Manpower Planning: Rostering

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Who am I?

- Anders Høeg Dohn, 29 years old.
- Cand. Polyt. in 2006 (from DTU).
 - Applied mathematics with specialization in Operations Research.
 - Master's thesis about a practical crew scheduling problem.
- PhD in Operations Research in 2010 (from DTU).
 - "Rostering and Task Scheduling Applications in Manpower Planning"
 - Supervisors: Jens Clausen and Jesper Larsen.
- From July 2010:
 - Operations Analyst at Copenhagen Airports A/S.
- Board member and treasurer of the Danish Operations Research Society (DORS)
 - www.dorsnet.dk

Outline of the Remaining Three Lectures



- Rostering
 - Introduction to Manpower Planning
 - What is Rostering?
 - Solving Rostering Problems with Column Generation
- Task Scheduling
 - What is Task Scheduling?
 - Solving Task Scheduling Problems with Column Generation
 - Your Assignment
- Other OR-Problems in Airports
 - "Being an Operations Analyst in an Airport is like being a kid in a candy store"!
 - Questioning session on the assignment

Scope

- During these lectures I will:
 - Go over some of the practical problems encountered in manpower planning.



- Task Scheduling
- Propose models that can be used to solve these problems, i.e. present case studies of these methods.
 - Integer Programming
 - Set Partitioning Formulations
 - Column Generation
 - Branch & Price

Motivation for using OR in manpower planning



- Large savings of structured manpower planning have been documented in the literature.
- The planning process consists of a number of steps and starts several months before the day of operation.
 - Each step is an interesting operations research problem, where structured forecasting, simulation and optimization may introduce significantly improved results.
- In many cases, rostering and task scheduling can be modeled and solved as well defined optimization problems.



- From: Long term strategic planning
- To: Disruption management on the day of operation





























11 **DTU Management Engineering,** Technical University of Denmark Rostering 22/11/10

Rostering



- Allocate employees to shifts
- Assuming that:
 - Shift demands exist
 - The staff is fixed
 - We don't specify tasks during shifts

Some general characteristics

- Fixed planning period.
- Fixed number of shifts.
- Time norm for each employee.
- Maximum number of days on in a week / on-stretch.
- A minimum rest period after a shift is required.
- Specific shift transitions are not allowed.
- Each employee may have individual preferences.

Demands



																Da	ys													
Shifts	Skills	Type	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	$26\ 27$,
А		=	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D		\geq	15	15	15	15	12			16	16	16	16	13			15	15	15	15	12			16	16	16	16	13		
D	2	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	1	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	3	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	4	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
L		=			2	2	2					2	2	2					2	2	2					2	2	2		
L	0	\geq	2	2	1	1	1			2	2	1	1	1			2	2	1	1	1			2	2	1	1	1		
Ν	0	=	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	$2 \ 2$	

A Roster

			Days
Nurse	Skills	Cost	$0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \ 13 \ \ 14 \ 15 \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ \ 14 \ 15 \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ \ 16 \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 16 \ 17 \ 18 \ 19 \ 10 \ 10 \ 10 \ 10 \ 10 \ 10 \ 10$
0	0, 1	0	D D D D D D D D D D D D D D D D D D D D
1	0, 2	13	N - D N D N - N D D D D L N - D D D - N
2	0, 3	2	- D D L N D D D D L D D D D D D D D D
3	0, 4	2	D D D D D D D N - D L A L D D D D - L N
4	0, 1	0	D D D D D D D D D D D D D D D D D D D D
5	0, 2	0	D D D D D D D D D D D D D D D D D D D D
6	0, 3	2	D D D D D D L D D L D - D D N D D D D D
7	0, 4	0	D D D D D D D D D D D D D D D D D D D D
8	0, 1	0	D D D D D D D D D D D D D D D D D D D D
9	0, 2	2	- D D D N D D D D D D D L D D D D D D
10	0, 3	15	D D D D D D D D A D - N - D - D D D D D D D
11	0, 4	2	D D D A L D D D D D D D D - N L D D D D
12	0, 1	2	D D D D D D D D D D D D D D L L D D N
13	0, 2	2	D D A D D L D D - N D D D D D D D D D L
14	0, 3	13	D N - L D N - D N N - D D D N - L N N
15	0, 4	12	D L N - D A D D D D D D N - D D N D N -
16	0, 1	2	A D D D L L D - D N D D D D A D D D D D
17	0, 2	12	D N - D D D N - D D A D D N D L D D - N -
18	0, 3	28	L D N N - L L - N - D - N - D D D N
19	0, 4	25	D A D N - D L - D D L - L N - D - L D
20	0, 1	25	- D D D N - D - L D L L D N L D D
21	0, 2	39	D D A N - D - D D D N - A - N - A N D
22	0, 3	36	N - D N D D A N - N D D D D D N
23	0, 4	31	- L L N L N - D N - L - N - L - N - D
24	0, 1	21	L - L D D D A D D N A N D A - A A
25	2	0	D D D D D D D D D D D D D D D D D D D D
26	3	0	D D D D D D D D D D D D D D D D D A D D
27	4	0	D D D D D D D A D D D D D D D D D D D

