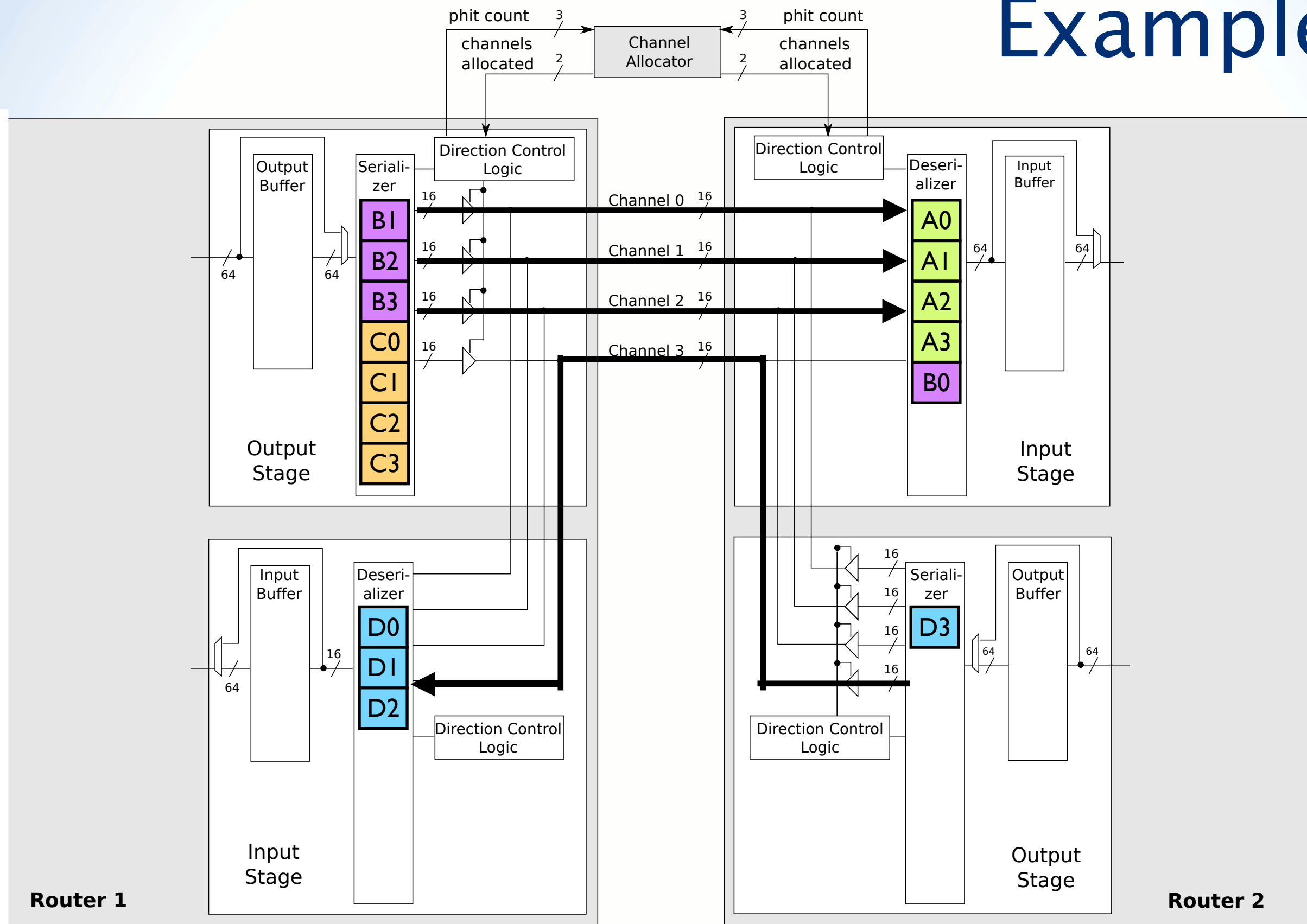


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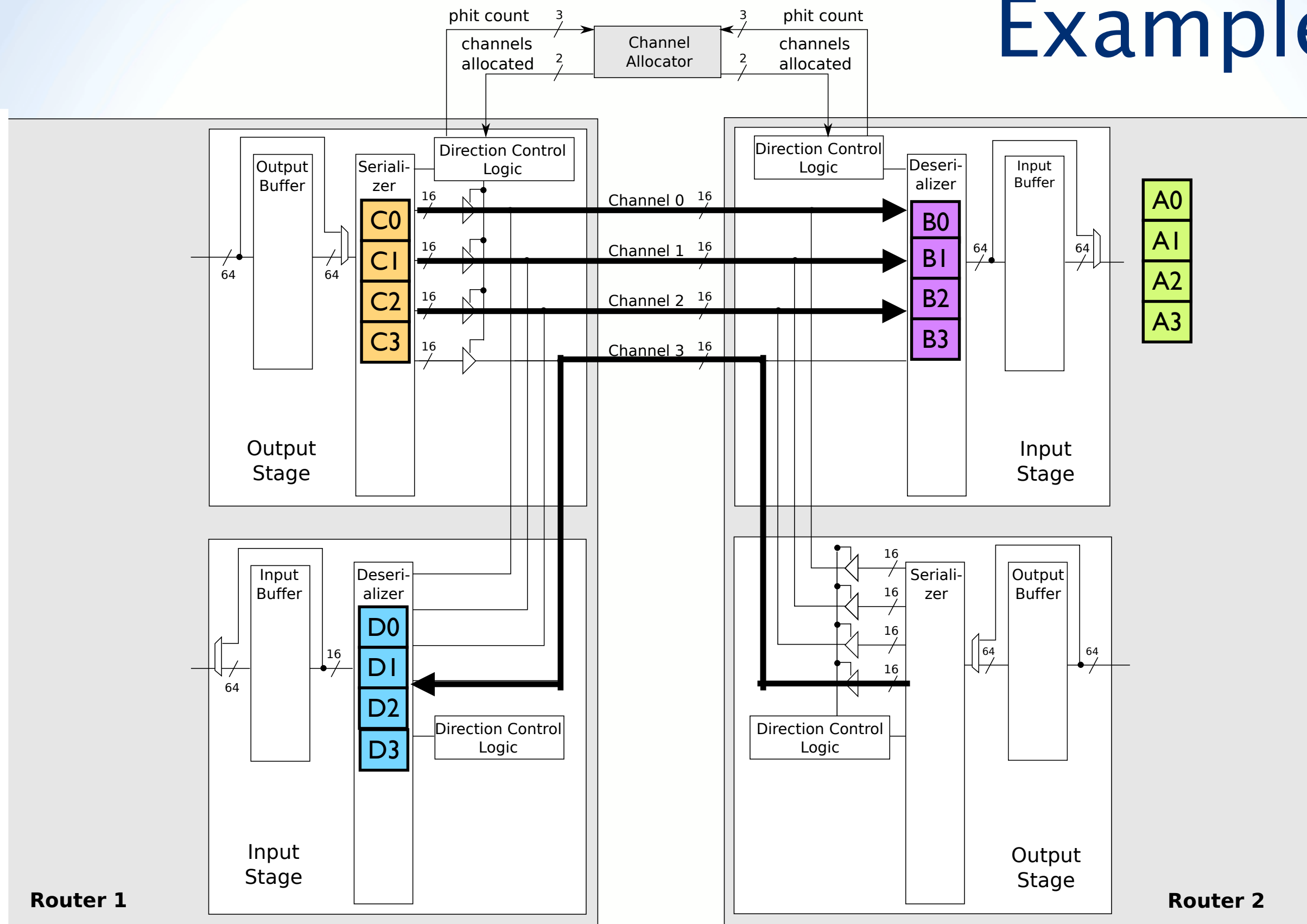
Router 2



