

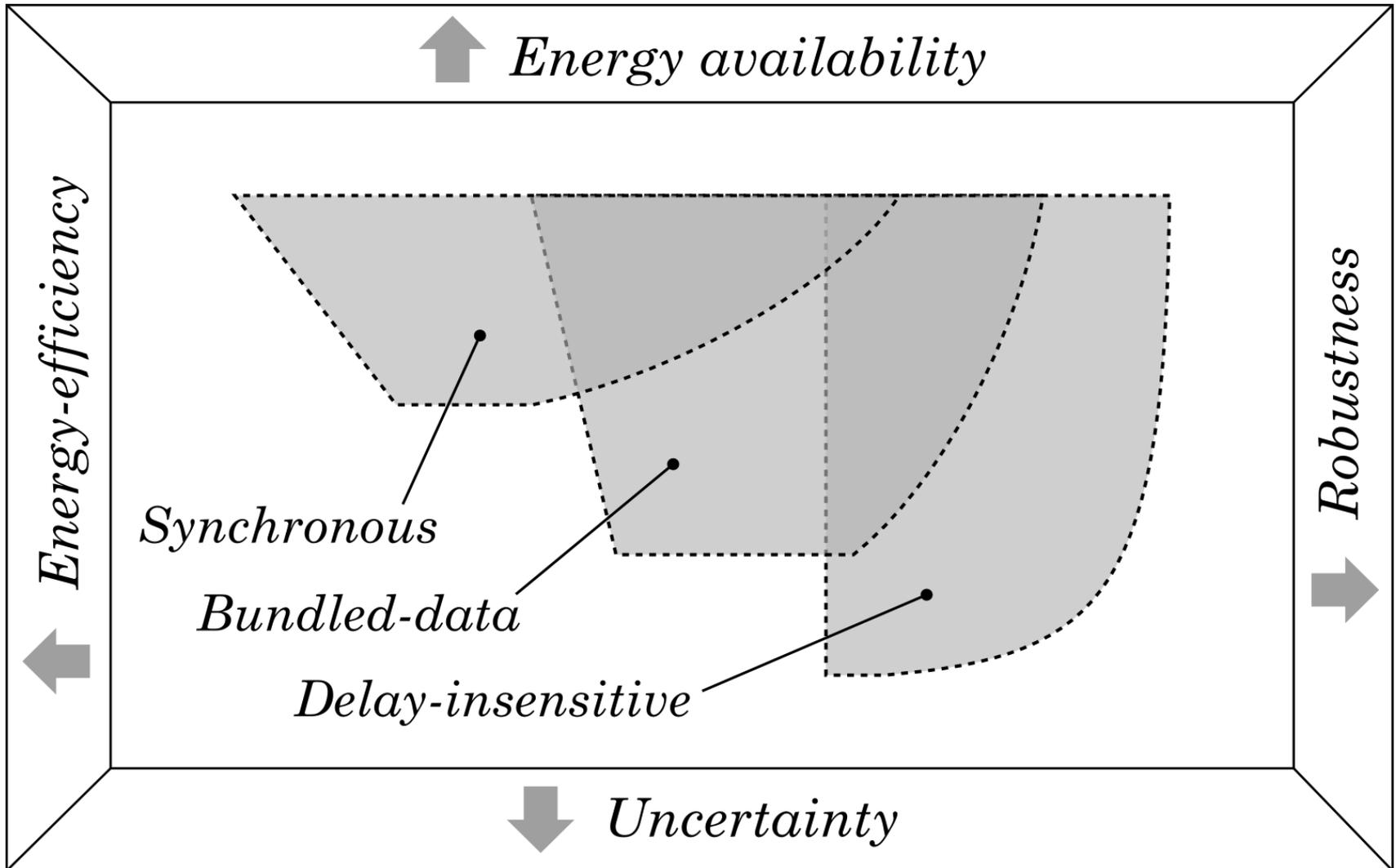
Adapting Asynchronous Circuits to Operating Conditions by Logic Parameterisation

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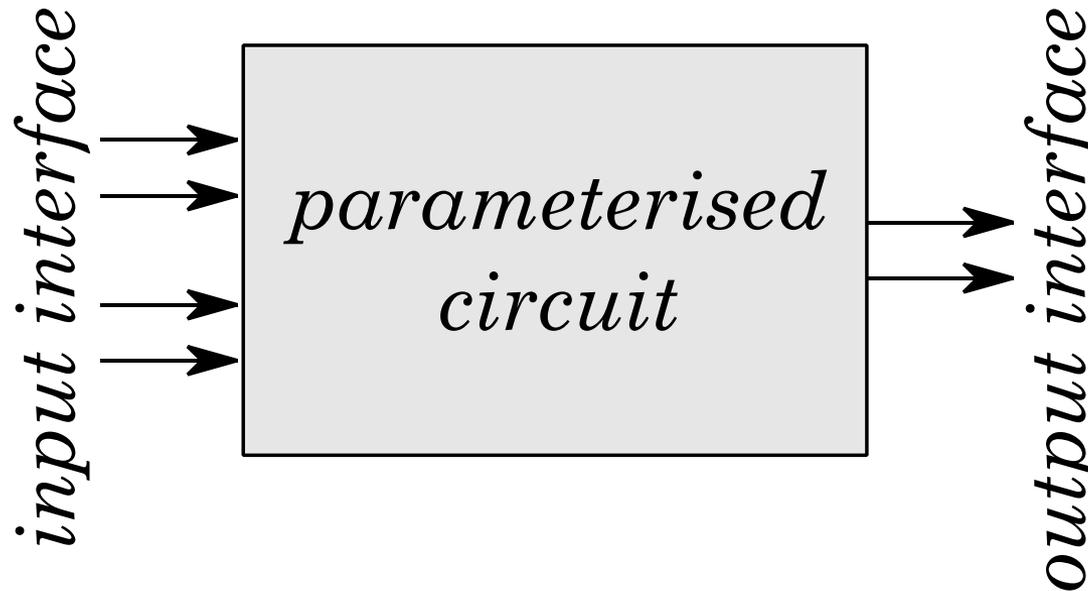
Designer's dilemma

- Contradicting requirements:
 - \uparrow performance, \downarrow energy consumption, \uparrow robustness
 - high performance \rightarrow high energy consumption
 - low energy consumption \rightarrow low robustness
 - high robustness \rightarrow low performance
- Competing implementation styles:
 - Synchronous
 - Asynchronous
 - Delay Insensitive (DI)
 - Speed Independent (SI) or Quasi Delay Insensitive (QDI)
 - Bundled data, relative timing assumptions (TA)
 - ...

Energy-efficiency v Robustness

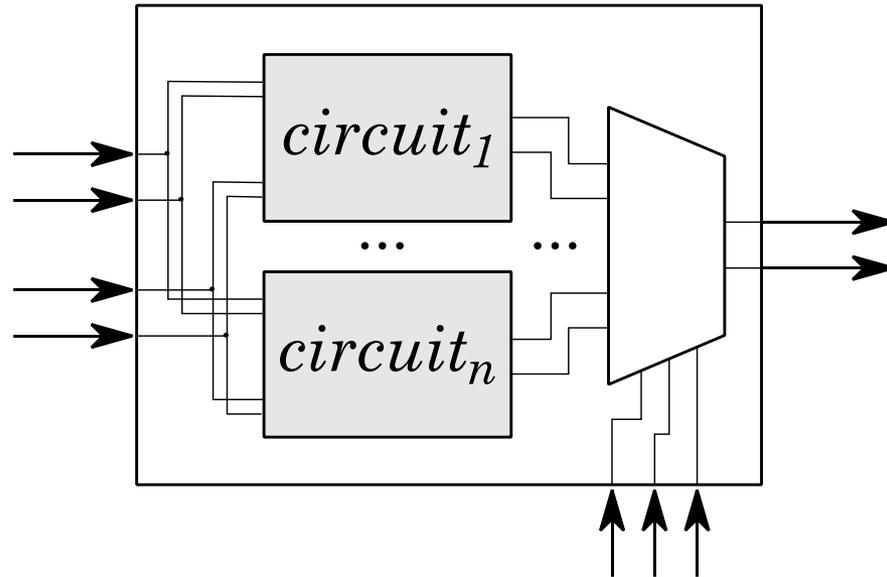


Parameterised Circuits



- Static parameters set at test/binning stages
- Dynamic parameters:
 - Power management controller
 - Maintenance mechanisms

Parameterised Circuits: Trivial Approach

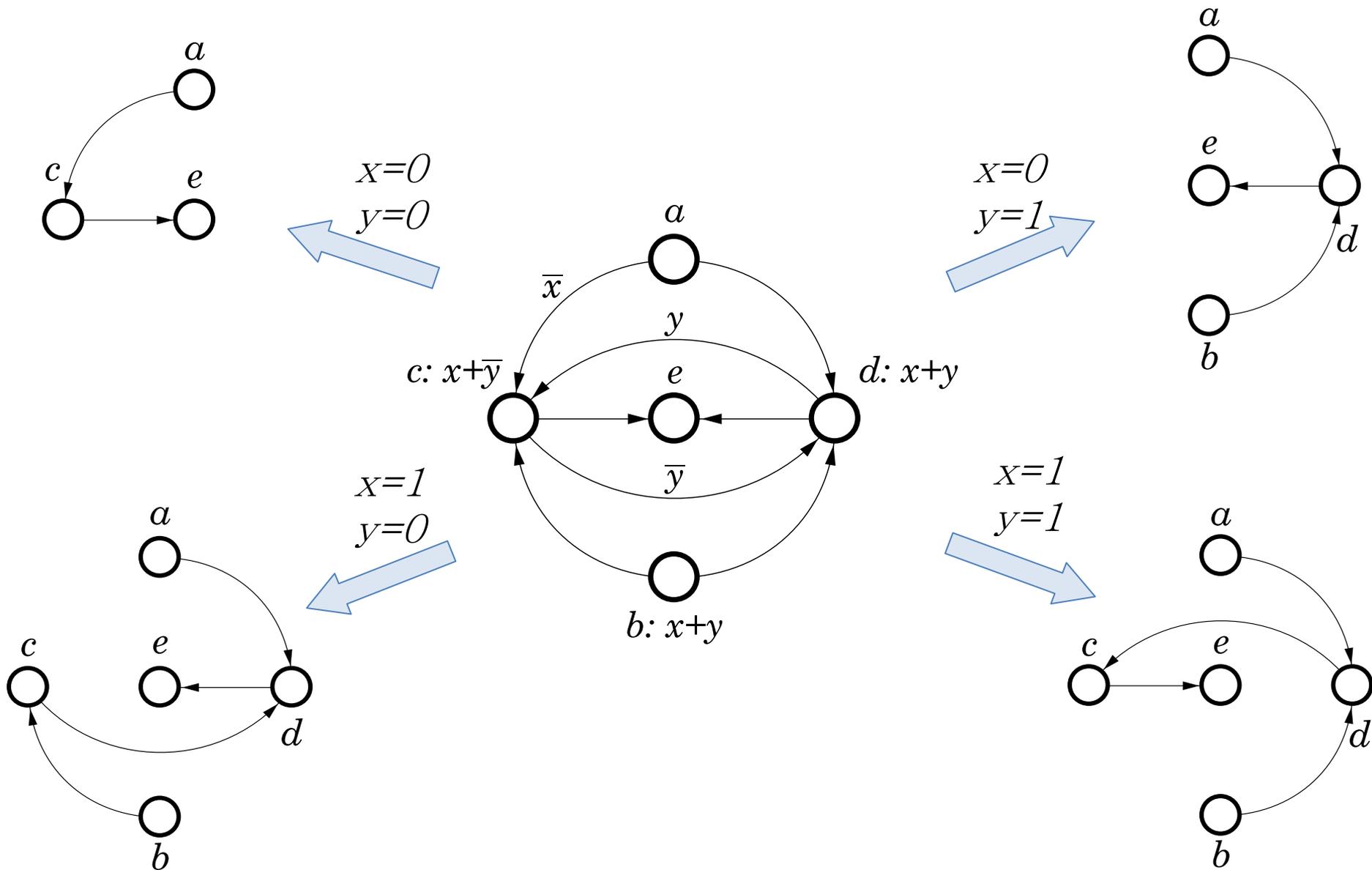


- Easy to design!
- Large overheads
- How can we do better?