A Roster-line

			Days
Nurse	Skills	Cost	$0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \ 13 \ \ 4 \ 15 \ 16 \ 17 \ 18 \ 19 \ 20 \ \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 26 \ 27 \ \ 26 \ 26 \ 26 \ 26 \ 26 \ 26 \ 26 $
0	0, 1	0	D D D D D D D D D D D D D D D D D D D D
1	0, 2	13	N - D N D N - N D D D D L N - D D D - N
2	0, 3	2	- D D L N D D D D L D D D D D D D D D
3	0, 4	2	DDDDDDDN-DLALDDDD-LN
4	0, 1	0	DDDDD DDDDD DDDDD DDDDD
5	0, 2	0	DDDDD DDDDD DDDDD DDDDD
6	0, 3	2	DDDDD DLDDL D - DDN DDDDD
7	0, 4	0	DDDDD DDDDD DDDDD DDDDD
8	0, 1	0	DDDDD DDDDD DDDDD DDDDD
9	0, 2	2	- DDDN DDDDD DDLDD DDDDL
10	0, 3	15	DDDDD DDDAD - N - D - DD DDDDD
11	0, 4	2	DDDAL DDDDD DDD - N LDDDD
12	0, 1	2	DDDDD DDDDD DDDDL LDDN
13	0, 2	2	DDADD LDD - N DDDDD DDDDL
14	0, 3	13	DN-LDN-DNN-LNN
15	0, 4	12	DLN-DADDDDDN-DDNDN-
16	0, 1	2	ADDDL LD - DN DDDDA DDDDD
17	0, 2	12	DN-DDDN-DDADDNDLDD-N-
18	0, 3	28	L D N N - L L - N - D - N - D D D N
19	0, 4	25	DADN-DL-DDL-L-N-D-LD
20	0, 1	25	- D D D N - D - L D L L D N L D D
21	0, 2	39	D D A N - D - D D D N - A - N - A N D
22	0, 3	36	N - D N D D A N - N D D D D D N
23	0, 4	31	- L L N L N - D N - L - N - L - N - D
24	0, 1	21	L - L D D D A D D N A N D A - A A
LN	- D	' - -	A D D D D = - D D N = D = - D N = - D N
27	4	0	DDDDD DDADD DDDDD DDDDD

A Roster-line





A generalized set partitioning formulation



- Model the problem as that of choosing the best feasible roster-line for each employee.
 - Meet the demands as well as possible with these roster-lines.
- A multi-objective problem
 - Meet demands
 - Maximize satisfaction
 - Minimize costs
- A subproblem generates the feasible roster-lines.
 - All rules concerning a single employee are handled in the subproblem.

Branch & Price

- Necessary considerations:
 - How do we model and solve the master problem?
 - How do we model and solve the subproblem?
 - How do we ensure integrality in the master problem?