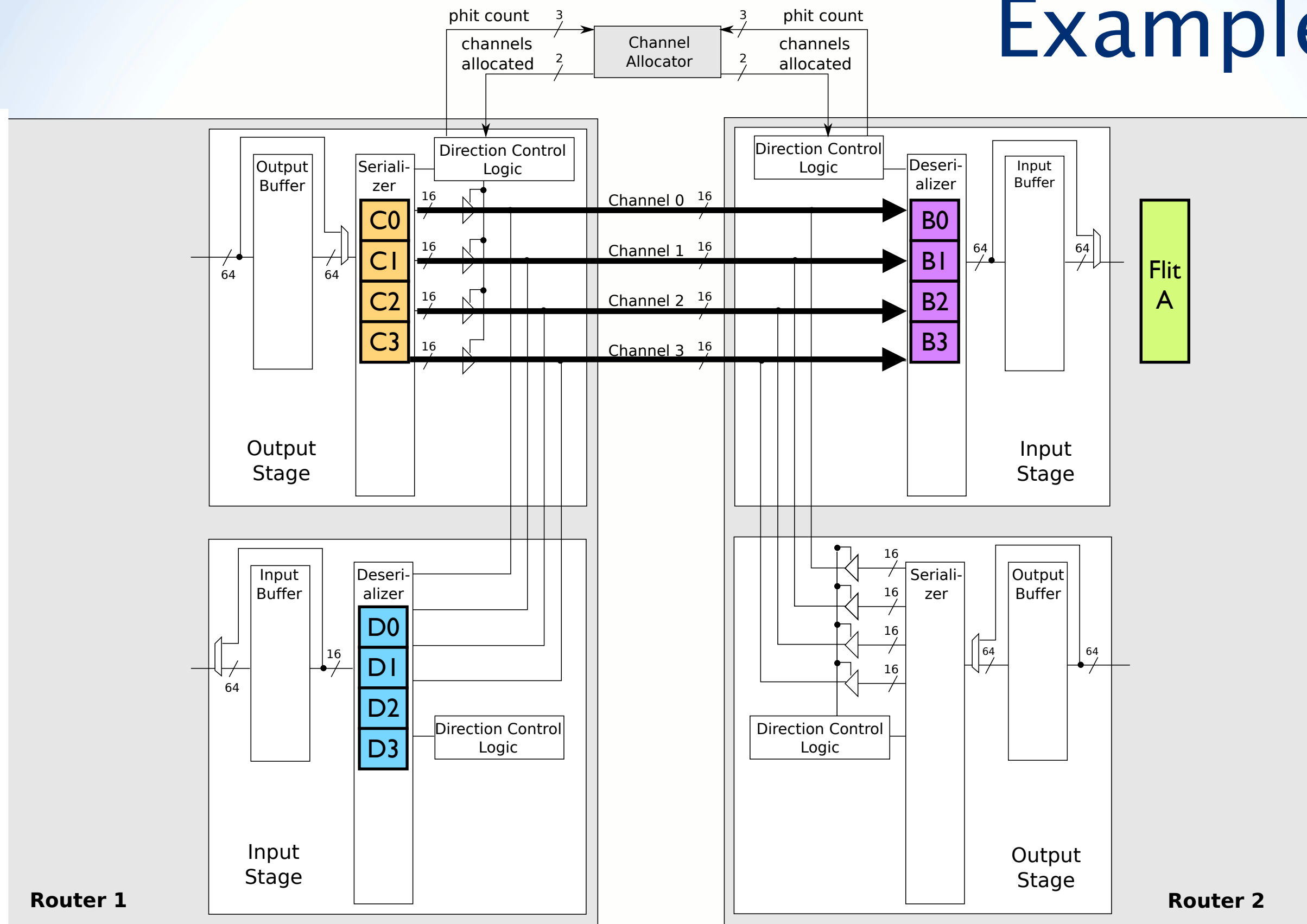
Example



Example



Example

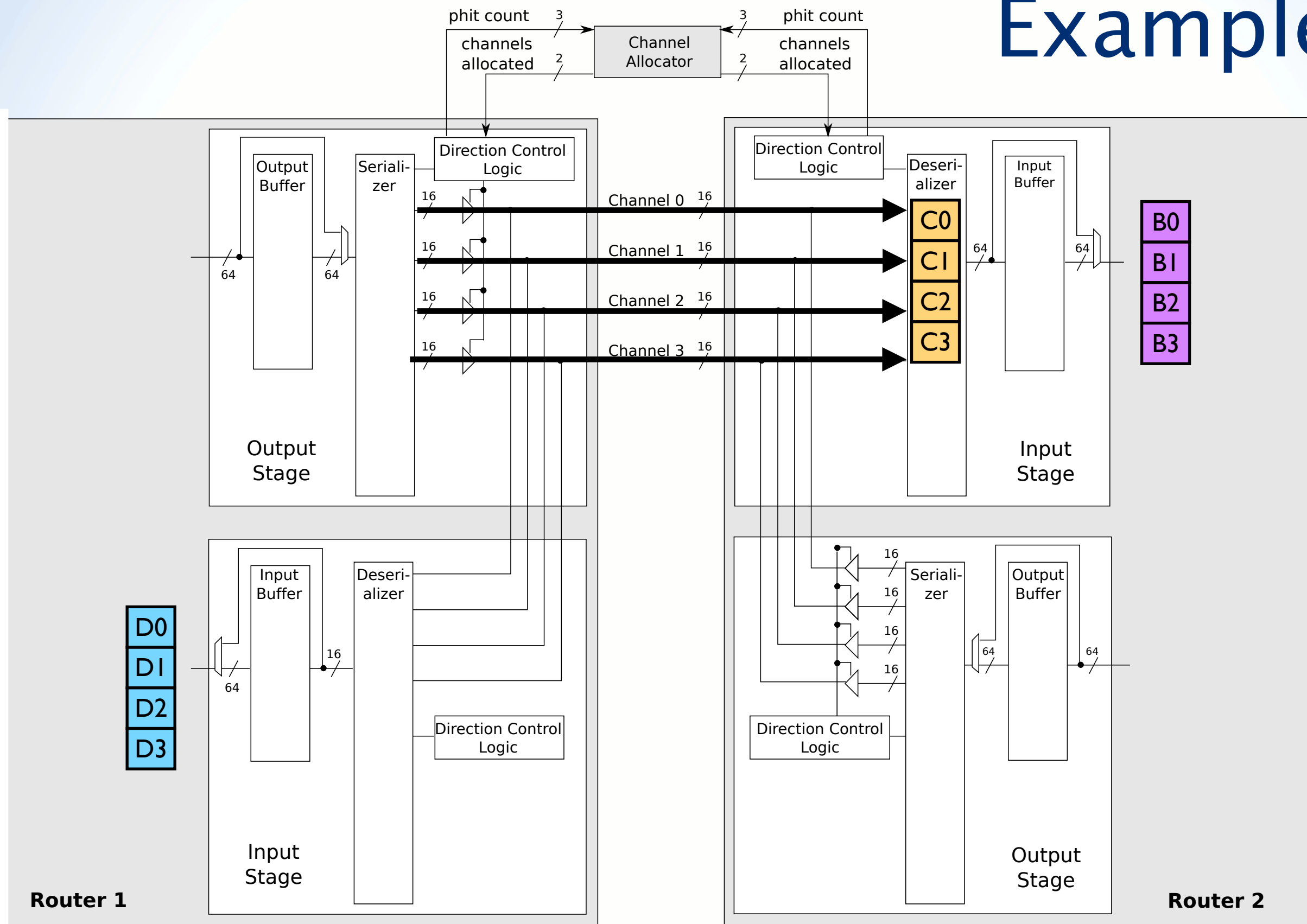


Router 1

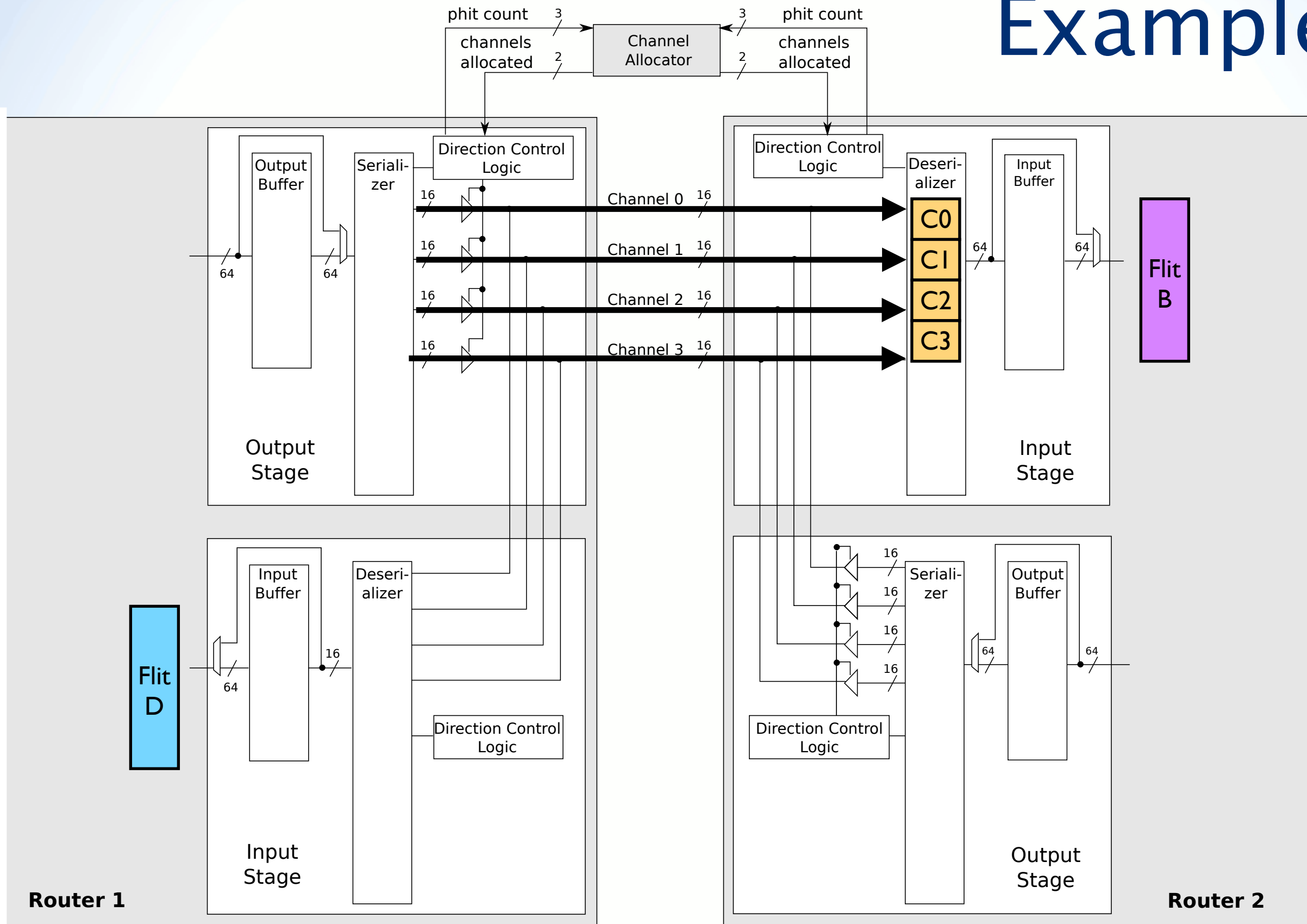
Router 2



Example



Example

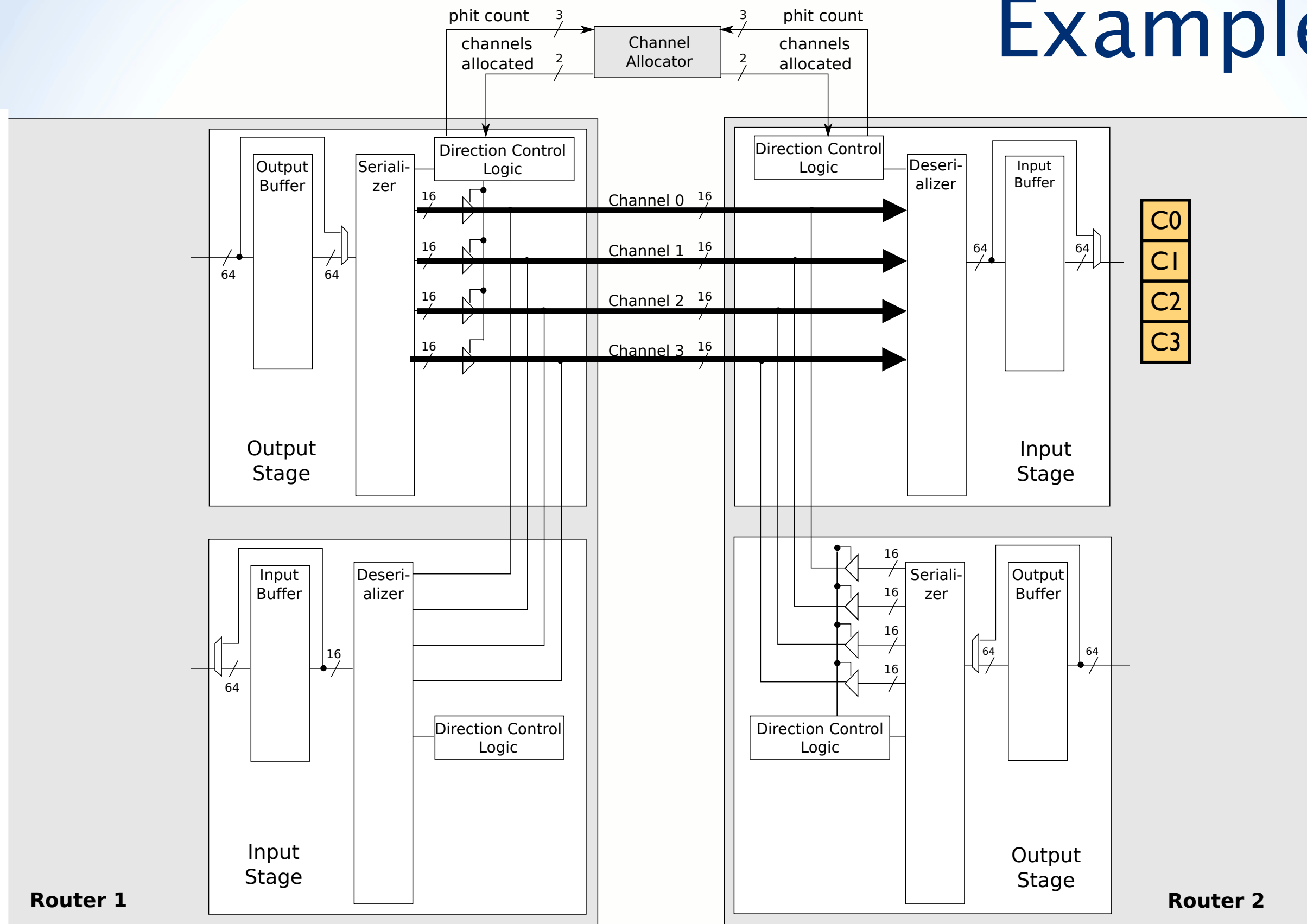


Router 1

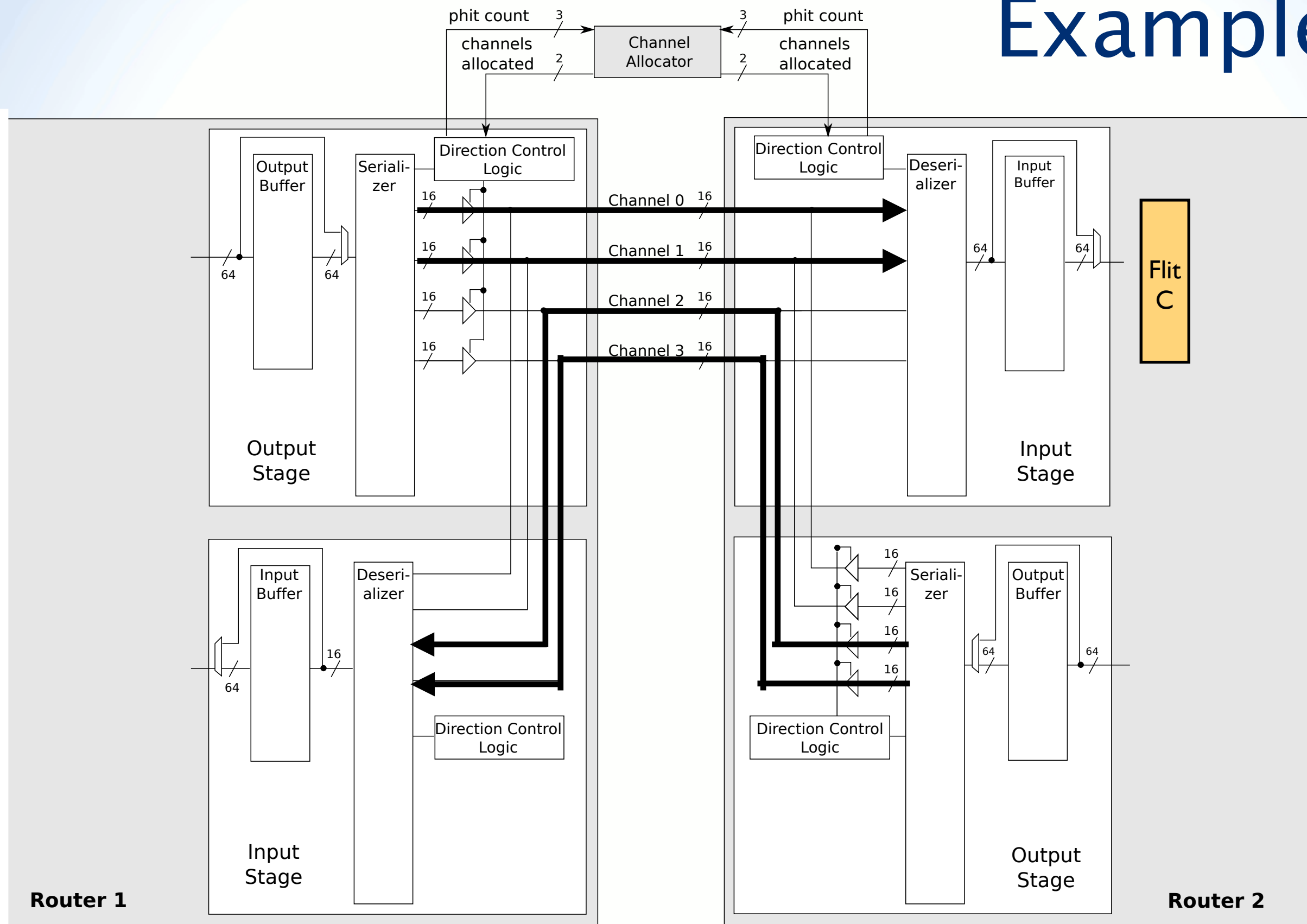
Router 2



Example



Example



Related Work



Related Work

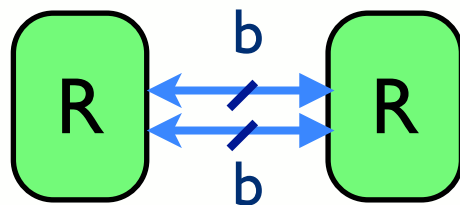
- Previous Work:



Related Work

- Previous Work:
 - BiNoC (Lan et al., NOCS 2009)
 - only coarse-grained BW adaptivity

BINOC



Link BW: $2*b$

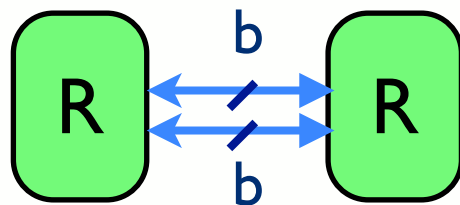
XBar ports: $2*P$



Related Work

- Previous Work:
 - BiNoC (Lan et al., NOCS 2009)
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 - Oblivious Routing in On-Chip Bandwidth-Adaptive Networks (Cho et al., PACT 2009)

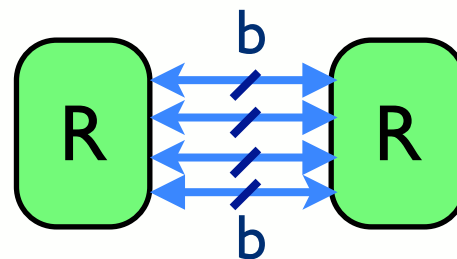
BINOC



Link BW: $2*b$

XBar ports: $2*P$

BWADAPTIVE



Link BW: $N*b$

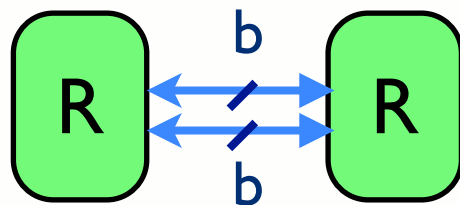
XBar ports: $N*P$



Related Work

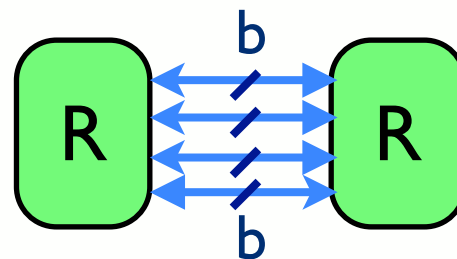
- Previous Work:
 - BiNoC (Lan et al., NOCS 2009)
 - only coarse-grained BW adaptivity
 - Oblivious Routing in On-Chip Bandwidth-Adaptive Networks (Cho et al., PACT 2009)
- Objective is different:

BINOC



Link BW: $2*b$
XBar ports: $2*P$

BWADAPTIVE



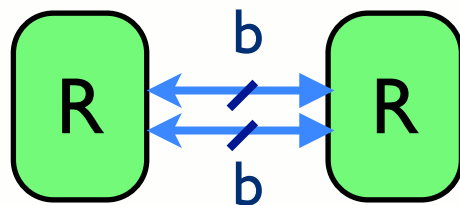
Link BW: $N*b$
XBar ports: $N*P$



Related Work

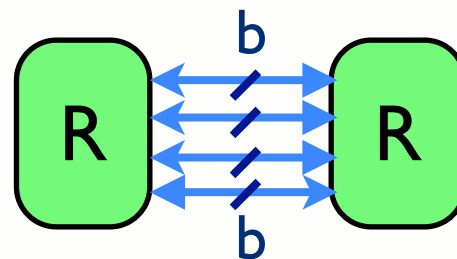
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- Objective is different:

BINOC



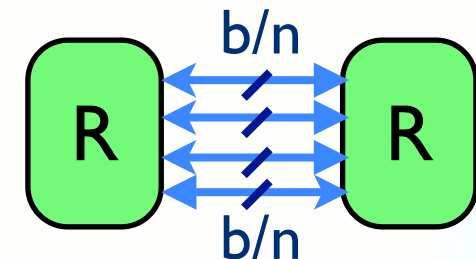
Link BW: $2*b$
XBar ports: $2*P$

BWADAPTIVE



Link BW: $N*b$
XBar ports: $N*P$

BAR



Link BW: $\leq b$
XBar ports: P

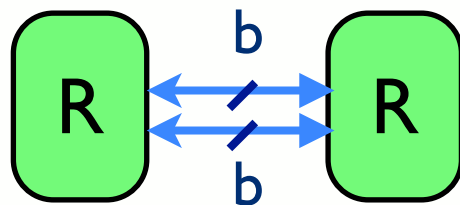


Related Work

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- Objective is different:

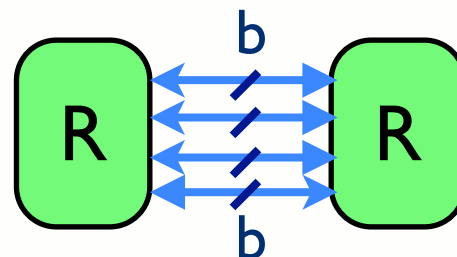
Remember: BW demands $\ll b$

BINOC



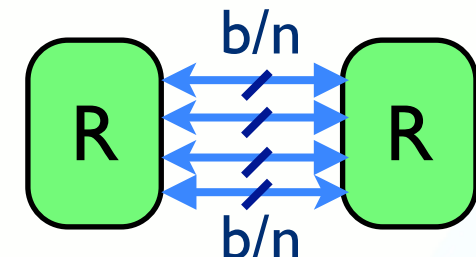
Link BW: $2*b$
XBar ports: $2*P$

BWADAPTIVE



Link BW: $N*b$
XBar ports: $N*P$

BAR



Link BW: $\leq b$
XBar ports: P



Evaluation

- Synthetic and real workloads (**PARSEC**)
- Comparing 4 router designs:
 - **BAR**: our BW adaptive router design
 - **STANDARD**: typical virtual-channel router
 - **BINOC**: BiNoC router
 - **BWADAPTIVE**: existing BW adaptive router



Area Comparison

May 10, 2012

15

University of Toronto



Area Comparison

- Orion 2.0 with 45nm process

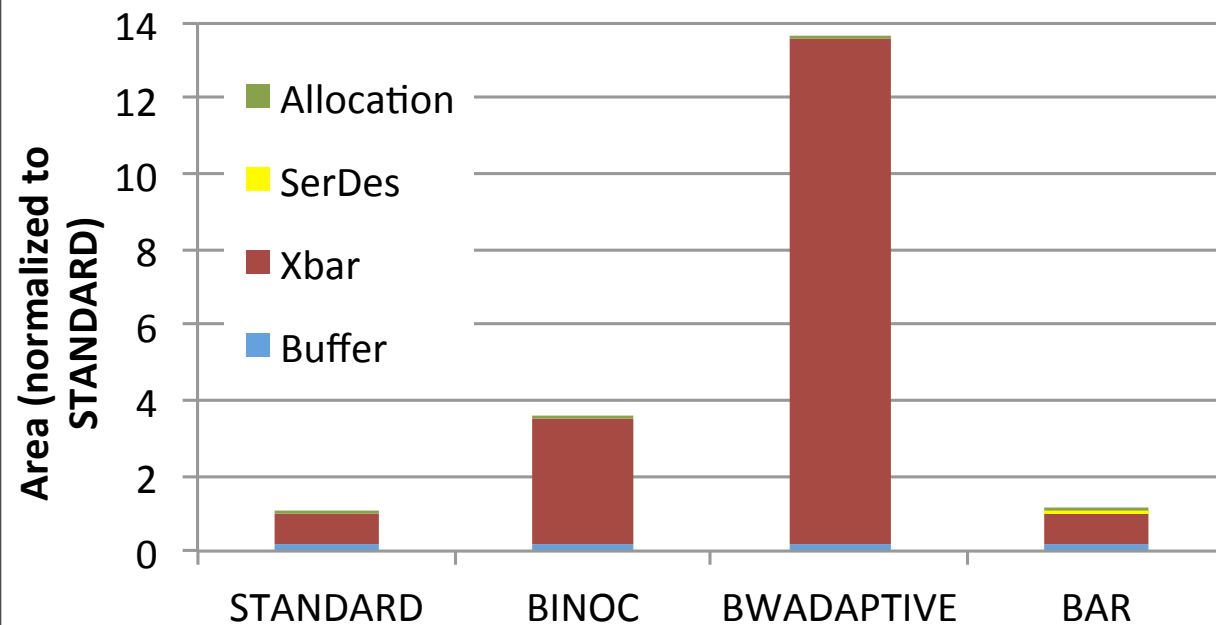
Architecture	STANDARD	BINOC	BWADAPT.	BAR
Total # of Buf.	5	10	20	10
Total Channels	5-in 5-out	10-inout	20-inout	20-inout
Each Buf. Size	32 flits	16 flits	8 flits	16 flits
Total Buf. Size	160 flits	160 flits	160 flits	160 flits
Crossbar	5x5	10x10	20x20	5x5
Flit width	8 byte	8 byte	8 byte	8 byte



Area Comparison

- Orion 2.0 with 45nm process

Initial Comparison



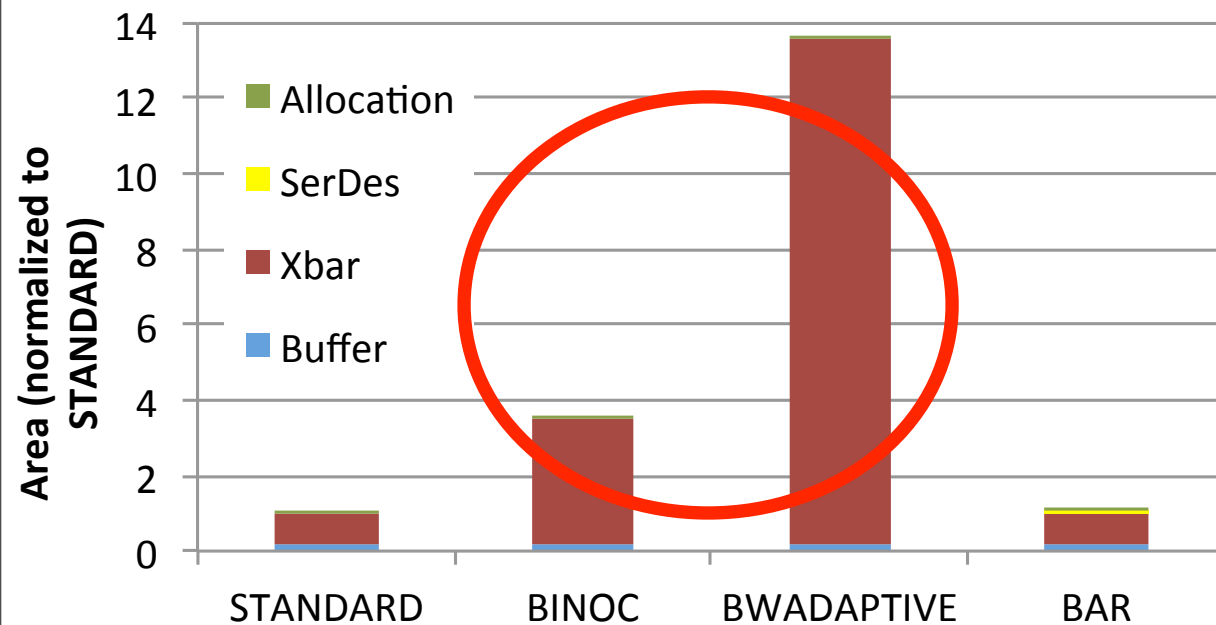
Architecture	STANDARD	BINOC	BWADAPT.	BAR
Total # of Buf.	5	10	20	10
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Area Comparison