Parameterised Circuits: Better Approach

- Goal:
 - Combine circuit implementations **efficiently**
- Key observations:
 - Externally: the circuits have the same interface
 - Internally: the circuits behave **similarly**
- Solution: use a model that can capture functional similarities at the circuit level
 - Conditional Partial Order Graphs **almost** fit

Conditional Partial Order Graphs

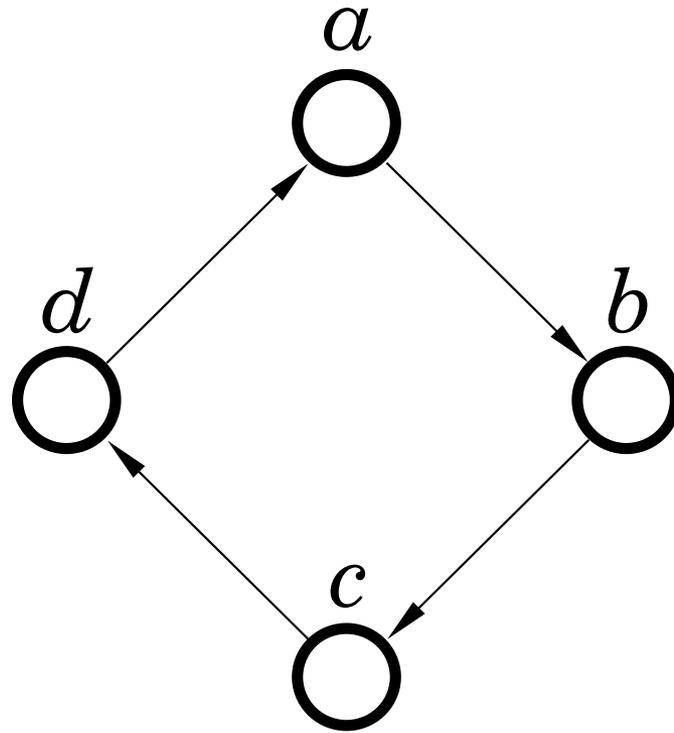


Conditional Partial Order Graphs

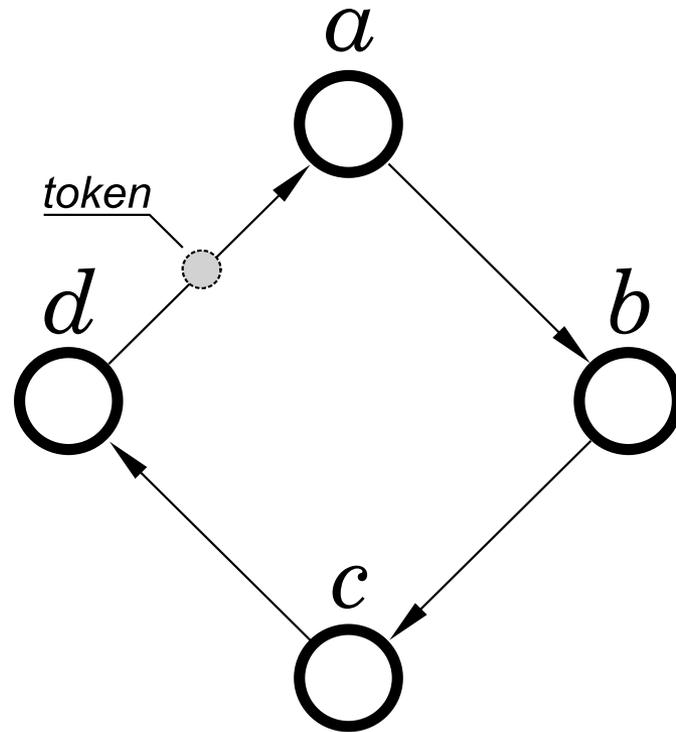
- 👉 Describe concurrency and causality
- 👉 Capture similarities in behaviours
- 👉 Represent families of **partial orders**
 - acyclic
 - not directly applicable to circuit specification
- 👉 Is it possible to specify cyclic behaviour using acyclic objects?

—Yes!

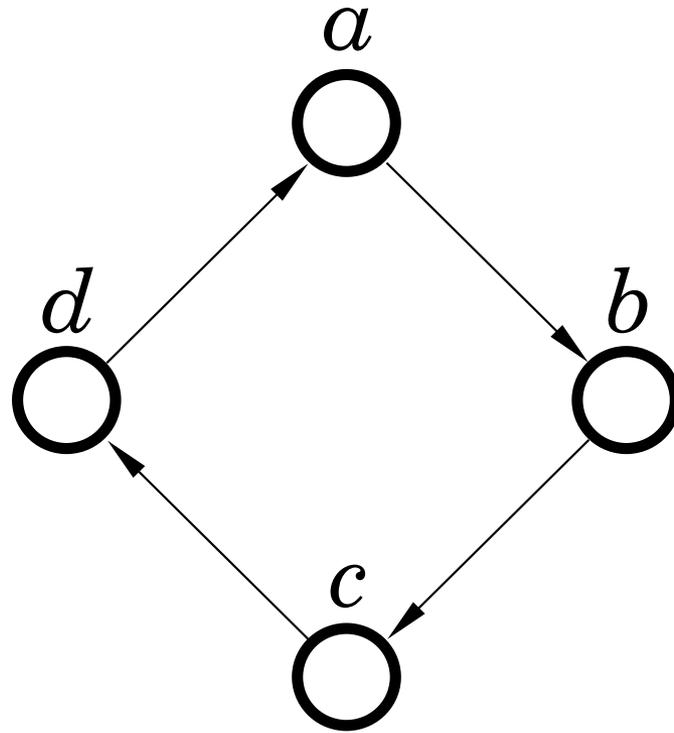
Describing cycles



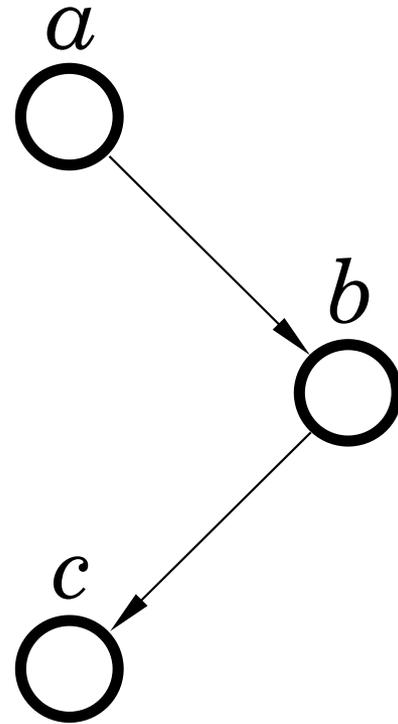
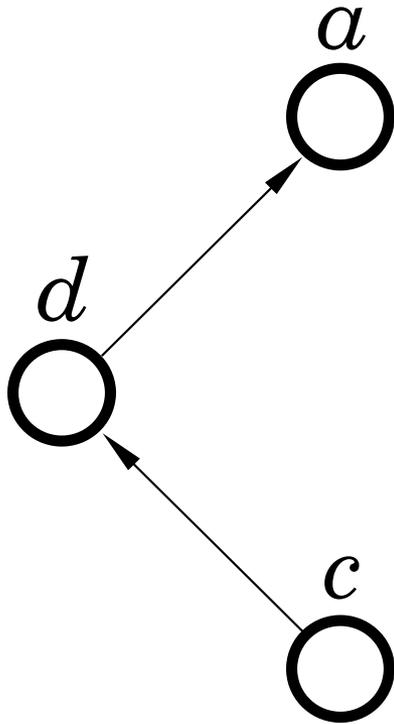
Describing cycles



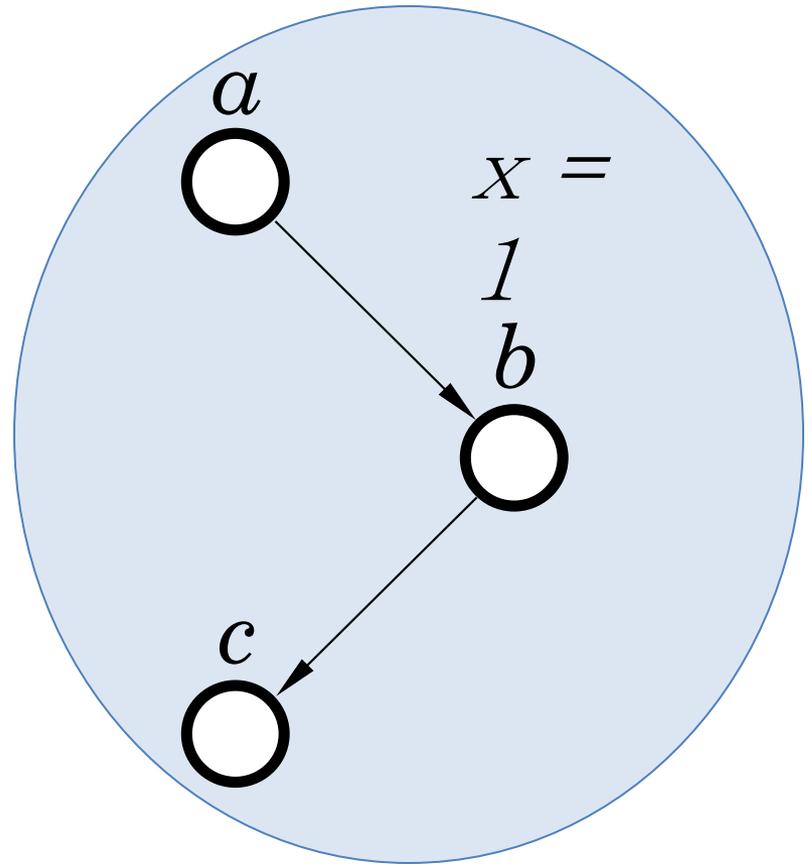
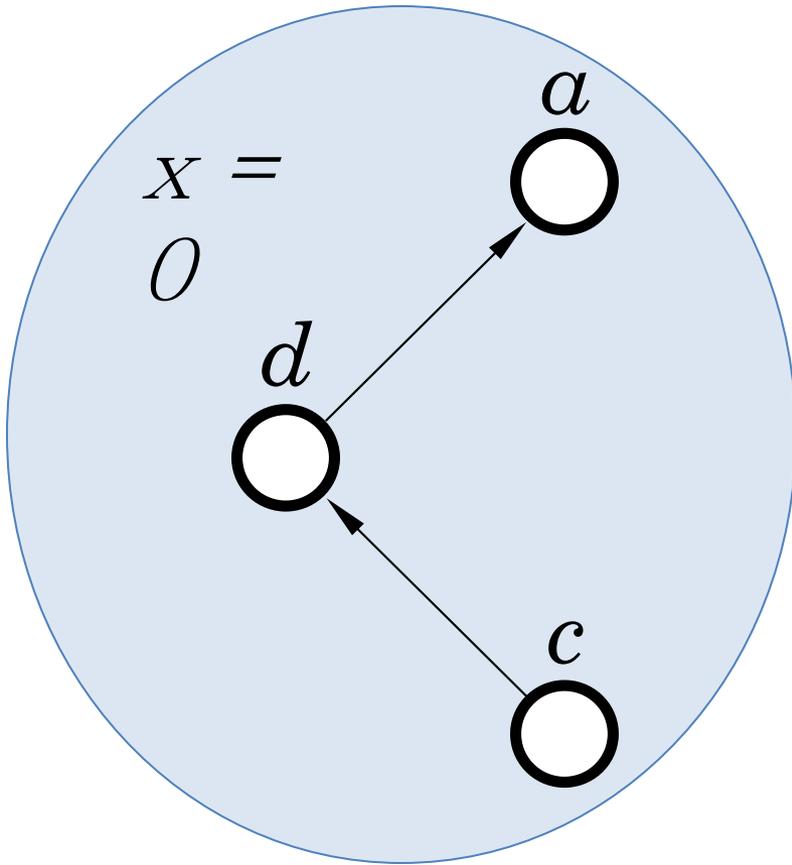
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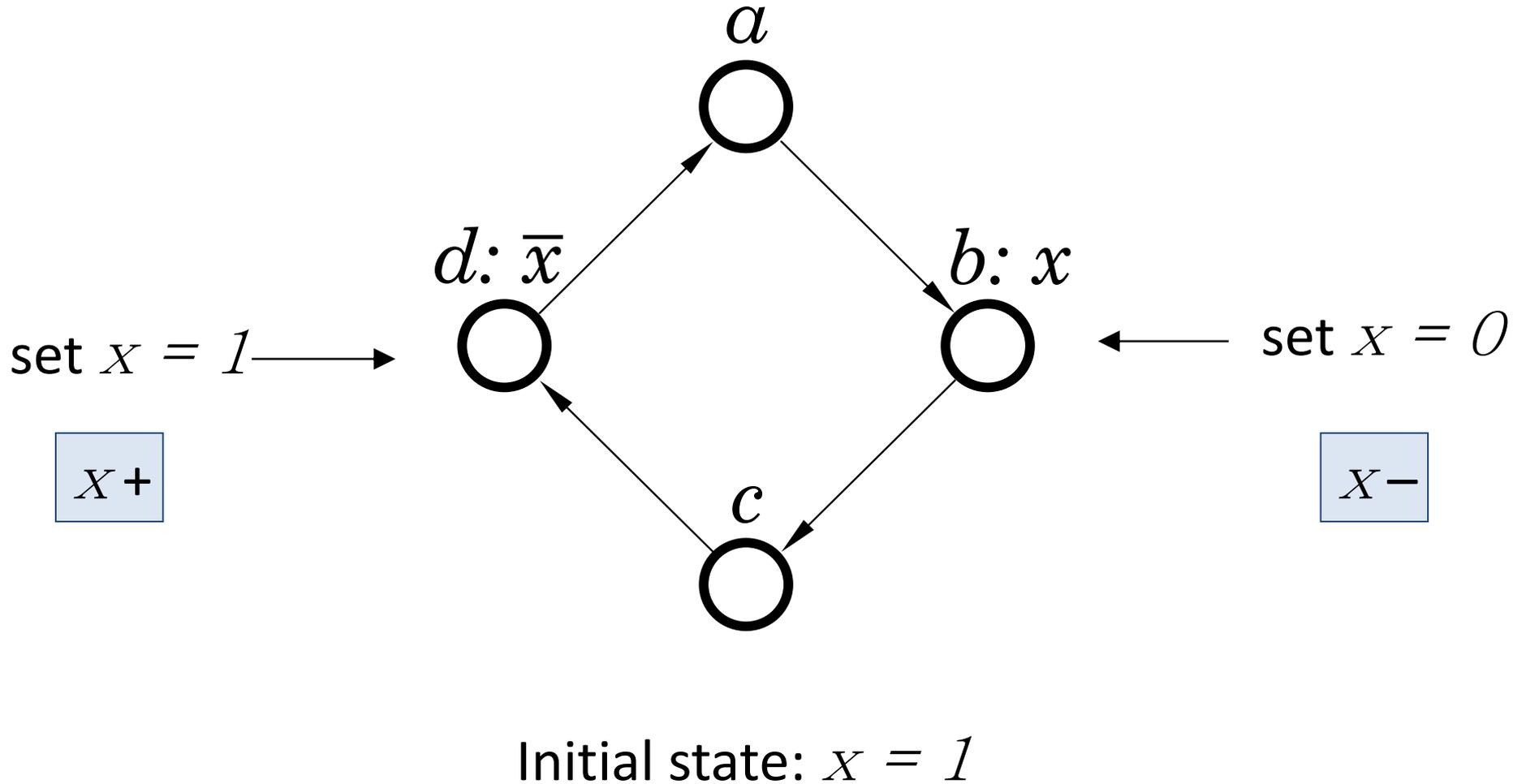
Describing cycles