A generalized set partitioning formulation



- \mathcal{E} Employees
- Demands (Shifts) \mathcal{D}
- \mathcal{R}'_{e} Roster-lines

- $\lambda_e^r = \begin{cases} 1, & \text{if roster-line } r \text{ is chosen for employee } e. \\ 0, & \text{otherwise.} \end{cases}$
- = Amount of under-coverage (slack) for demand d. = Amount of over-coverage (surplus) for demand d.
- $\begin{array}{c} c_e^r \\ c_d^- \\ c_d^+ \end{array}$ Cost of roster-line r of employee e.
- Cost of under-coverage for demand d.
- Cost of over-coverage for demand d.
- a_{ed}^r =1, if demand d is covered by roster-line r of employee e.
- Required number of employees for demand d. b_d

$$\min \sum_{e \in \mathcal{E}} \sum_{r \in \mathcal{R}'_e} c_e^r \lambda_e^r + \sum_{d \in \mathcal{D}} \left(c_d^- s_d^- + c_d^+ s_d^+ \right)$$
$$\sum_{r \in \mathcal{R}'_e} \lambda_e^r = 1 \qquad \forall e \in \mathcal{E}$$
$$\sum_{e \in \mathcal{E}} \sum_{r \in \mathcal{R}'_e} a_{ed}^r \lambda_e^r + s_d^- - s_d^+ = b_d \qquad \forall d \in \mathcal{D}$$
$$\lambda_e^r \in \{0, 1\} \qquad \forall e \in \mathcal{E}, \forall r \in \mathcal{R}'_e$$
$$s_d^- \ge 0, s_d^+ \ge 0 \qquad \forall d \in \mathcal{D}$$

A generalized set partitioning formulation



- The number of variables is HUGE.
 - Any legal combination of shifts constitutes a variable.
 - It is impossible to generate all varibles a priori.
 - Therefore, solving the model using column generation is a natural choice.
- Column generation is used directly on the generalized set partitioning formulation.
 - We are not decomposing a compact formulation.
 - The compact model can be "reverse engineered", but it is not easy and the model usually isn't very nice, anyway.

Branch & Price

DTU

- Necessary considerations:
 - How do we model and solve the master problem?
 - How do we model and solve the subproblem?
 - How do we ensure integrality in the master problem?

Column generation



Master Problem



1 0 0 1 0 0 1

Subproblem

Column Generation - Overview



Column generation

- Solve an LP-relaxation of the problem.
- One subproblem for each employee
 - The subproblems can be solved independently, in parallel or sequentially.
 - To terminate column generation, all subproblems must be unable to return new columns.
 - The subproblem can be solved as a Resource Constrained Shortest Path Problem.

Dynamic Programming

- Resources
 - Anything that can affect feasibility or cost of a label must be considered when dominating labels.
 - Such measures are represented by *resources*.
- Domination
 - For each resource, we must specify what a "better" value is.
 - Domination must consider both feasibility and cost.
 - Sometimes, it is not better to have neither a smaller nor a larger value. Then domination only applies if values are equal.
 - All resources must allow domination (and cost must be lower) for one label to dominate another.

Dynamic Programming









- Assume that we have two resources: number of worked hours and number of night shifts.
- Assume that both D, L and N have a length of 8 hours.
- Is it better to have a smaller/larger number of hours?

Dynamic Programming

I = (cost, hours, nightshifts)



- I_3 does not dominate I_1 because it has a larger number of night shifts.
 - Assuming that lower is better.
 - Why is lower better? Must be defined in the individual case for each of the resources.
 - Due to feasibility
 - Due to costs which depend on resource values
 - What if it is the other way around?

Utilizing the roster-line structure



- Combine shifts into on-stretches
- Combine on-stretches and off-stretches into work-stretches
- Combine work-stretches into roster-lines
- Use dynamic programming to find only the non-dominated entities in each step.