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Initial Comparison



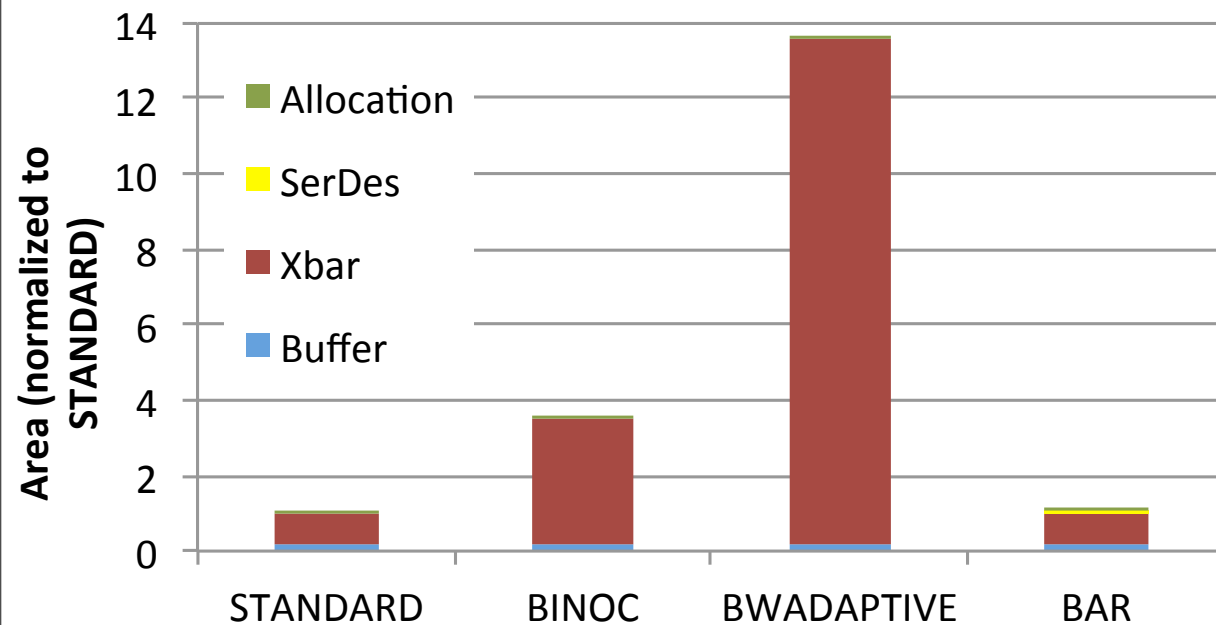
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Crossbar	5x5	10x10	20x20	5x5
Flit width	8 byte	8 byte	8 byte	8 byte



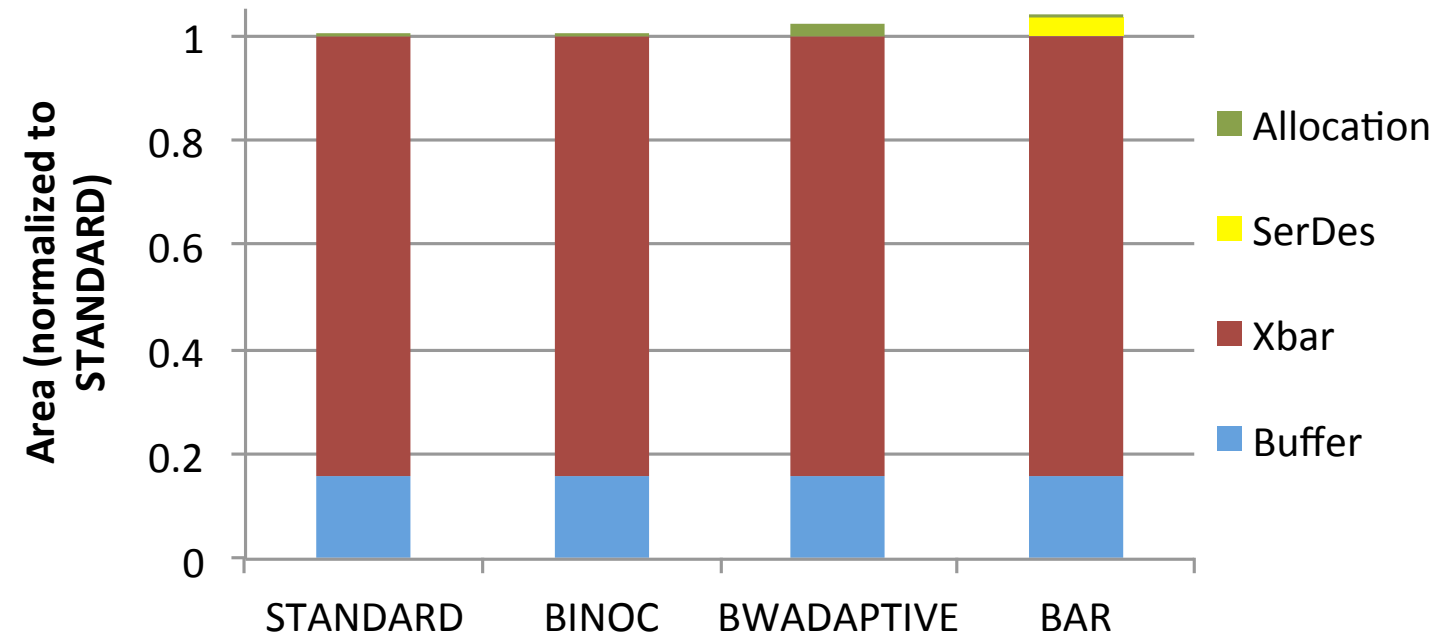
Area Comparison

- Orion 2.0 with 45nm process

Initial Comparison



Equalized Area

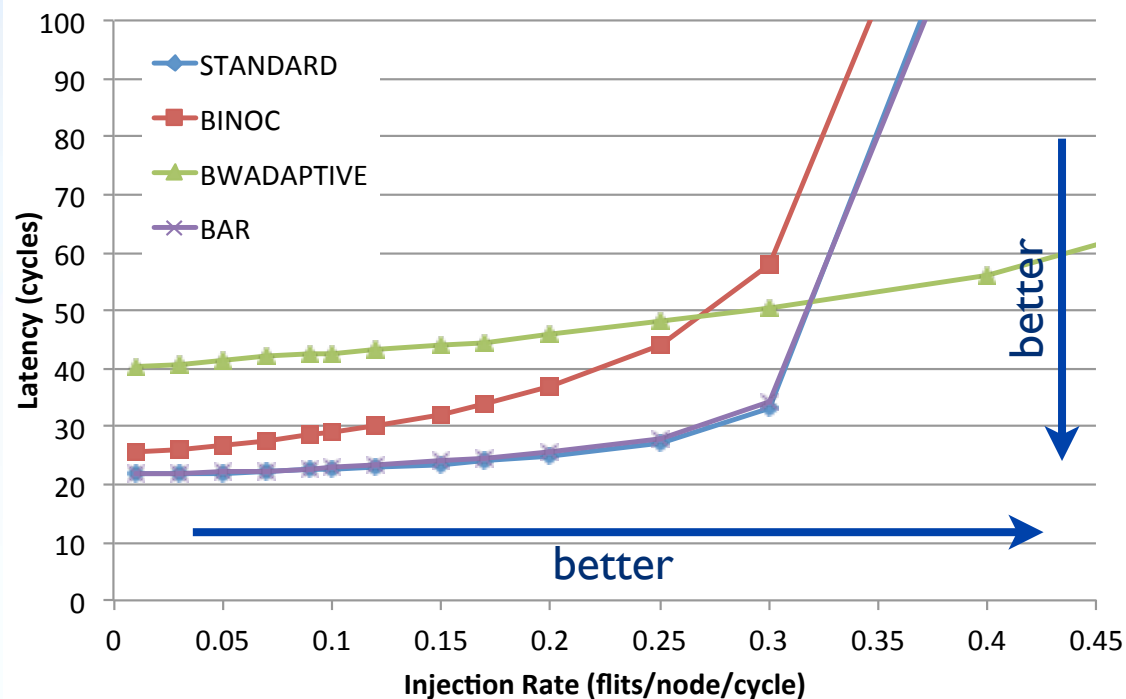


Architecture	STANDARD	BINOC	BWADAPT.	BAR
Total # of Buf.	5	10	20	10
Total Channels	5-in 5-out	10-inout	20-inout	20-inout
Each Buf. Size	32 flits	32 flits	32 flits	16 flits
Total Buf. Size	160 flits	160 flits	160 flits	160 flits
Crossbar	5x5	10x10	20x20	5x5
Flit width	8 byte	4 byte	2 byte	8 byte