Describing cycles

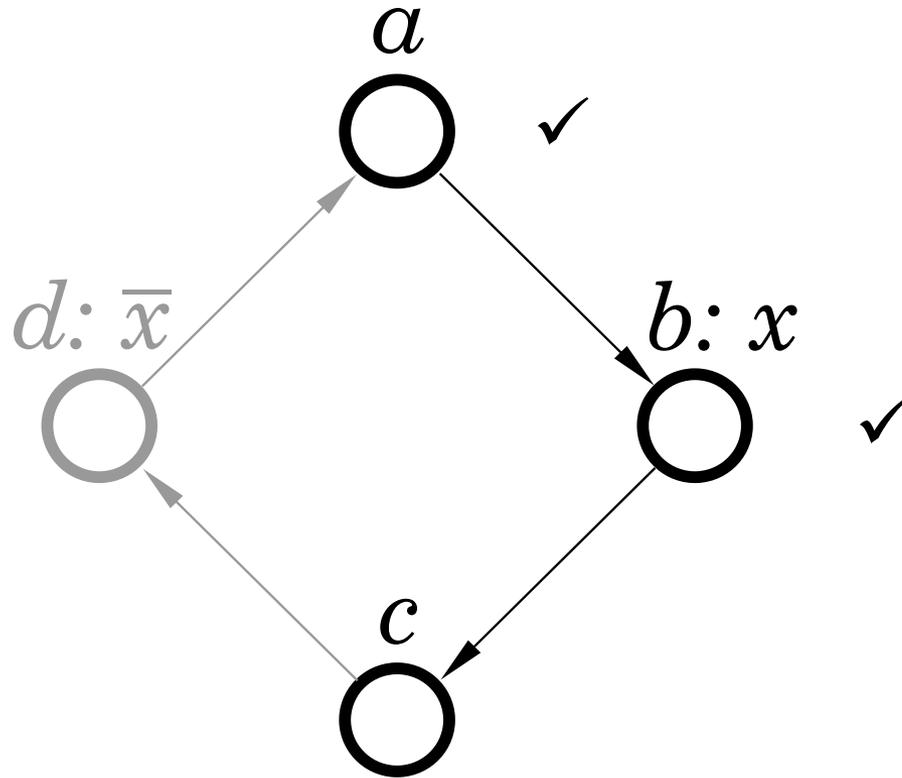


Describing cycles



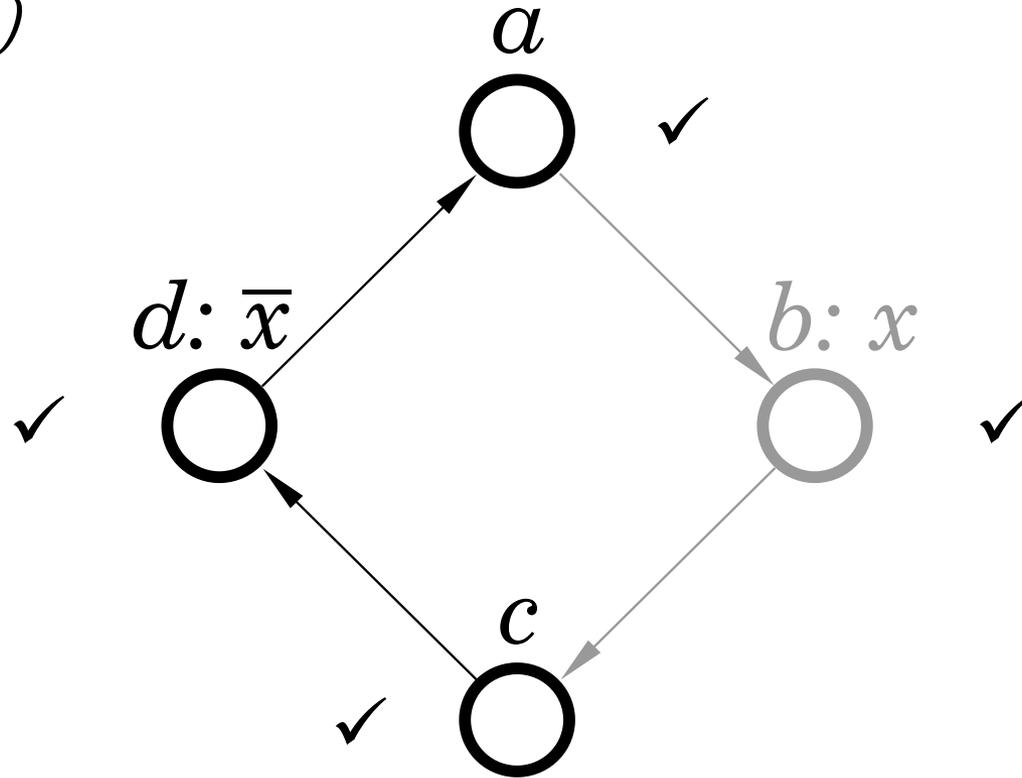
Describing cycles

$$x = 1$$



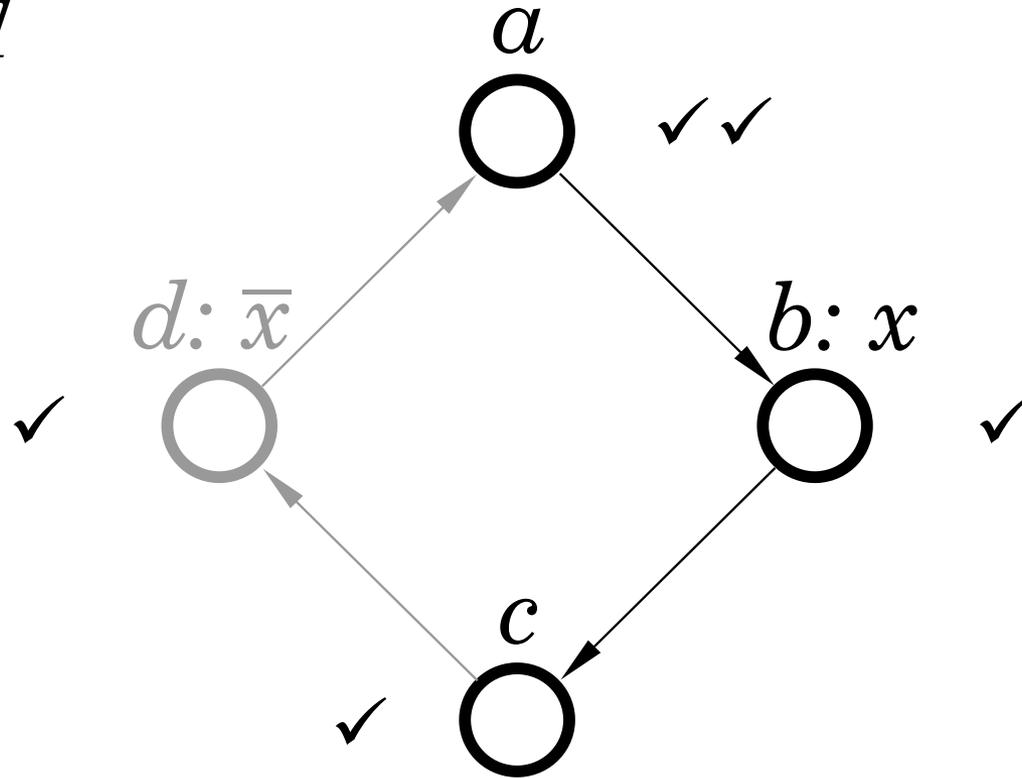
Describing cycles

$$x = 0$$

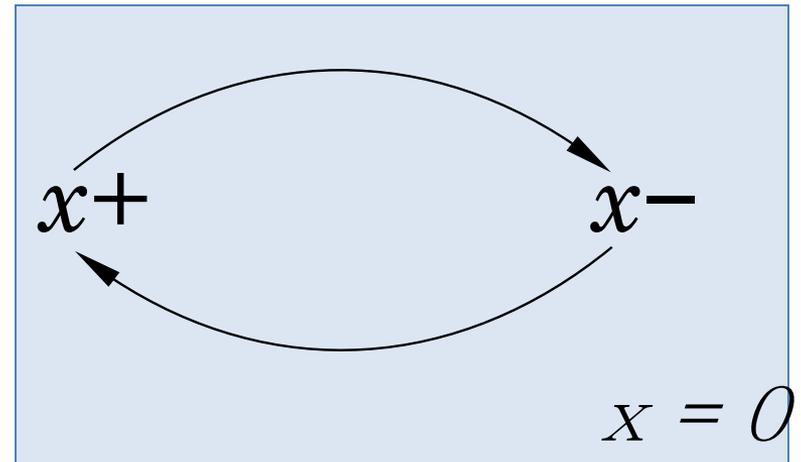
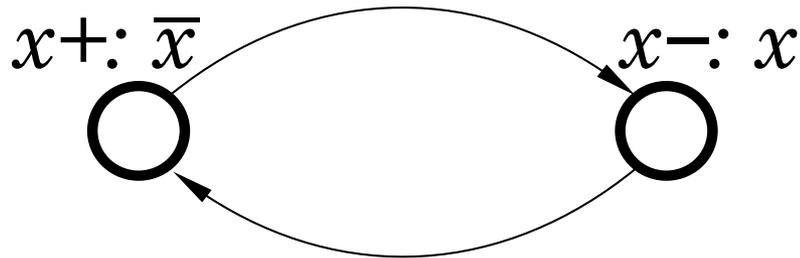