Utilizing the roster-line structure



Utilizing the roster-line structure



An optimal LP-solution

																Dε	ays													
Nurse	х	Cost	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27
0	0.378191	2	D	D	D	D	D	-	-	D	\mathbf{L}	D	D	D	-	-	-	L	L	D	Ν	-	-	D	Ν	-	D	D	-	-
0	0.621809	0	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-
1	0.966728	2	D	\mathbf{L}	D	D	D	-	-	D	L	-	D	Ν	-	-	Ν	-	L	D	D	-	-	D	D	D	D	D	-	-
1	0.0332715	5	D	D	Ν	-	\mathbf{L}	-	-	D	D	D	D	D	-	-	Α	Α	D	Α	\mathbf{L}	-	-	Ν	-	Ν	-	-	-	Ν
2	0.843555	0	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-
2	0.0764406	2	D	D	D	D	Α	-	-	L	D	D	Α	D	-	-	D	D	D	D	\mathbf{L}	-	-	D	Α	D	-	Ν	-	-
2	0.0800048	2	N	-	Ν	-	\mathbf{L}	-	-	L	D	-	D	Ν	-	-	L	D	L	D	D	-	-	D	D	\mathbf{L}	D	D	-	-
3	0.474511	0	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-
3	0.23604	2	L	D	D	L	D	-	-	-	N	-	Α	Ν	-	-	D	L	D	D	D	-	-	D	D	D	D	D	-	-
3	0.18799	5	N	-	D	D	D	-	-	A	D	N	-	D	-	-	D	Ď	L	D	D	-	-	N	-	D	-	D	-	Ν
3	0.0410925	2	-	L	D	D	N	-	-	L	D	Đ	D	D	-	-	A	A	D	D	D	-	-	D	D	L	N	-	-	-
3	0.0603669	12		- NT	IN	- NT	А	-	-	IN D	-	L	D	D	-	-	A	D	D	D	D	-	-	D	D	- T	D	D	IN	- N
4	0.260215	5		N	- D	N	- D	-	-	D	D	D	D	A	-	-	D	IN D	- D	D	D	-	-	D	D	L	D	- D	-	IN
4	0.071949	10		D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D N	D	D	-	-		D	D	D	D	-	- N
4	0.451775	13		- N	- T	IN D	- D	-	-		Б	D T	D	L	-	-	D	D	IN D	- D		-	-		- N	IN	-	D N	-	IN
4	0.210001					D	D	-	-	A			D	D	-	-	D	D	D	D	A D	-	-			- D	- D	D	-	-
5	0.707095	5				D	D	-	-	D	D D	D	D N	D	-	-			D	D	D	-	-	D	D N	D	D N	D	-	- N
5	0.106765	0			T	T	D	-	-		T	D	D	- -	-	-	Б П	D A	D	D	D	-	-	N	IN	-		- N	-	IN
6	0.0033403			D	Ď	D D	D	-	-	D	Б.	D	D	D	-	-	D	D D	D D	D	D	-	-	IN	Ē	-	D	D	N	-
7	0 125366	9		D	D	D	D	_	_	D	D	D	D	L	_	_	L	D	D	D	N	-	-	D	Δ	D	Δ	D	11	-
7	0.397964	2	D	Ď	Ď	Ď	Ľ	_	_	D	Ď	Ľ	Ă	Ď	_	_	Ď	Ď	D	Ď	A	_	_	D	Ď	-	Â	N	_	_
7	0.0576867	2	Ď	Ď	-	Ď	Ň	_	_	Ľ	Ď	D	Ĺ	Ă	_	_	D	D	Ď	D	D	_	_	Ľ	D	D	Ň	-	_	_
7	0.418983	Õ	Ď	Ď	D	Ď	D	-	-	Ď	Ď	Ď	Ď	D	_	_	Ď	Ď	Ď	Ď	Ď	_	_	D	Ď	Ď	Ď	D	-	_
8	0.422656	5	D	Ď	Ď	Ľ	Ď	_	_	D	Ň	-	N	-	_	_	Ď	Ď	Ď	Ď	Ď	_	_	-	Ď	Ň	-	Ă	-	Ν
8	0.577344	ŏ	D	Đ	Đ	D	Đ	-	-	D	Ď	D	D	D	-	-	D	Đ	Đ	D	Đ	-	-	D	Đ	D	D	D	-	_
9	0.0772873	13	D	D	Ā	N	_	-	_	D	D	N	_	D	-	_	N	_	Ā	D	D	-	-	N	_	N	_	Ĺ	-	Ν
9	0.722712	0	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	$\overline{\mathrm{D}}$	-	-
9	0.128979	5	D	D	\mathbf{L}	Ν	-	-	-	D	Ν	-	Ν	-	-	-	D	D	D	D	D	-	-	D	Α	-	D	D	-	Ν
9	0.0710214	2	A	D	D	D	Α	-	-	L	D	\mathbf{L}	Ν	-	-	-	Α	-	D	D	Ν	-	-	L	D	D	D	D	-	-
10	0.301887	2	N	-	-	D	Ν	-	-	L	D	\mathbf{L}	\mathbf{L}	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-
10	0.0220762	2	D	D	Α	\mathbf{L}	D	-	-	D	D	D	\mathbf{L}	D	-	-	Α	Ν	-	D	D	-	-	-	\mathbf{L}	D	D	Ν	-	-
10	0.174596	2	D	Ν	-	D	Α	-	-	D	D	D	D	D	-	-	\mathbf{L}	L	D	D	D	-	-	D	D	-	\mathbf{L}	Ν	-	-
10	0.501441	0	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D	-	-
11	0.753929	2	D	-	D	D	Ν	-	-	D	D	D	D	L	-	-	Ν	-	D	D	L	-	-	\mathbf{L}	D	D	\mathbf{L}	D	-	-
11	0.246071	5	A	D	Ν	-	D	-	-	D	D	D	D	D	-	-	\mathbf{L}	D	D	D	D	-	-	Ν	-	-	Ν	-	-	Ν
12	0.947893	14	D	D	D	D	D	-	-	D	D	D	D	-	Ν	-	D	D	-	D	D	-	-	D	D	D	D	D	-	-
12	0.0229911	13	N	-	Ν	-	D	-	-	D	Ν	-	D	L	-	-	D	Α	Ν	-	D	-	-	A	D	Ν	-	D	-	Ν
12	0.0291159	4	A	A	Ν	-	D	-	-	D	D	L	D	D	-	-	D	D	L	Ν	-	-	-	-	-	D	\mathbf{L}	Α	Ν	-
13	0.230343	12	D	A	D	D	D	-	-	D	N	-	L	D	-	-	A	Ņ	-	D	D	-	-	D	D	Ν	-	-	Ν	-
13	0.069953	2	A	L	D	D	D	-	-	D	A	Ν	-	D	-	-	D	A	D	-	Ν	-	-	L	D	L	D	D	-	-
13	0.179915	12	N	-	D	Ν	-	-	-	D	Ν	-	D	D	-	-	D	L	D	А	D	-	-	D	-	D	D	D	Ν	-