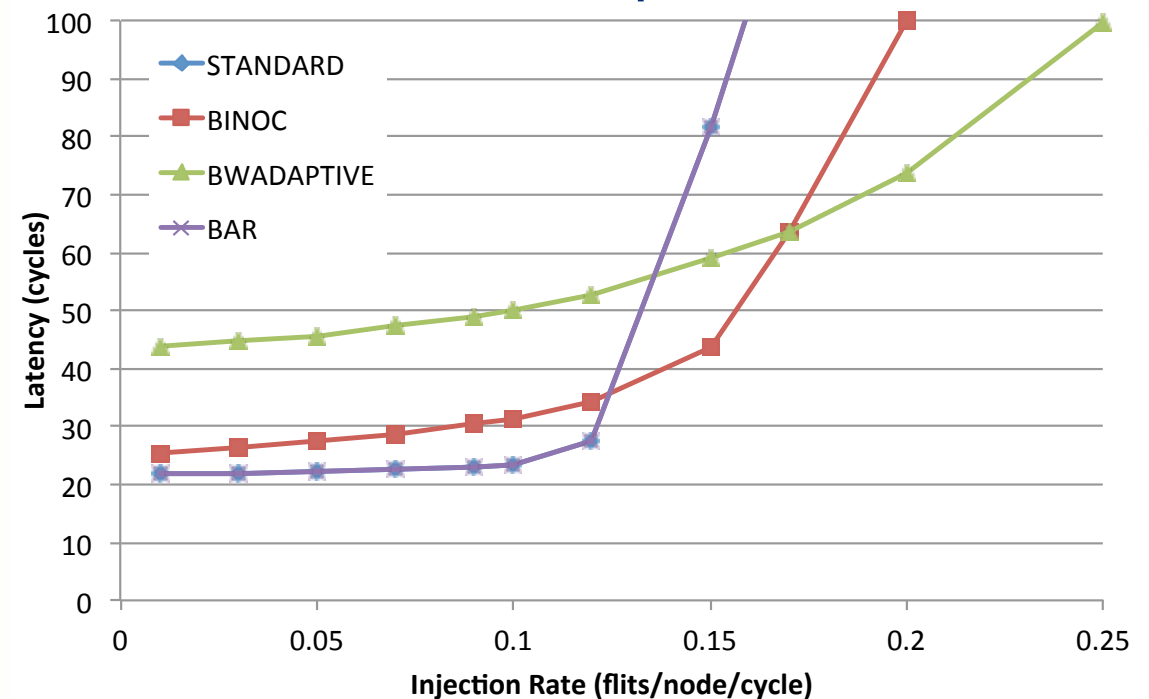


Static Network Performance

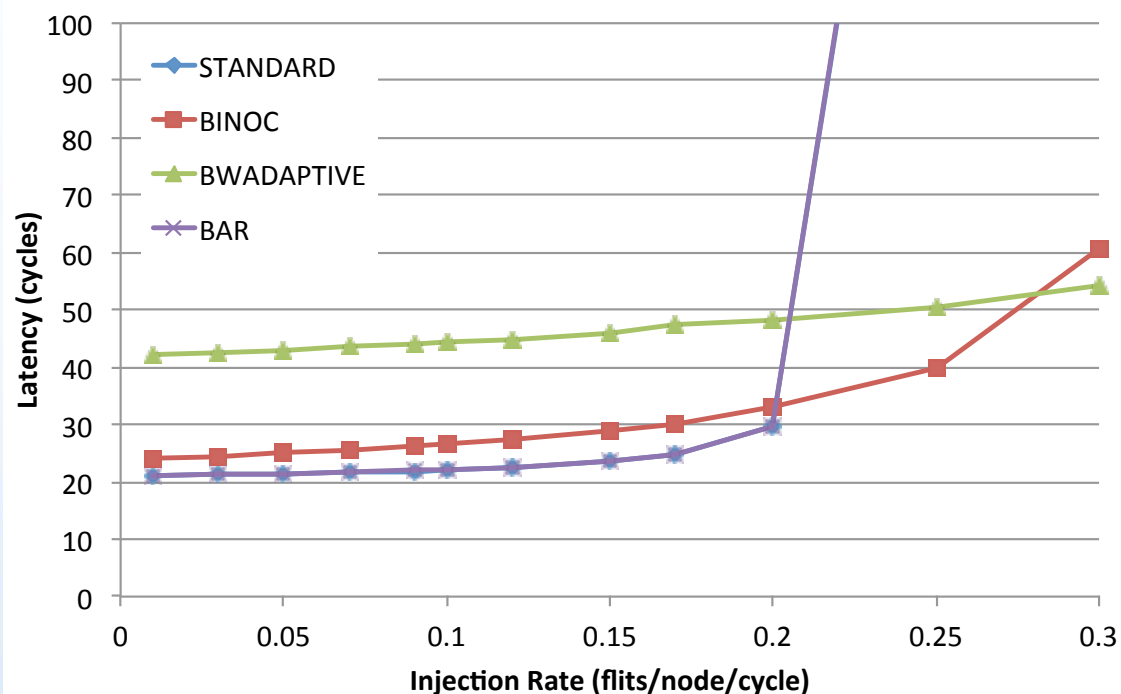
Uniform Random



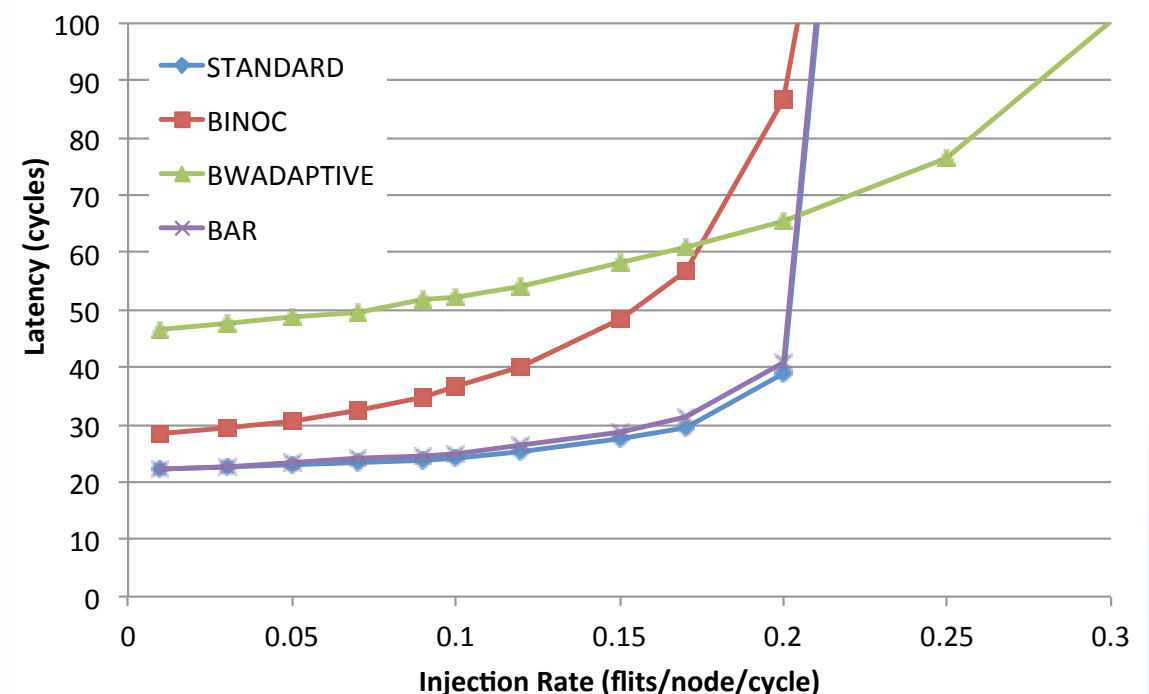
Transpose



Shuffle

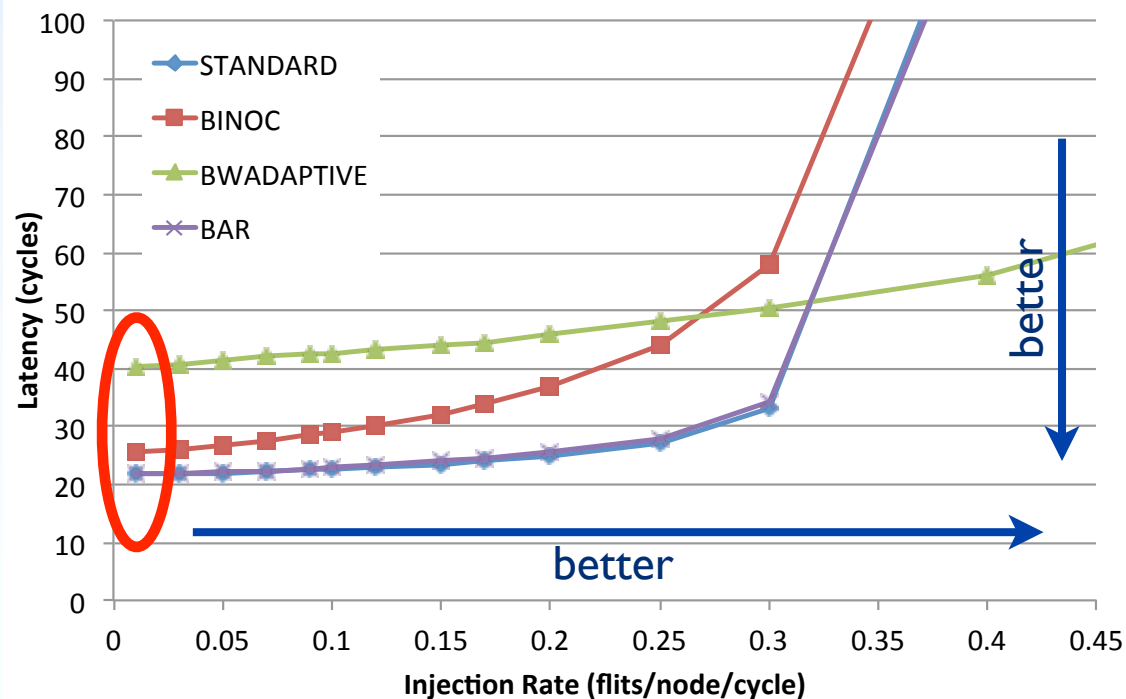


Bit-Complement

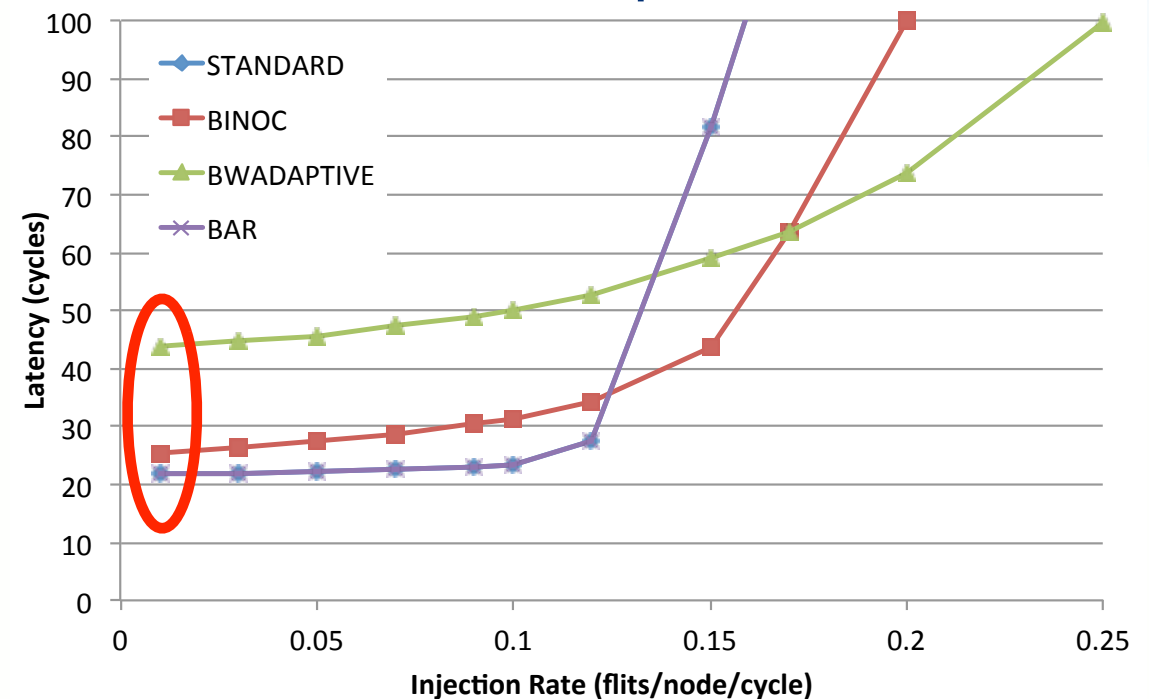


Static Network Performance

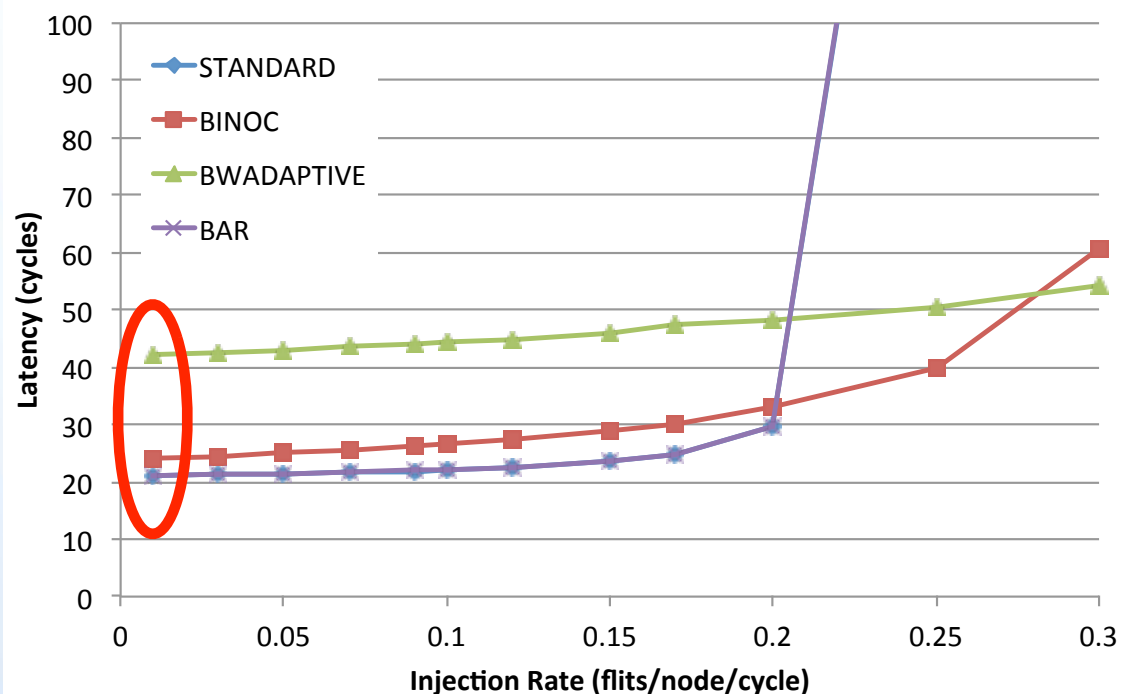
Uniform Random



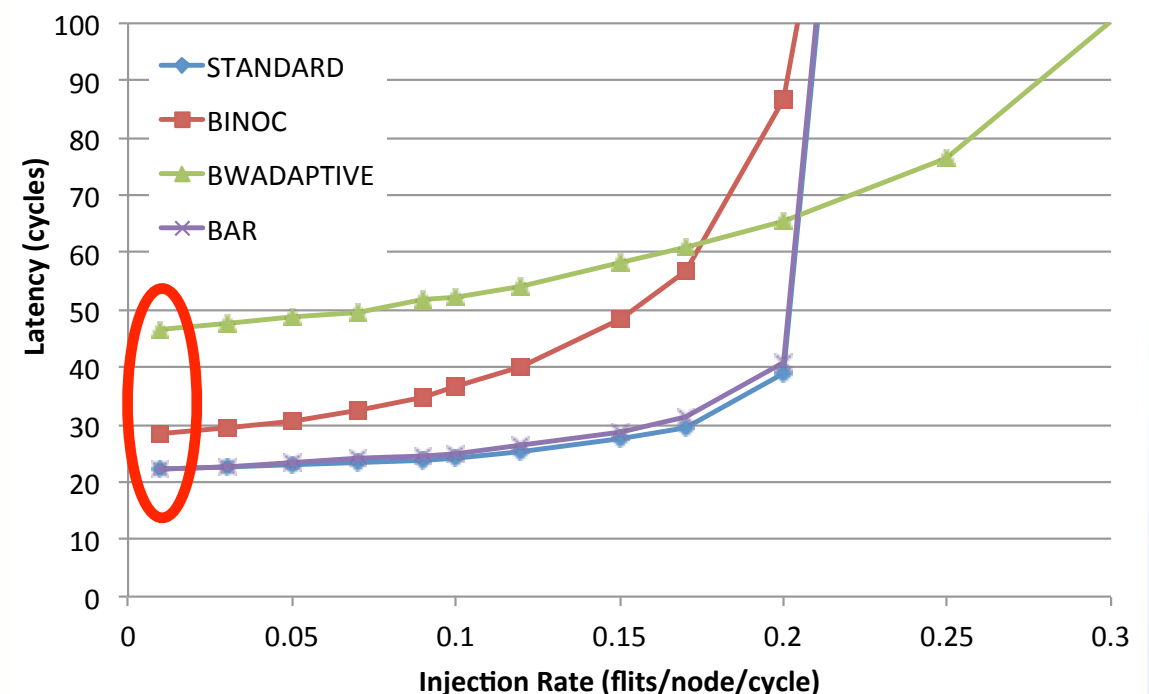
Transpose



Shuffle

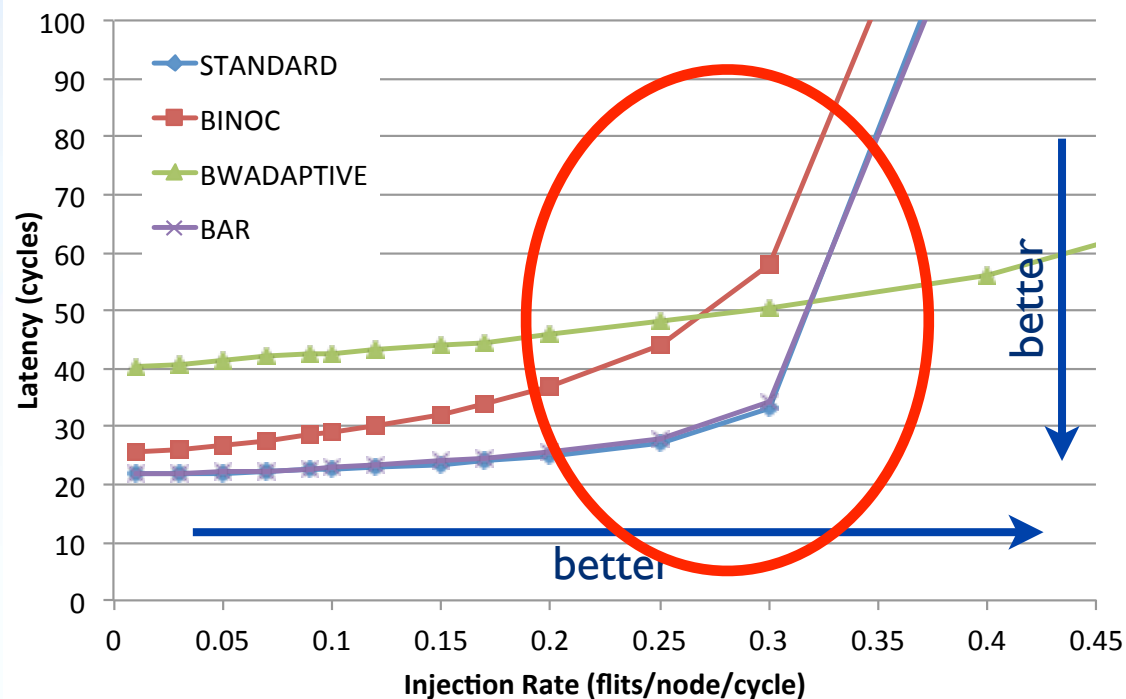


Bit-Complement

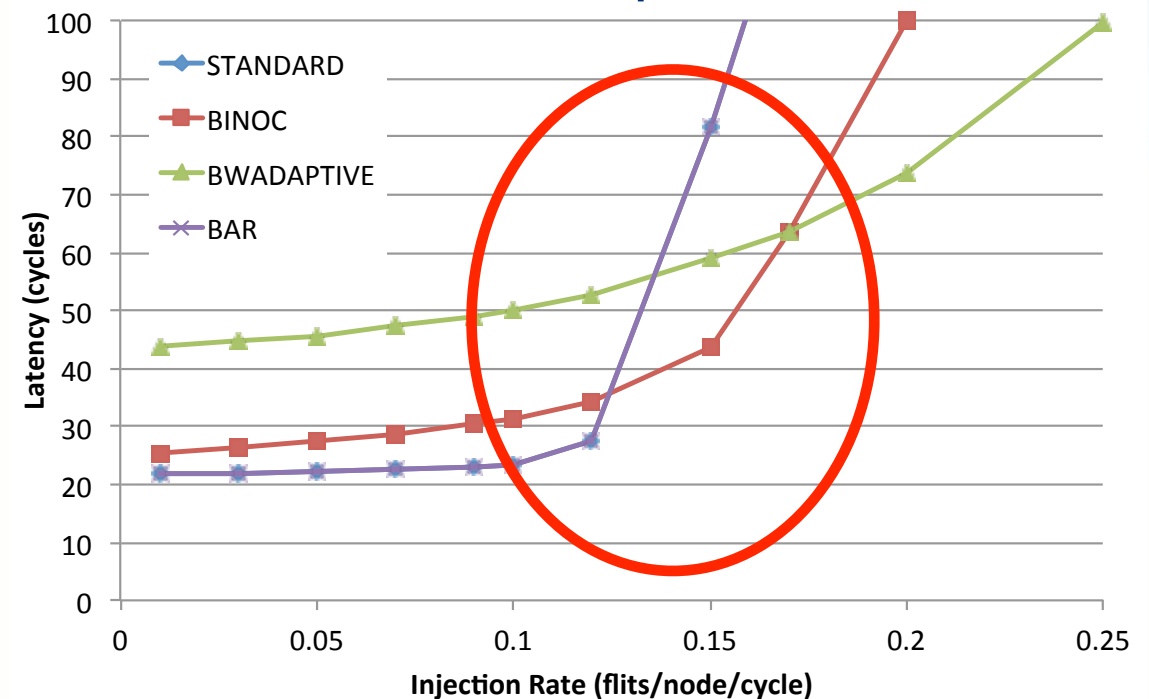