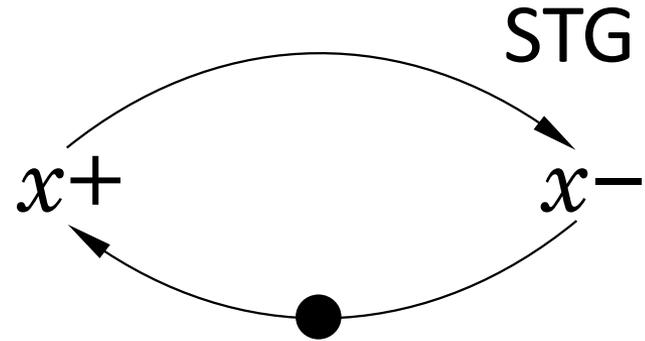
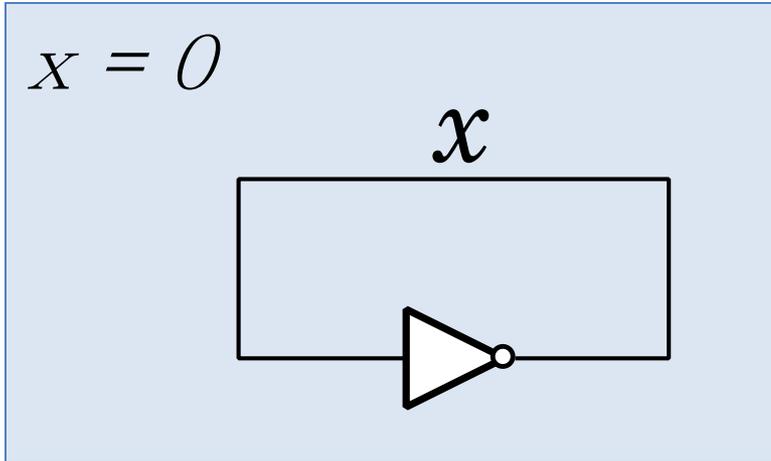


Describing cycles

$$x = 1$$

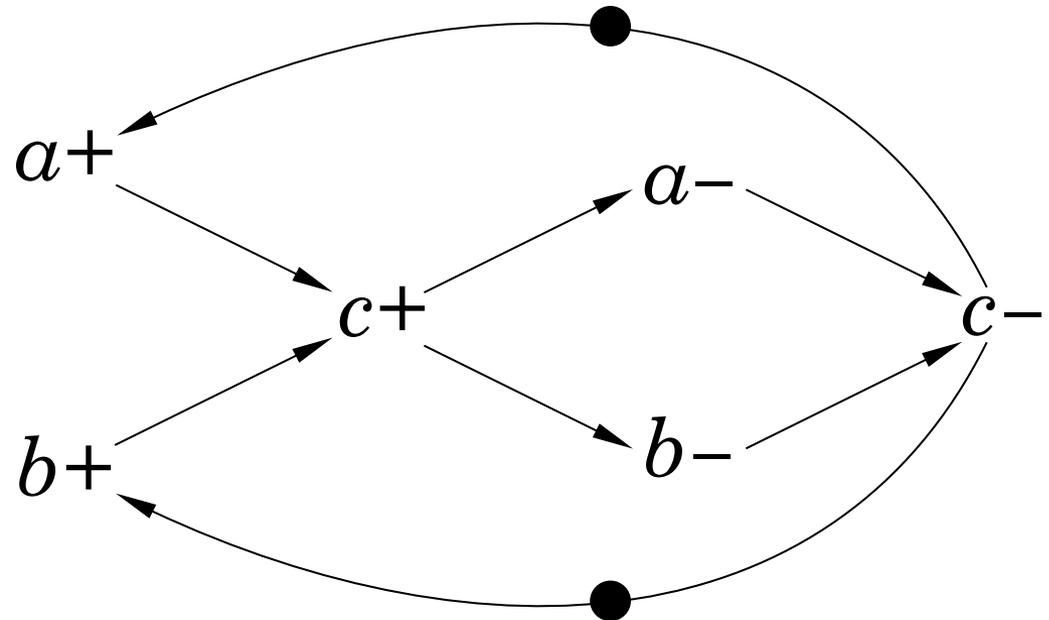
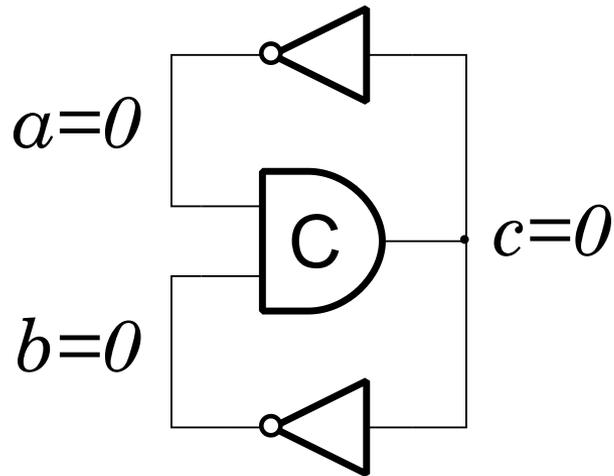


Describing circuits: oscillator



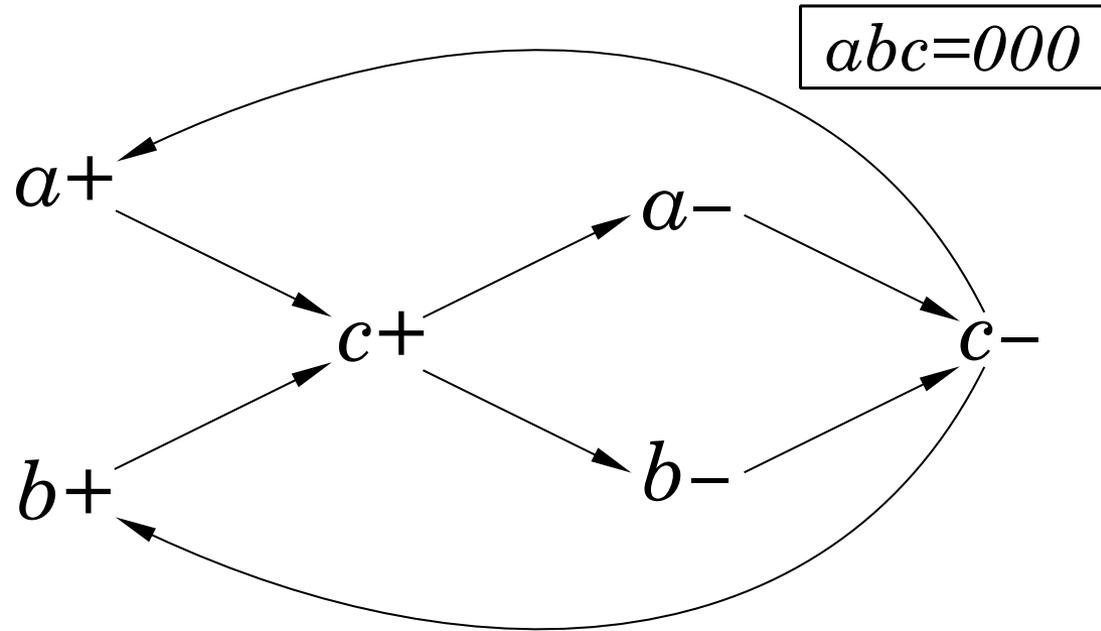
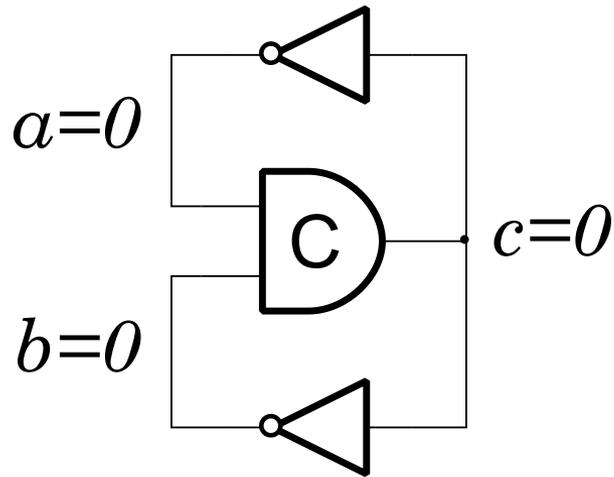
Conditional Signal Graphs (CSGs)