Branch & Price

DTU

- Necessary considerations:
 - How do we model and solve the master problem?
 - How do we model and solve the subproblem?
 - How do we ensure integrality in the master problem?

Branch & Price - Overview





Branch & Price - Overview



Branching

- Variable branching is usually not a possibility in column generation.
 - How do you prohibit a variable from reappearing?
- Branch on a fractional shift allocation.
- When all shift allocations are integer the full master problem solution is integer.
 - As the constraint matrix for each employee is a perfect matrix.
 - For more on this: Follow the course "42122 The Set Partitioning Optimization Model and its Application in Practical Scheduling Problems".

Branching

															Da	ys												
Nurse	x	Cost	0	1 1	2 3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24 5	$25 \ 2$	5 27
$\begin{pmatrix} 0 \end{pmatrix}$	0.378191	2	D	DΙ) D	D	-	-	D (L)	D	D	D	-	-	-	L	L	D	Ν	-	-	D	Ν	-	D	D -	-
0	0.621809	0	D	DI) D	D	-	-	D	Đ.	D	D	D	-	-	D	D	D	D	D	-	-	D	D	D	D	D -	-
1	0.966728	2	D		ע נ	D	-	-	D	L	-	D	N	-	-	N	-	L	D	D	-	-	D	D	D	D	D -	- N
$\frac{1}{2}$	0.0332715	0 0			- v	D D	-	-	D	D	D	D	D	-	-	A D	A D	D	D A	Б	-	-	D	Ē	D	- D		IN
2	0.0764406	2	D	DI	\hat{D}	Ă	_	_	L	D	D	Ă	D	_	_	D	D	D	Ď	L	_	_	D	Ă	D	-	N -	_
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Branch & Price