Static Network Performance

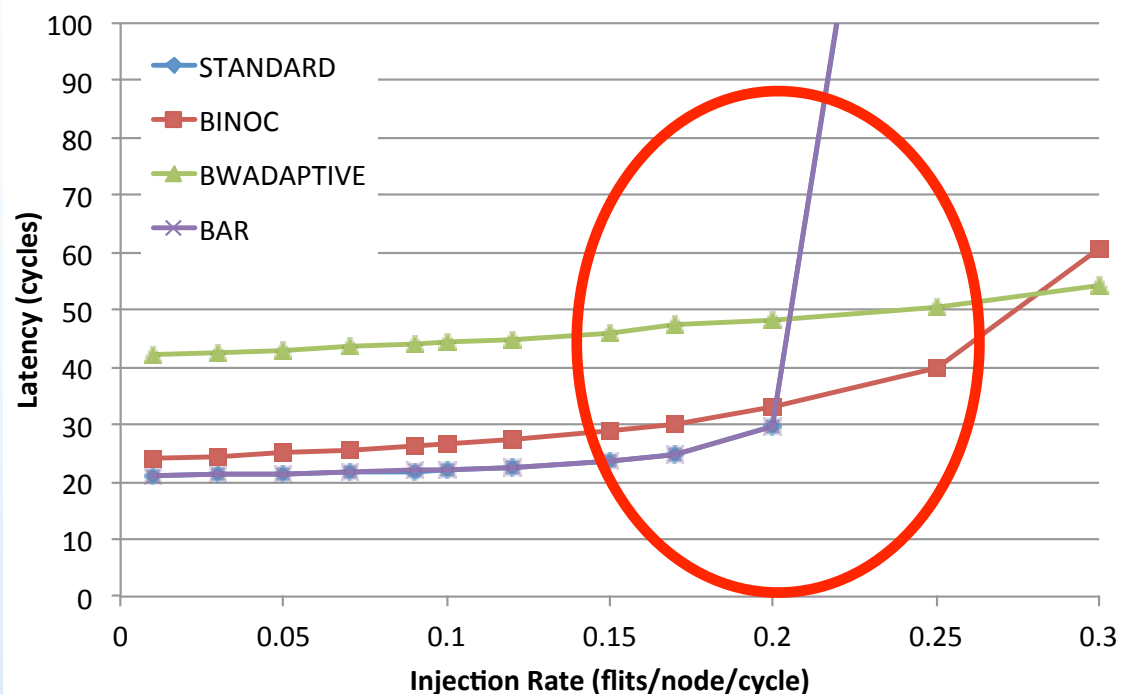
Uniform Random



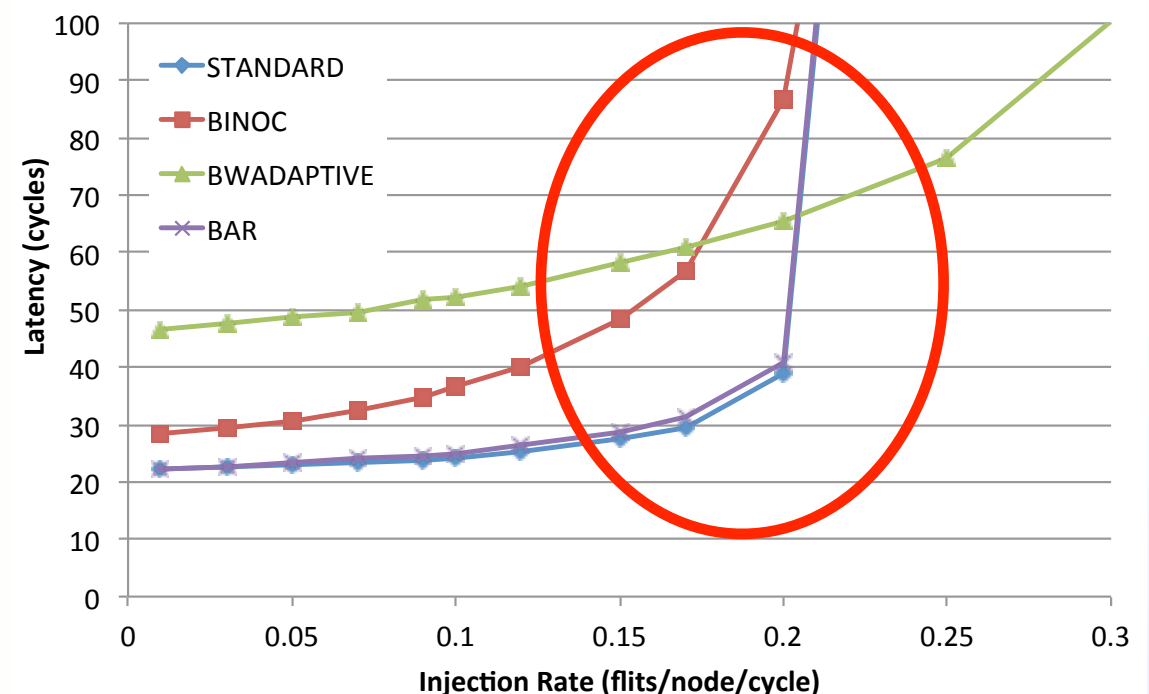
Transpose



Shuffle

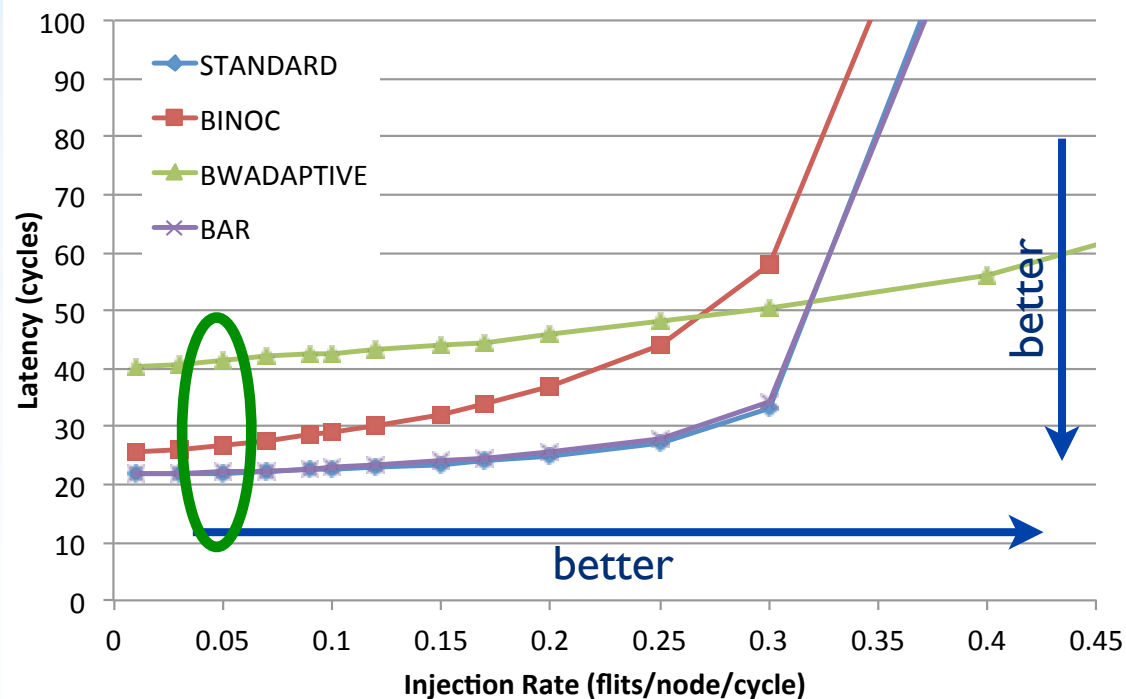


Bit-Complement

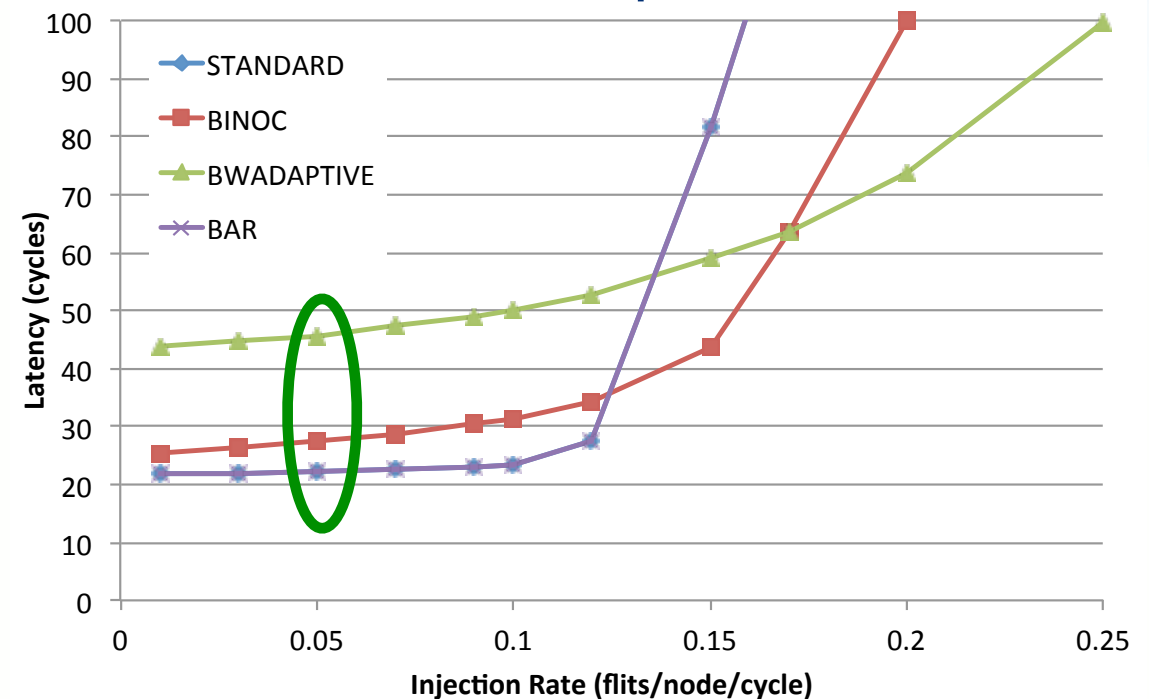


Static Network Performance

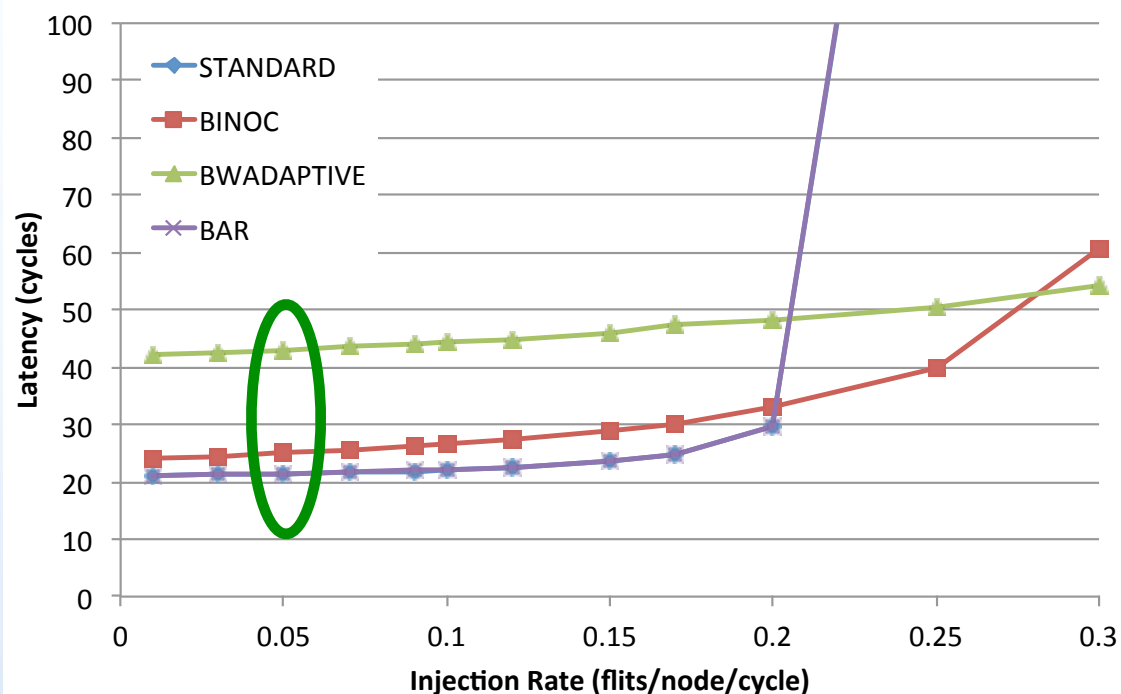
Uniform Random



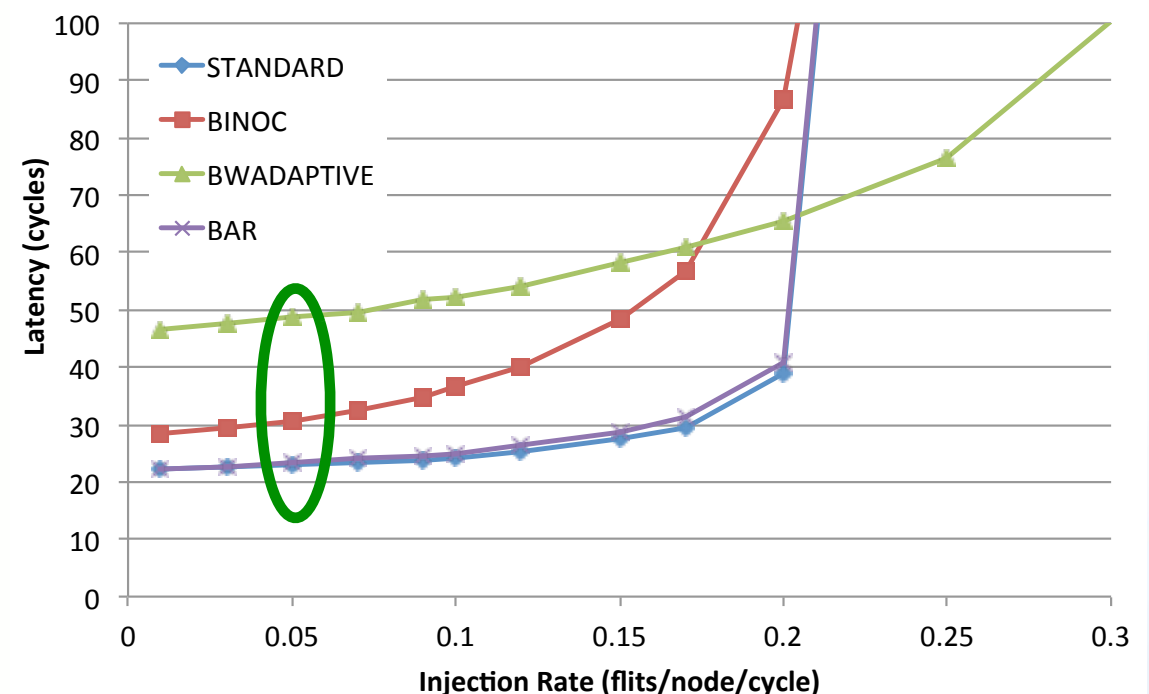
Transpose



Shuffle



Bit-Complement



Overall System Performance

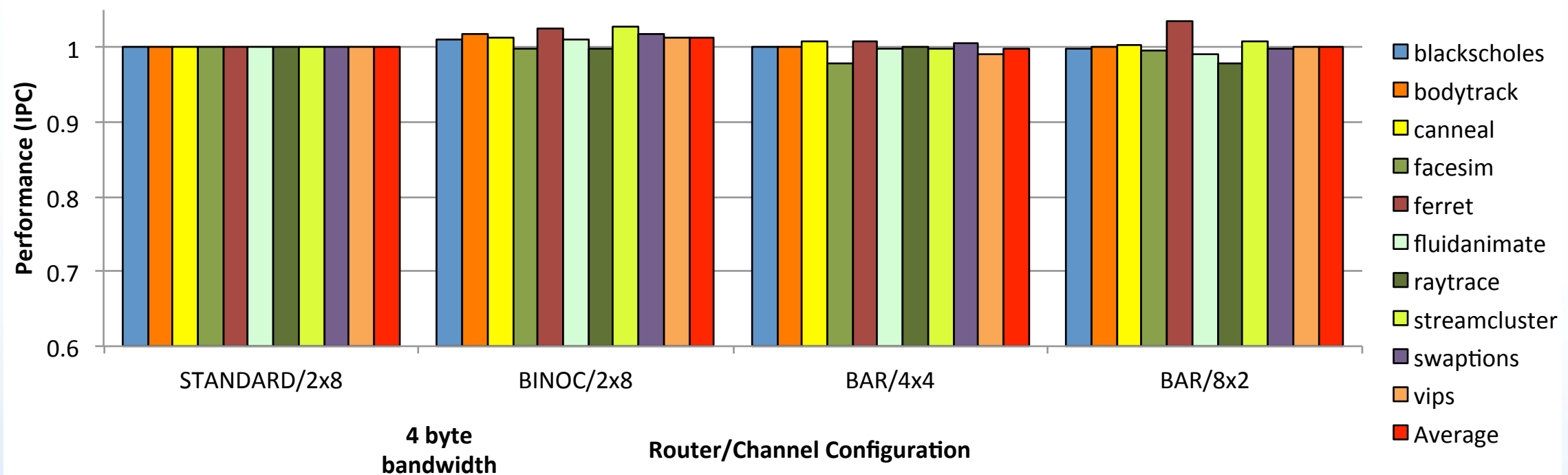
- Full-system simulation:
 - Cycle accurate x86 simulation: FeS2 + BookSim