Describing circuits: C-element



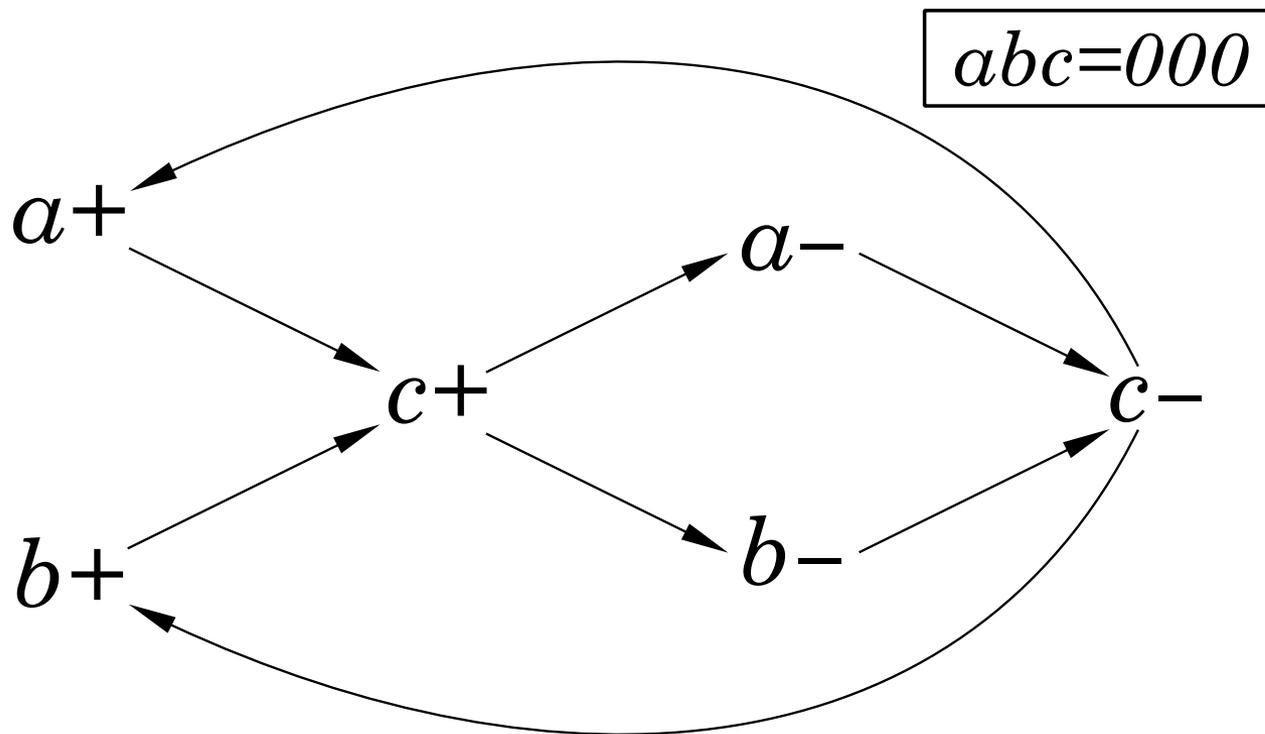
STG

Describing circuits: C-element

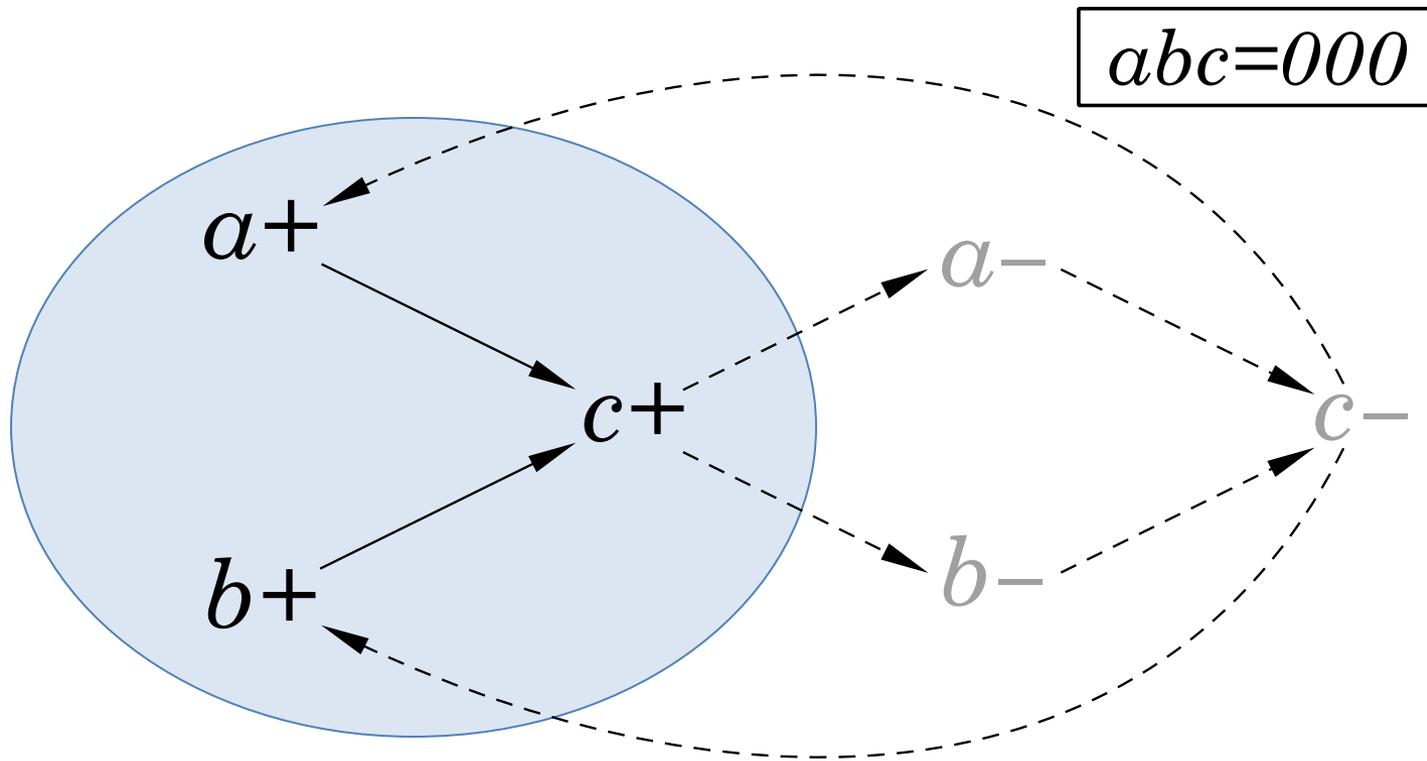


CSG

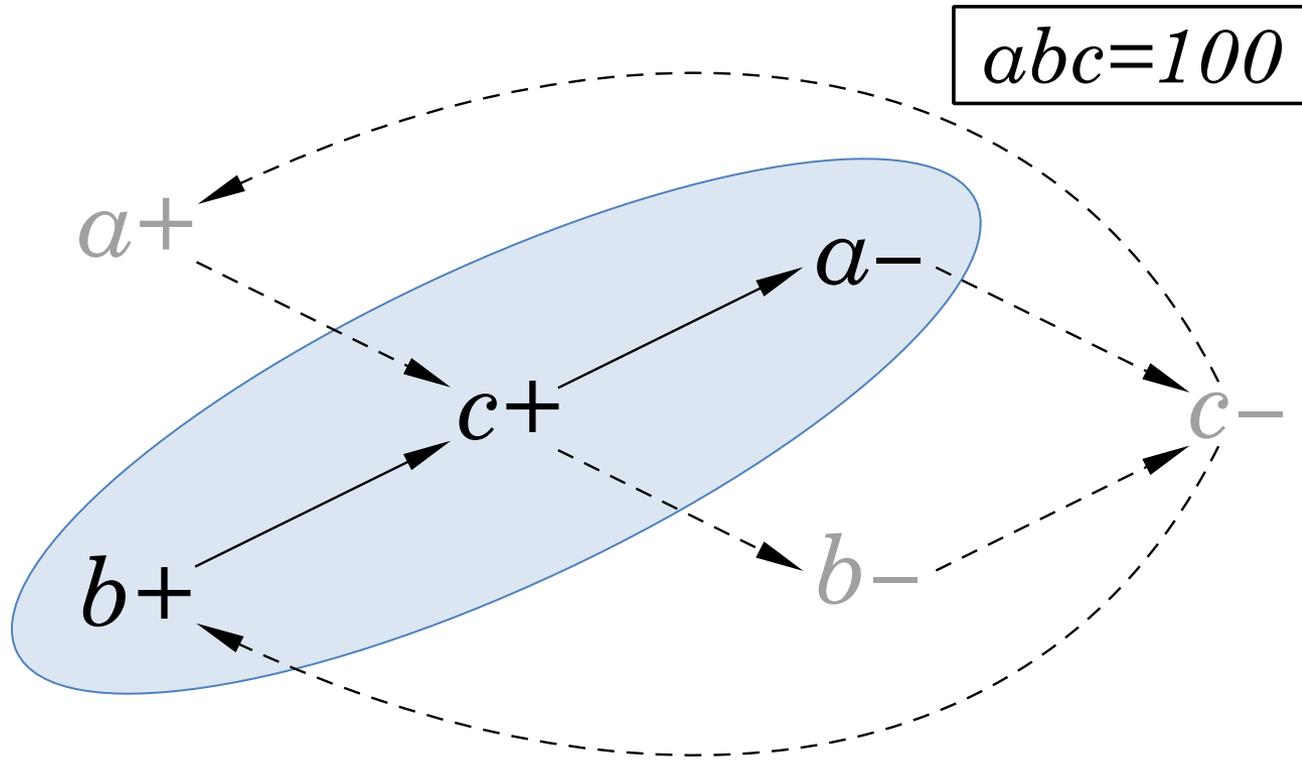
Describing circuits: C-element



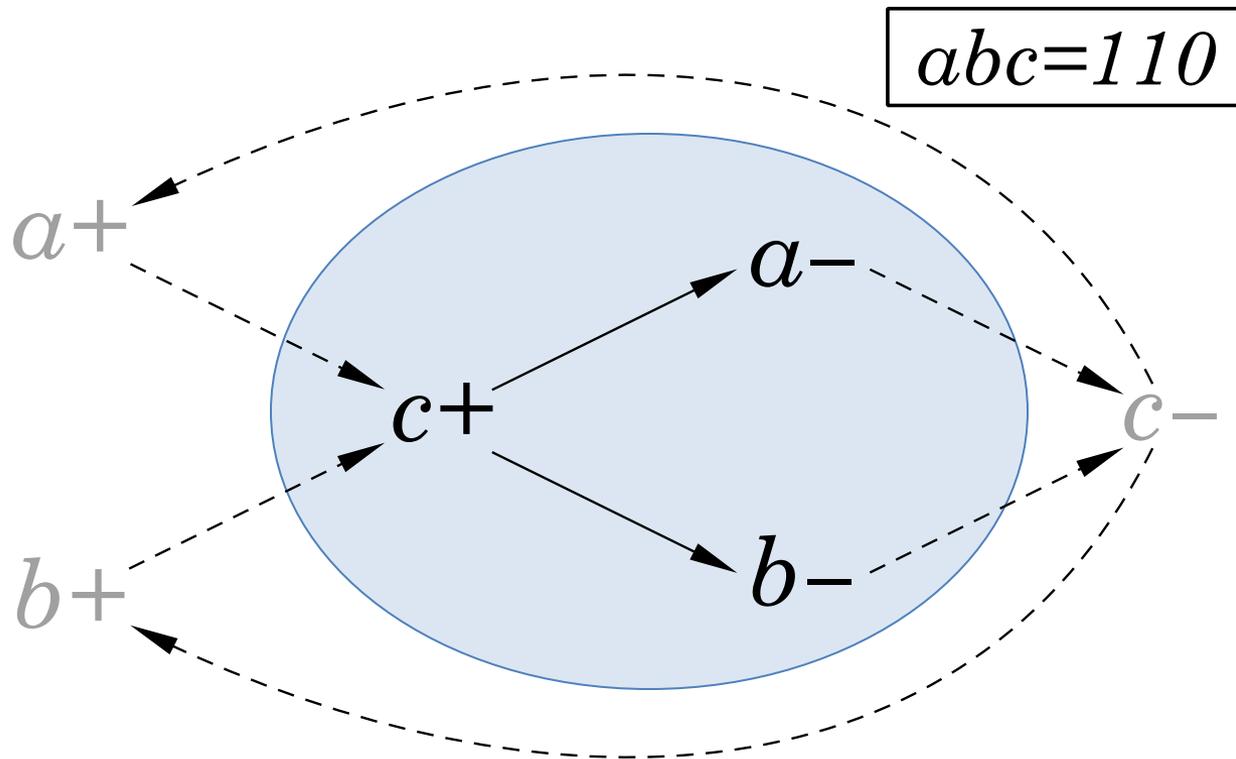
Describing circuits: C-element



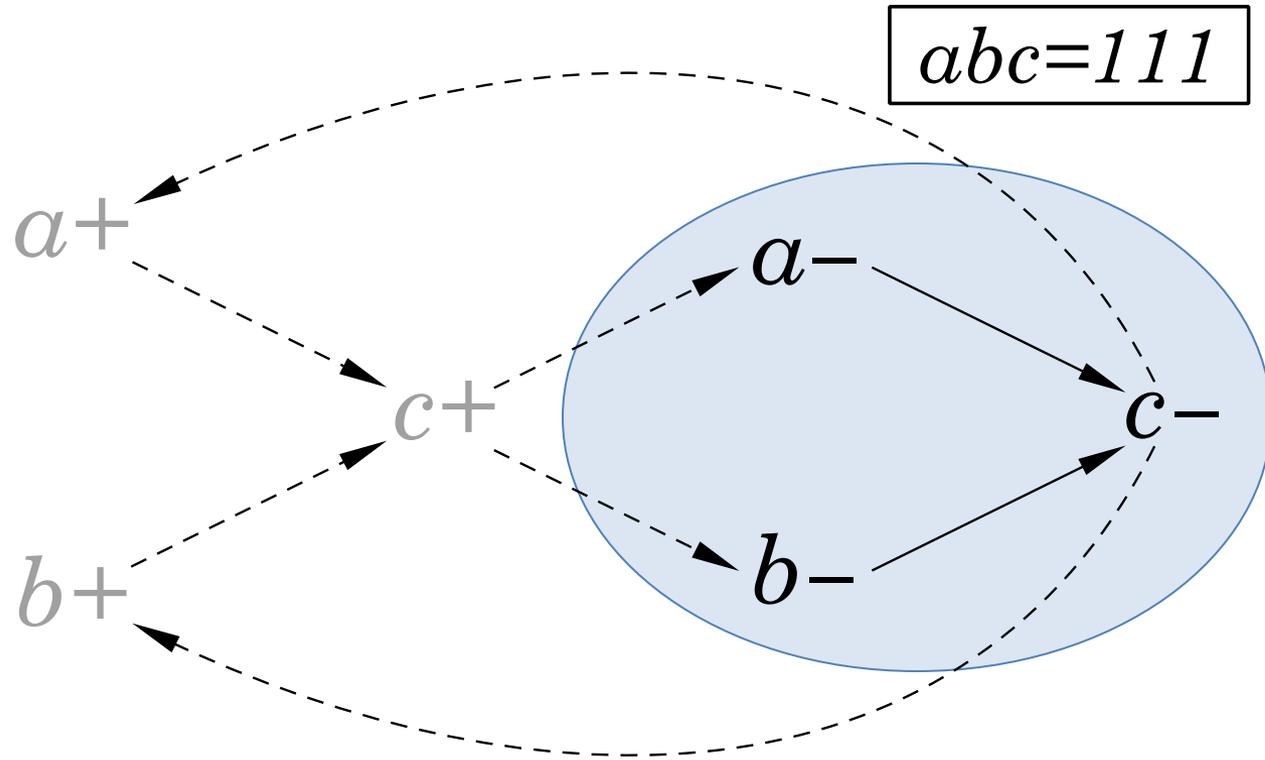
Describing circuits: C-element