DTU

- Necessary considerations:
 - How do we model and solve the master problem?
 - How do we model and solve the subproblem?
 - How do we ensure integrality in the master problem?
- Other considerations:
 - Master problem
 - Solve it to optimality every time?
 - Use dual stabilization?
 - Subproblem
 - Use heuristics?
 - In what order should the individual subproblems be solved?
 - Branching



- How do we search the branch-and-bound tree?
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Searching the Branch-and-Bound Tree

- Best first
 - Theoretically the most efficient
- Depth first
 - Works very well in practice for rostering problems.
 - Efficient because many nodes have the same lower bound.
 - Often, a solution exists in the bottom of the tree with solution value equal to the lower bound.





Branch & Price

DTU

- Necessary considerations:
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- How do we search the branch-and-bound tree?
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An optimal solution



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An optimal solution: slack / surplus



																Da	ys													
Shifts	Skills	Type	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	$26\ 27$	
А		=	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D		\geq	15	15	15	15	12			16	16	16	16	13			15	15	15	15	12			16	16	16	16	13		
D	2	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	1	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	3	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
D	4	\geq	1	1	1	1	1			1	1	1	1	1			1	1	1	1	1			1	1	1	1	1		
L		=			2	2	2					2	2	2					2	2	2					2	2	2		
L	0	\geq	2	2	1	1	1			2	2	1	1	1			2	2	1	1	1			2	2	1	1	1		
Ν	0	=	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	$2 \ 2$	

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А		=	0	-	-	-	-	-			-	-	-	-	-		-	-	-	-	-		-	-	-	-	-	
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D	1	\geq	0	3	5	5	6	4			4	5	4	5	4		4	4	5	4	2		4	3	4	5	4	
D	2	\geq	0	4	4	3	4	3			5	3	4	4	4		4	5	4	4	3		4	3	5	5	3	
D	3	\geq	0	3	4	4	2	3			4	3	4	2	1		3	2	3	5	5		5	5	3	3	3	
D	4	\geq	0	5	3	4	2	3			3	5	2	3	6		4	3	2	2	3		3	4	2	2	5	
L		=	0			-	-	-					-	-	-				-	-	-				-	-	-	
L	0	\geq	0	-	-	1	1	1			-	-	1	1	1		-	-	1	1	1		-	-	1	1	1	
Ν	0	=	0	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	-		-	-	-	-	-	

Rotation Patterns

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- Some companies use rotation patterns.
 - Because that is how they have always done it.
 - Because they are easier to create and maintain manually.
 - Because they are easier to discuss with unions and managers.
- Rotation patterns are more restrictive than individual roster-lines.
 - Hence, less desirable from an optimization point of view.

	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday	Sunday
Employee 1	D	D	-	-	А	Ν	-
Employee 2	L	А	L	D	D	-	-
Employee 3	Ν	-	D	D	L	L	L
Employee 4	-	-	-	L	Ν	-	-
Total	D, L, N	D, A	D, L	2 x D, L	D, L, A, N	L, N	L

Rotation Patterns

- We can no longer create an individual roster-line for each employee.
- Several employees share one rotation.
 - The rotation can be generated by solving the subproblem from before with some modifications.
 - The pattern must have a feasible wrap.
 - Demands are given for each weekday (not for each calendar day).
 - Each rotation contributes with several shifts for each day in the roster.
- Works badly when individual skills are important and when we have varying demands on each individual day.

	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday	Sunday
Employee 1	D	D	-	-	А	Ν	-
Employee 2	L	А	L	D	D	-	-
Employee 3	Ν	-	D	D	L	L	L
Employee 4	-	-	-	L	Ν	-	-
Total	D, L, N	D, A	D, L	2 x D, L	D, L, A, N	L, N	L

Summary

- What you should be able to remember from this lecture:
 - Manpower planning consists of several steps
 - Each is an interesting optimization problem in its own right.
 - Rostering
 - The practical problem
 - The set partitioning formulation
 - Characteristics and variations
 - Cost structure
 - Rotations
 - Solution method
 - Column generation
 - Dynamic programming
 - Branching strategy and tree search strategy