 - PARSEC benchmarks (16 threads)
 - 16 P4-like CPUs, 4x4 mesh NoC



Overall System Performance

- Full-system simulation:
 - Cycle accurate x86 simulation: FeS2 + BookSim
 - PARSEC benchmarks (16 threads)
 - 16 P4-like CPUs, 4x4 mesh NoC

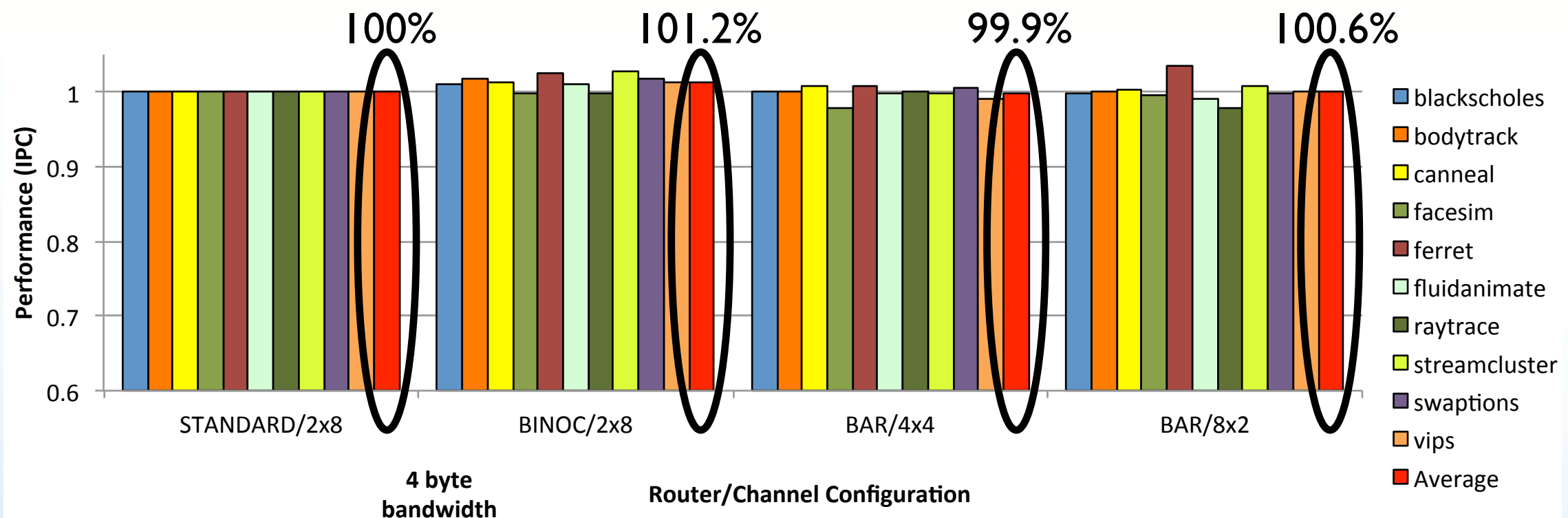
100% channel resources



Overall System Performance

- Full-system simulation:
 - Cycle accurate x86 simulation: FeS2 + BookSim
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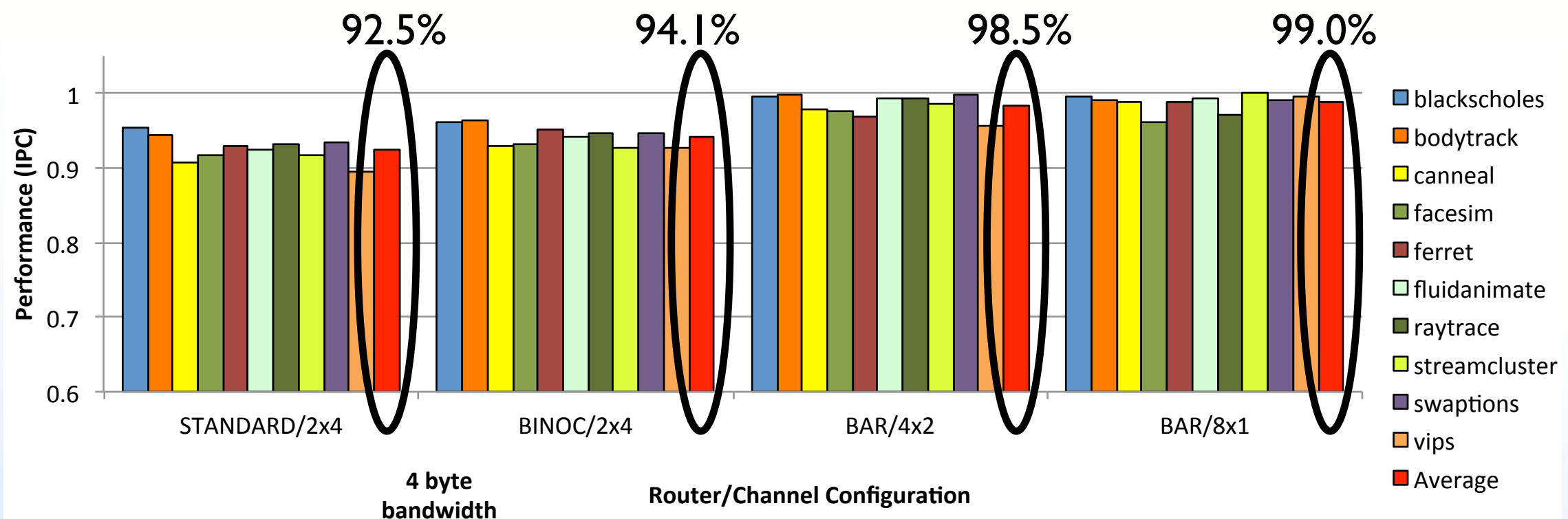
100% channel resources