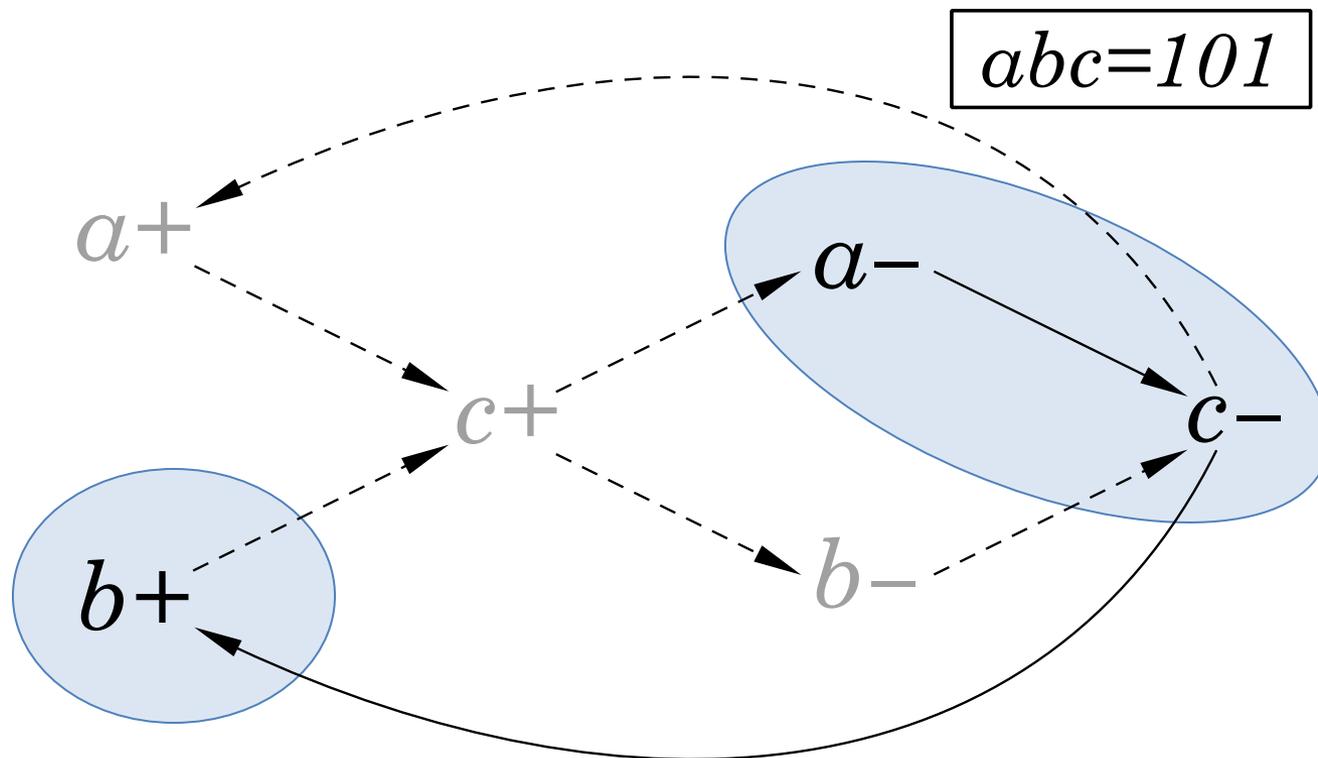
Describing circuits: C-element



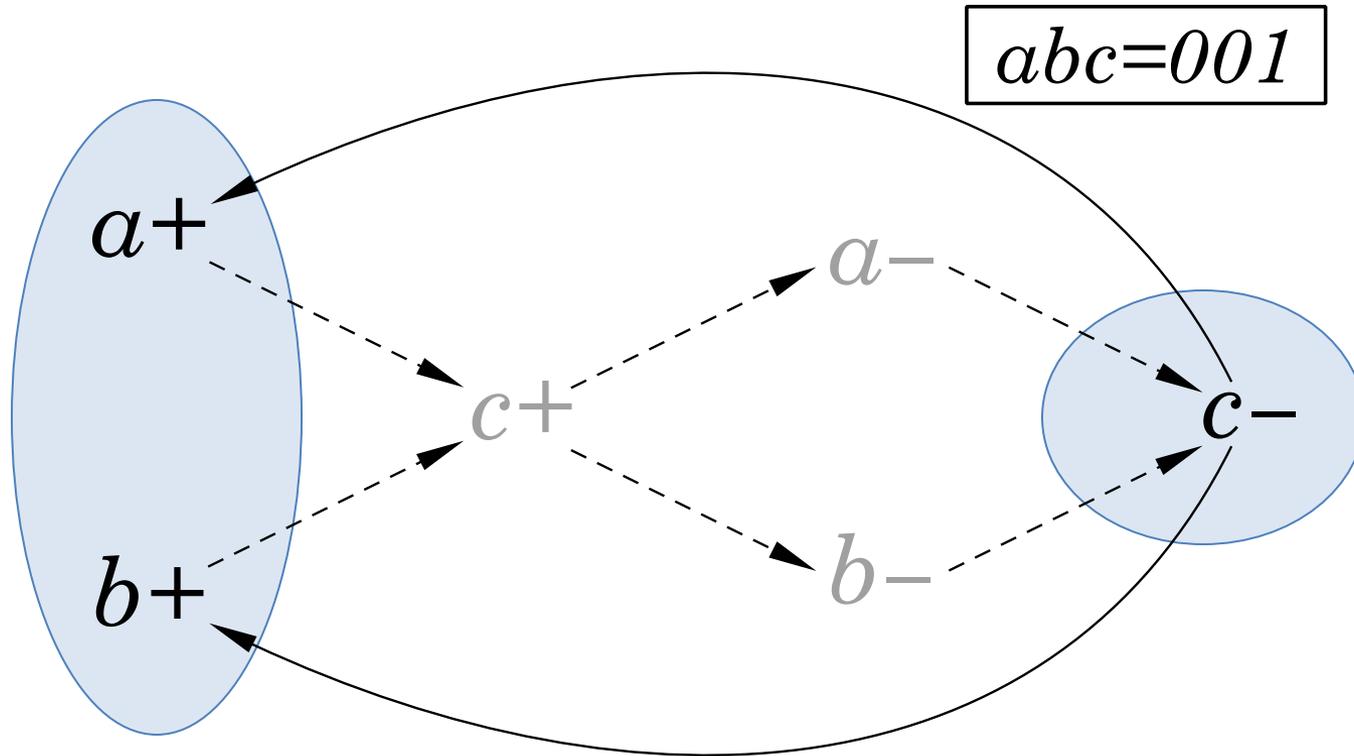
Describing circuits: C-element



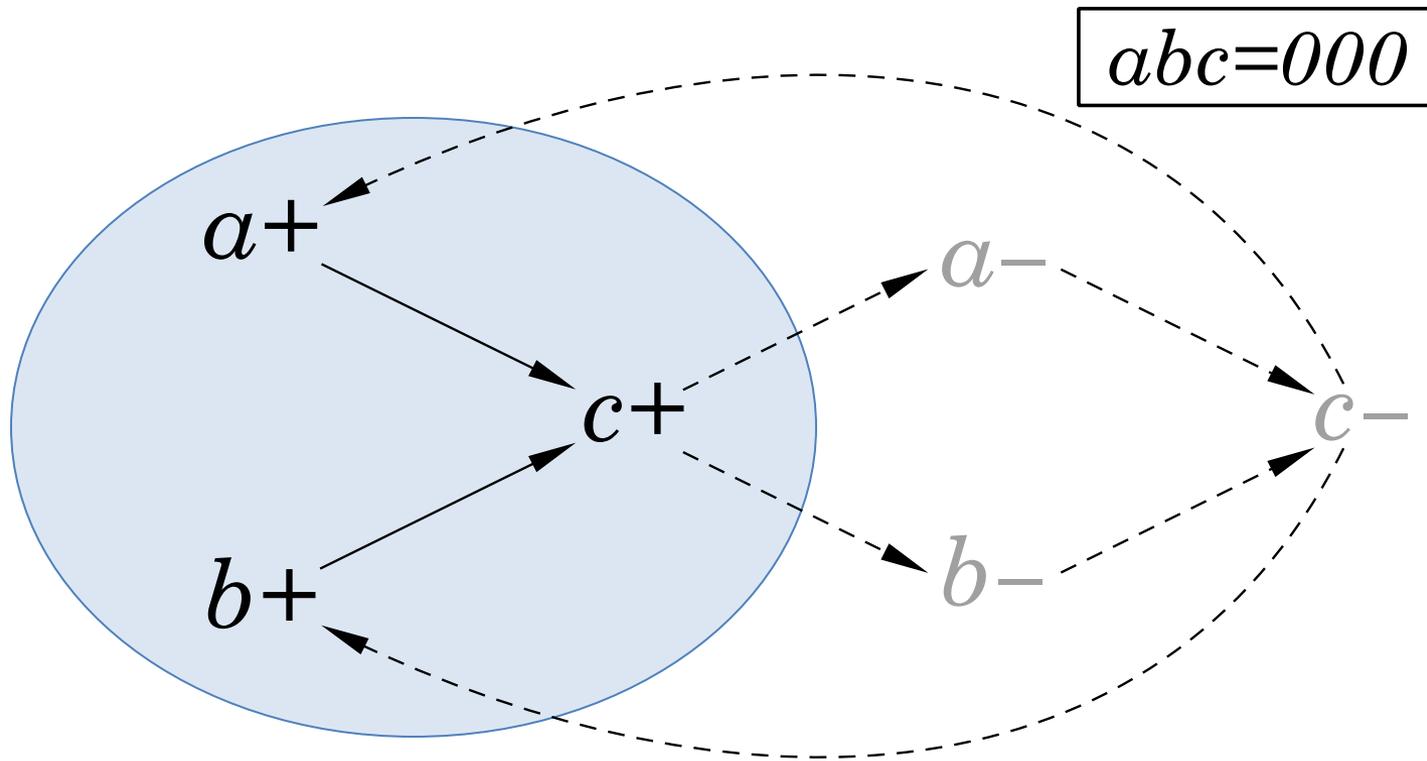
Describing circuits: C-element



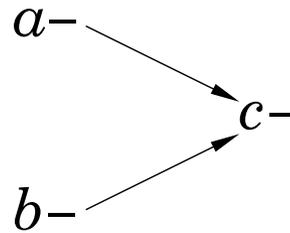
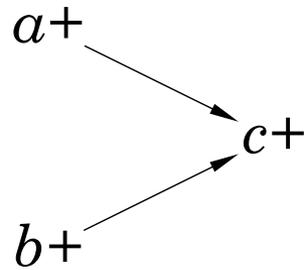
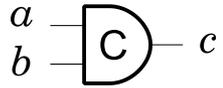
Describing circuits: C-element



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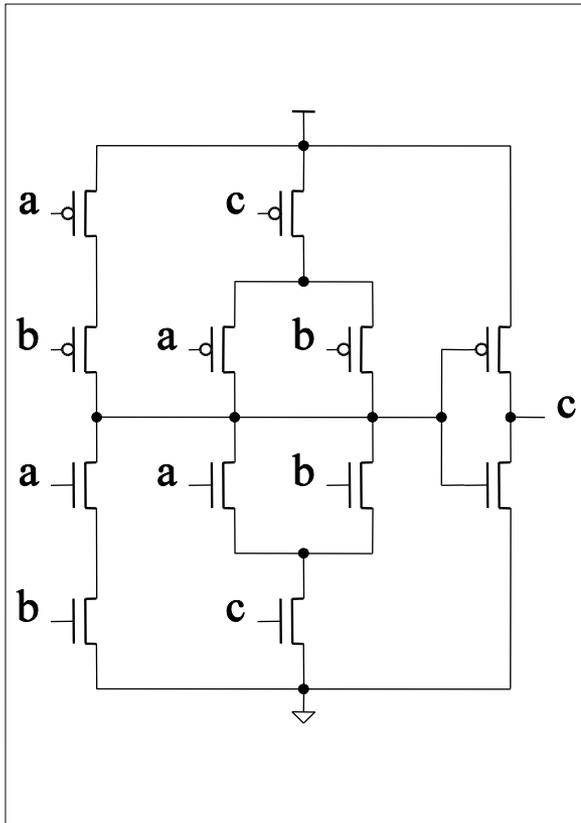
Circuit composition



C-element

$$c = ab + c(a + b)$$

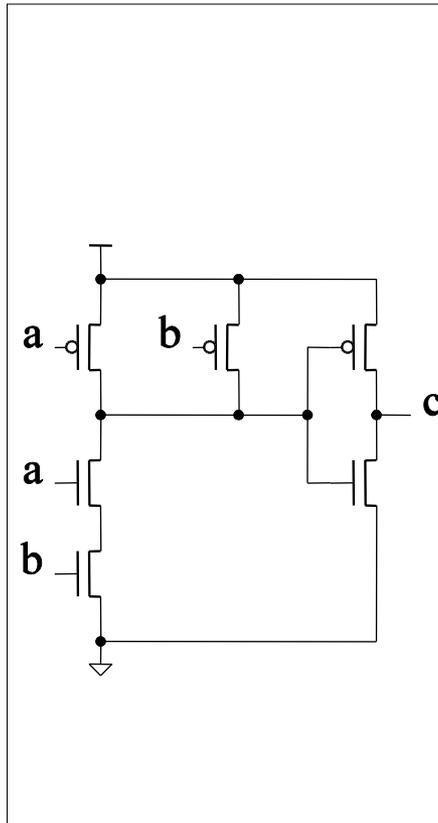
Circuit composition



C-element:

$\uparrow \max(a\uparrow, b\uparrow)$

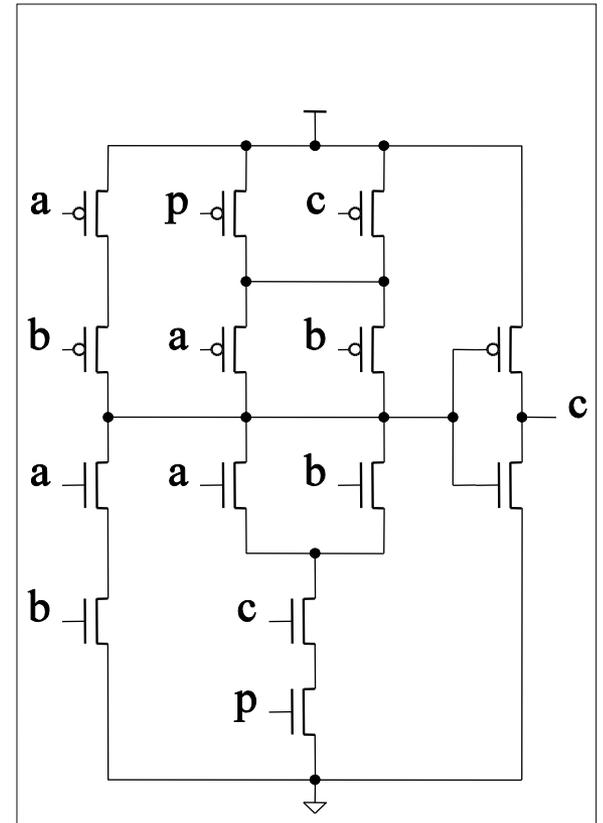
$\downarrow \max(a\downarrow, b\downarrow)$



AND-gate:

$\uparrow \max(a\uparrow, b\uparrow)$

$\downarrow \min(a\downarrow, b\downarrow)$



C/AND-element:

can be chosen

in run-time!

Circuit composition

- Algebraic operations with CSGs:

- Addition (overlay): $G_1 + G_2$

- Scalar multiplication (encoding): $f \cdot G$

- Composition of C-element and AND-gate:

$$G = p \cdot G_C + \bar{p} \cdot G_{AND}$$

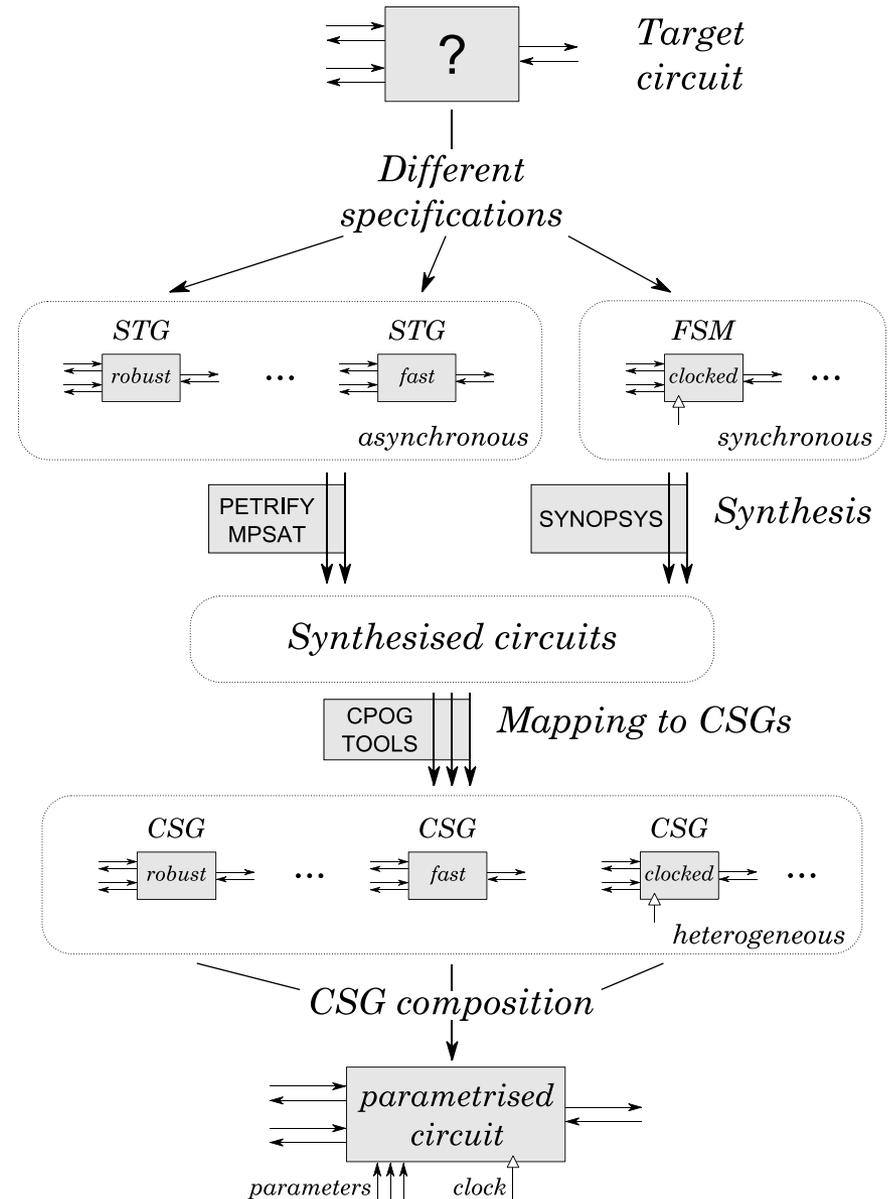
- General composition of n circuits:

$$G = f_1 \cdot G_1 + f_2 \cdot G_2 + \dots + f_n \cdot G_n$$

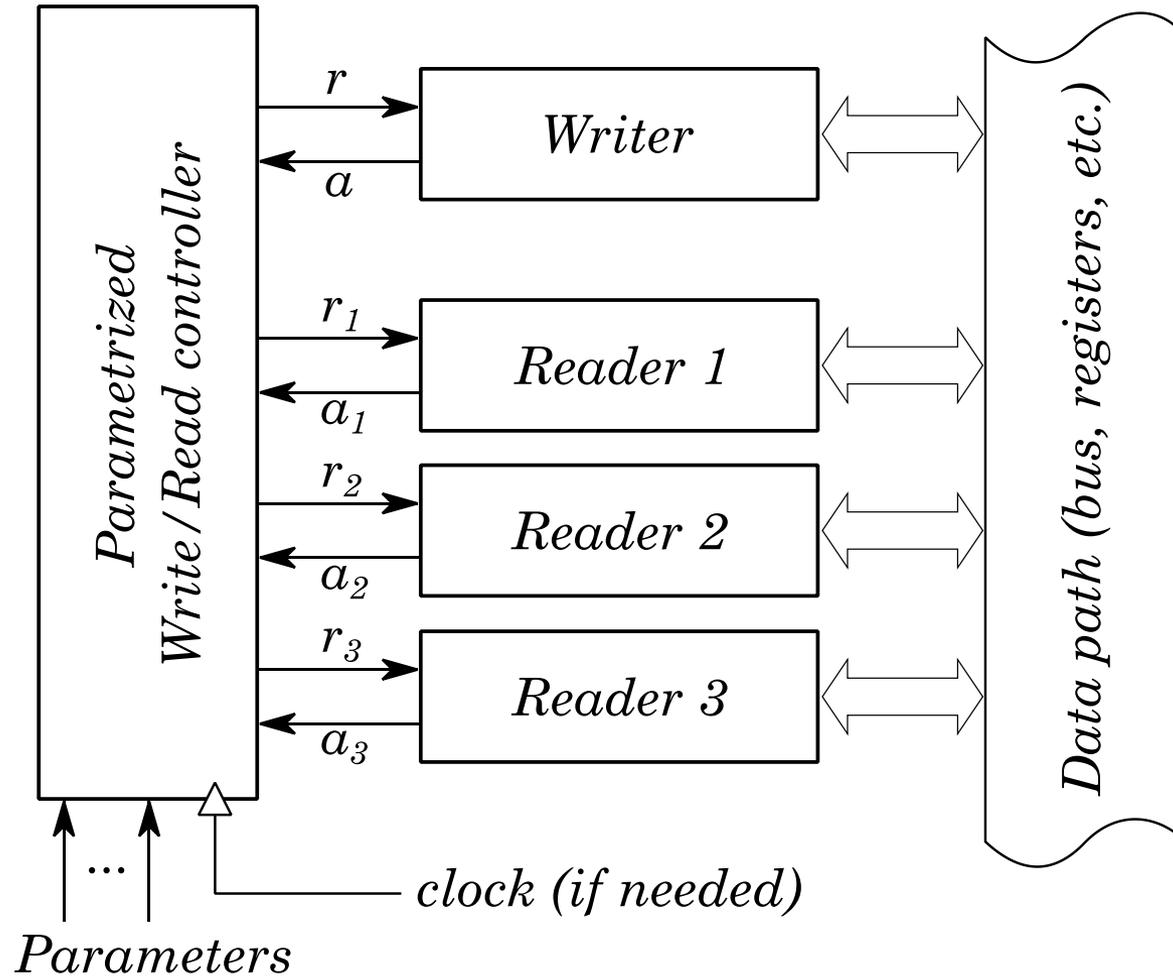
- Fast structural operation (if encoding is given)