Overall System Performance

- Full-system simulation:
 - Cycle accurate x86 simulation: FeS2 + BookSim
 - PARSEC benchmarks (16 threads)
 - 16 P4-like CPUs, 4x4 mesh NoC

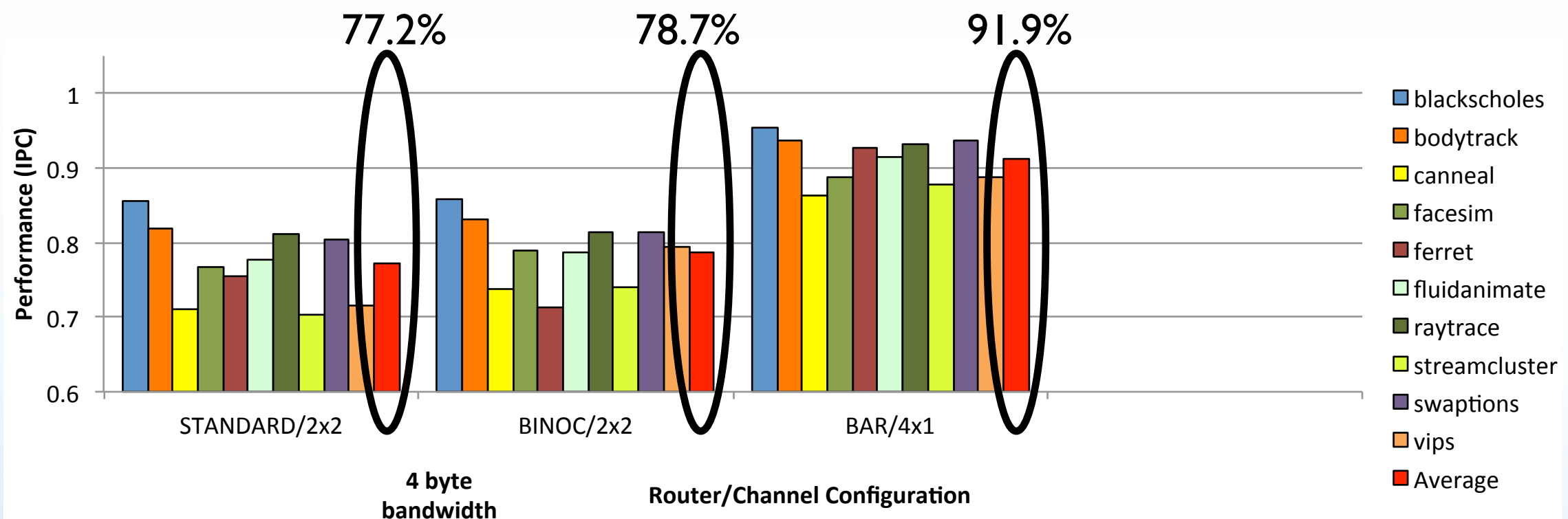
50% channel resources



Overall System Performance

- Full-system simulation:
 - Cycle accurate x86 simulation: FeS2 + BookSim
 - PARSEC benchmarks (16 threads)
 - 16 P4-like CPUs, 4x4 mesh NoC

25% channel resources



Summary



Summary

- We introduce fine-grained BW adaptivity



Summary

- We introduce fine-grained BW adaptivity
- Phit-serial communication to decouple flit width from channel width



Summary

- We introduce fine-grained BW adaptivity
- Phit-serial communication to decouple flit width from channel width
- Improved channel utilization without significantly increasing latency



Summary

- We introduce fine-grained BW adaptivity
- Phit-serial communication to decouple flit width from channel width
- Improved channel utilization without significantly increasing latency
- 50% reduction channel resources: 99% perf.



Summary

- We introduce fine-grained BW adaptivity
- Phit-serial communication to decouple flit width from channel width
- Improved channel utilization without significantly increasing latency
- 50% reduction channel resources: 99% perf.
- 75% reduction in channel resources: 92% perf.

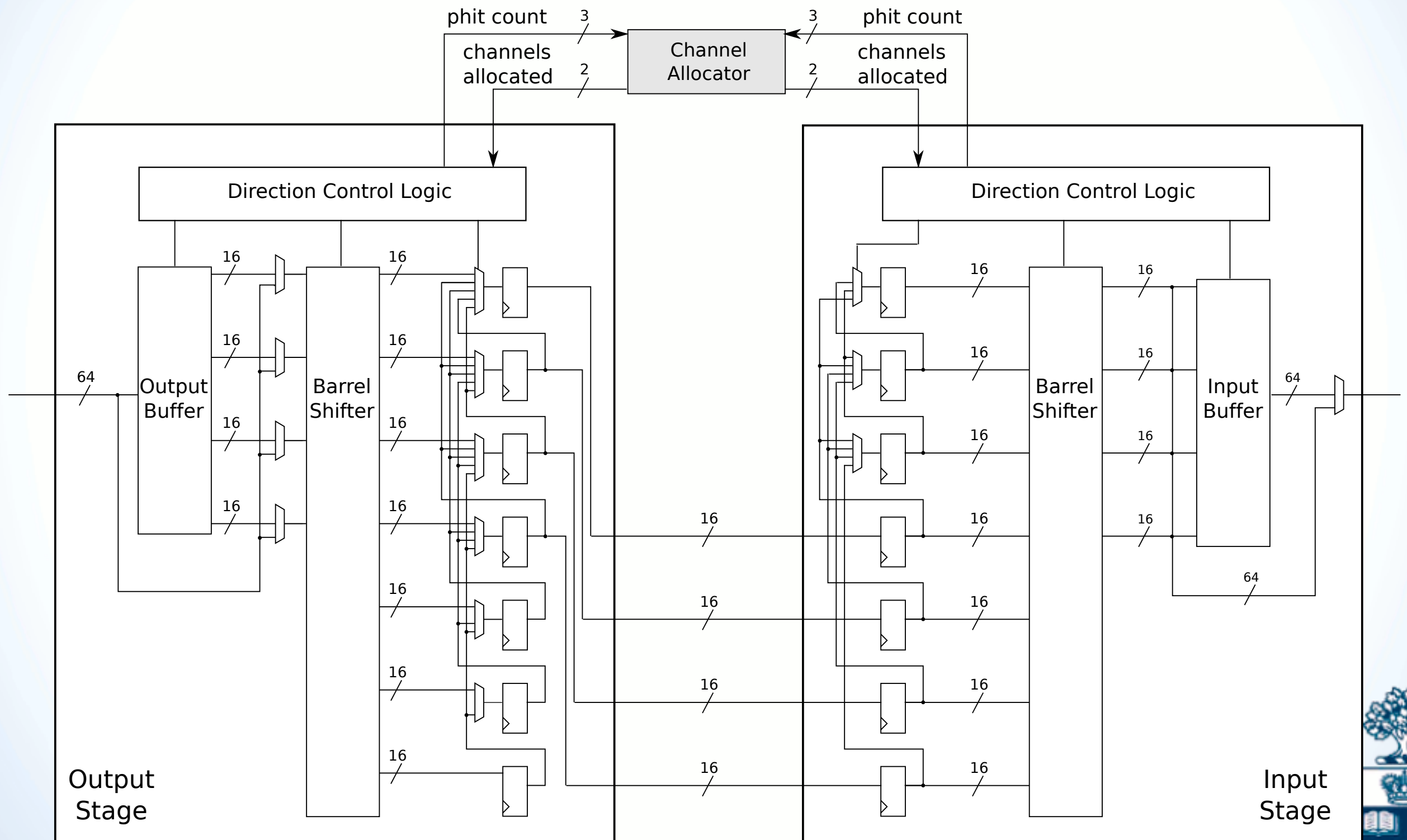


Thank you!

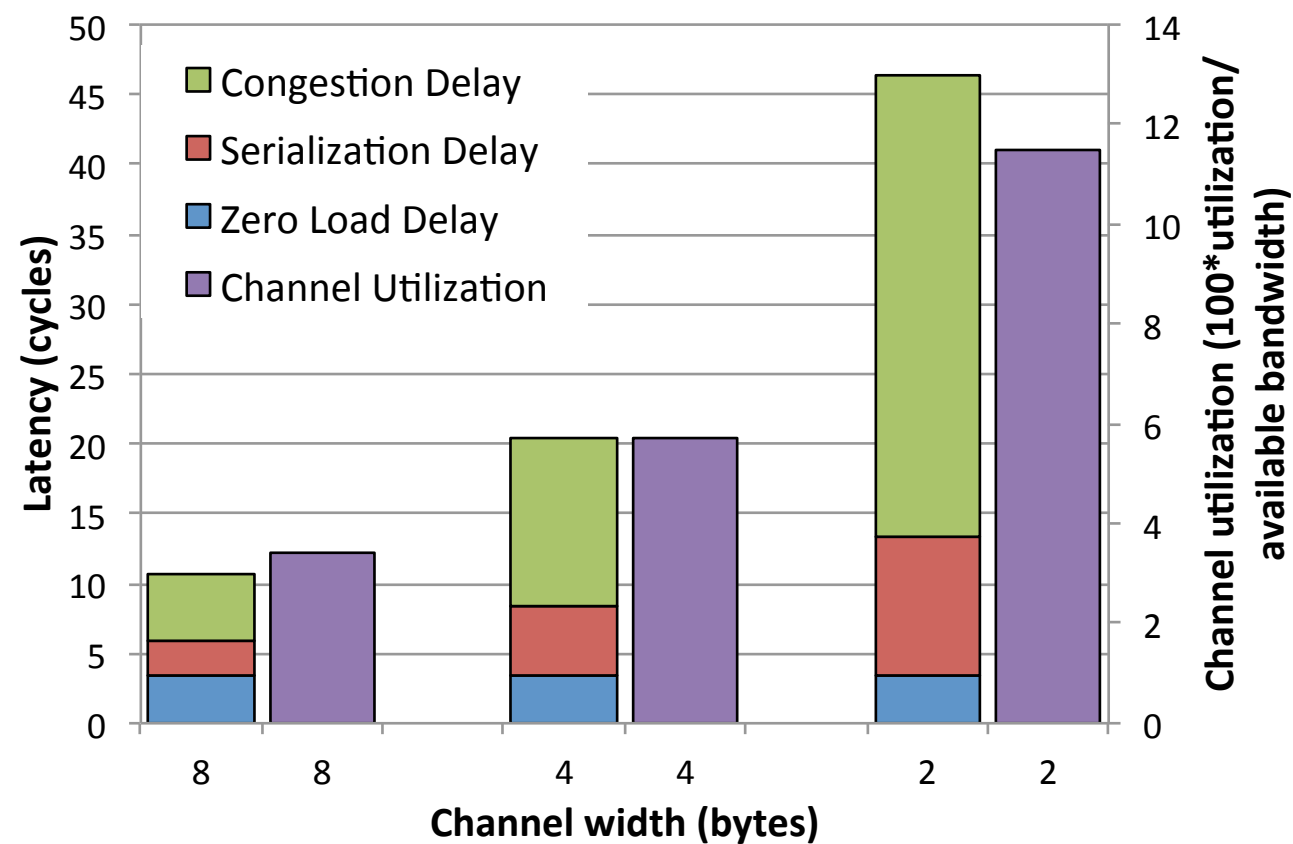
robert.hesse@utoronto.ca



Hardware Implementation



Additional Delays



Implementation Results

Router	Area (μm^2)	Power (mW)	Crit. path (ns)	Wires
STANDARD	316194	17.31	1.1	640
BAR	319358	17.63	1.1	345