Design flow

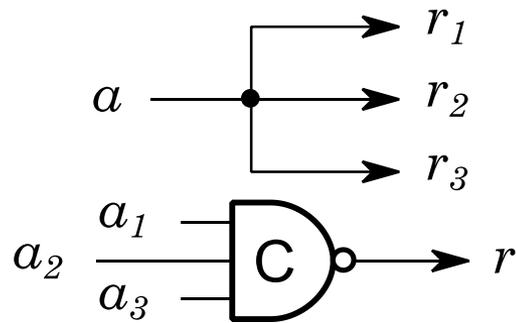
- Heterogeneous models and implementation styles
- Not computationally expensive:
 - composition is fast
 - exact encoding is slow but there are fast approximate methods



Example: Read/Write controller

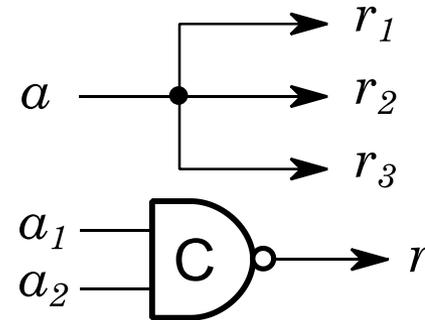


Example: possible implementations



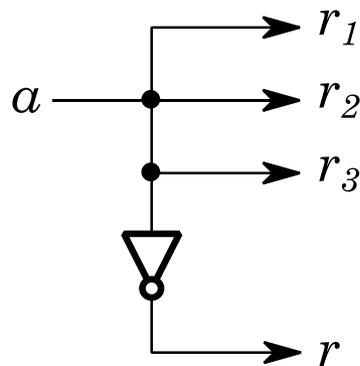
Delay Insensitive (DI)

$$T = d + \max\{d_1, d_2, d_3\} + C_3$$



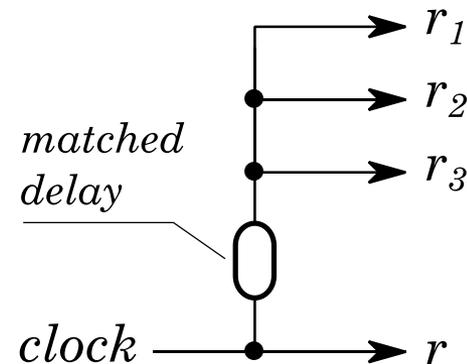
Partial Acknowledgement (PA)

$$T = d + \max\{d_1, d_2\} + C_2$$



Timing Assumptions (TA)

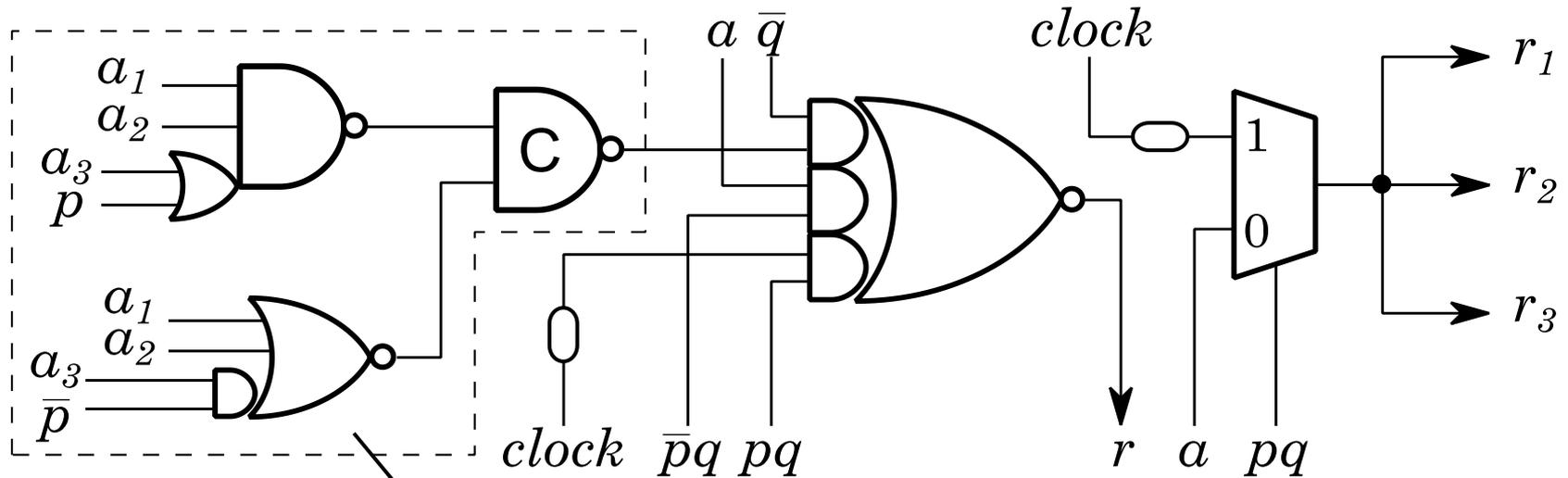
$$T = d + l$$



Synchronous (CL)

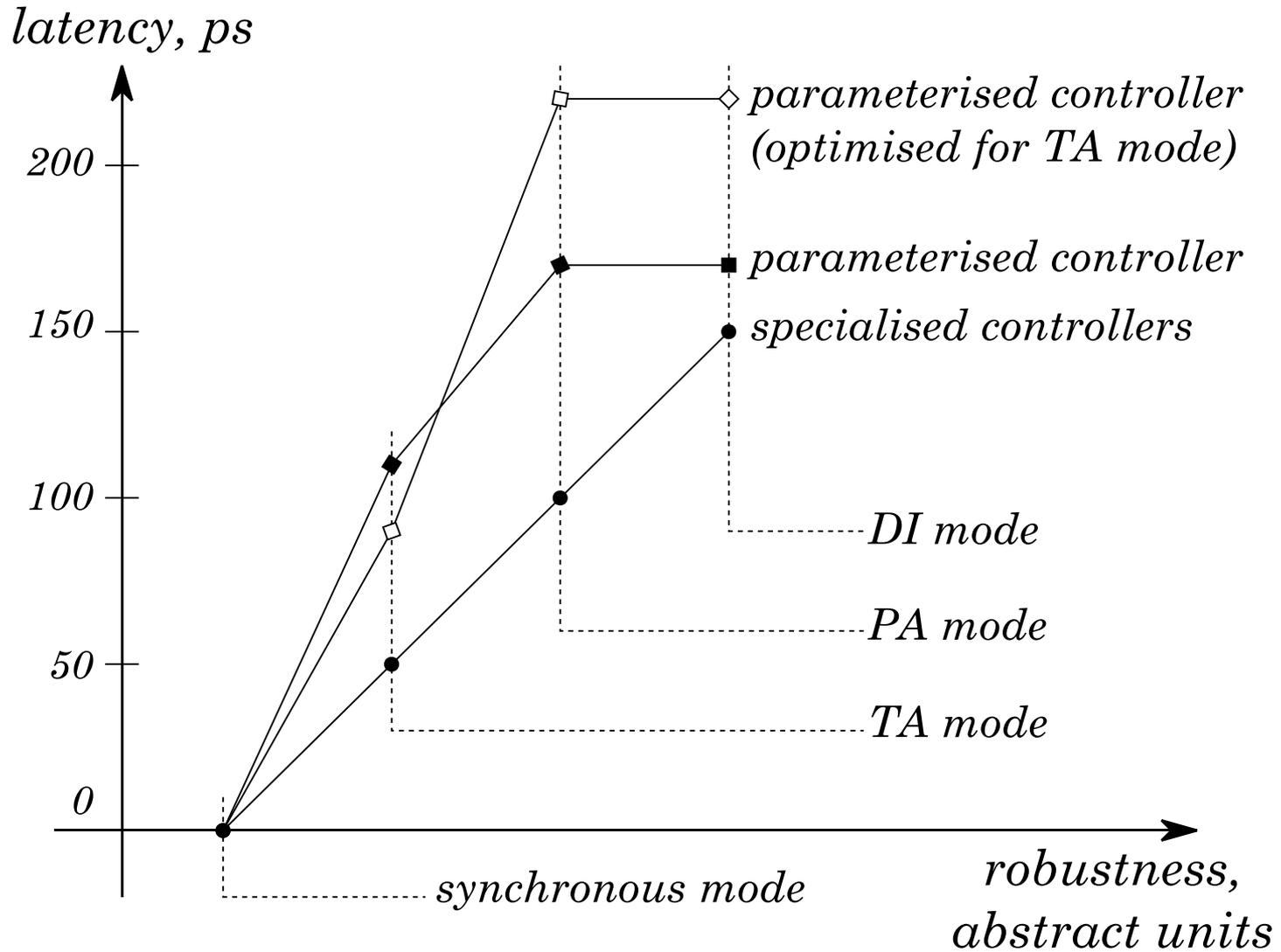
$$T = t_{\text{clock}}$$

Example: parameterised controller



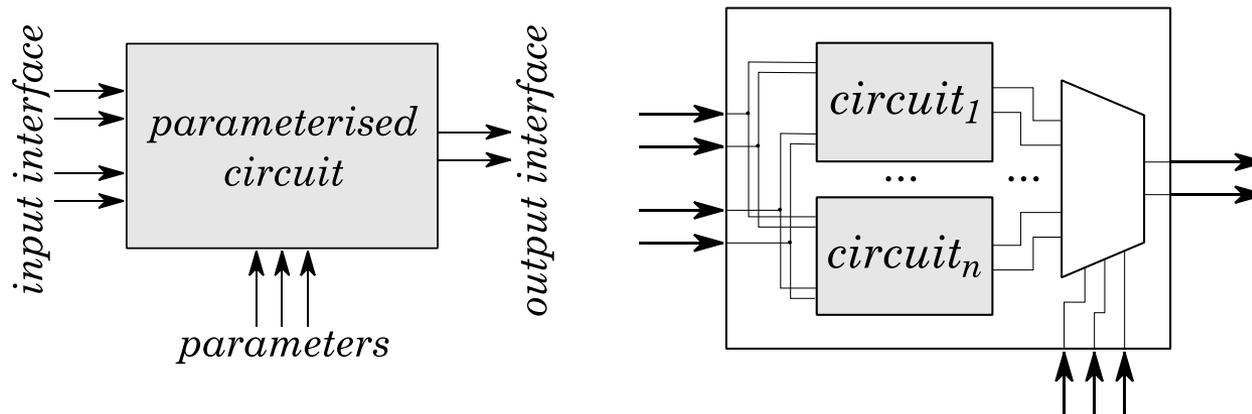
Can be switched off by power-gating in TA/CL modes

Example: simulation results



Cost of Parameterisation

- Trivial approach:
 - Latency overhead: 120-182%, **122%** on average
 - Energy overhead: 151-330%, **201%** on average



- Our approach:
 - Latency overhead: 0-40%, **19%** on average
 - Energy overhead: 0-98%, **23%** on average
- Latency/energy **savings** at the system level!

Conclusions & Future work

- *“Don’t oppose different implementation styles, take advantage of their benefits!”* (reviewer X)
- New model for circuit-level description and composition
- Cost of parameterisation is not prohibitive
- Future work:
 - Tool prototyping
 - CSC signals awareness
 - Power-gating mechanisms

Questions?