



02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.1		
Anne Haxthausen, IMM/DTU	Spring 2002	

Lists Contents

• what is a list	2
• list type expressions	2
• list value expressions	4
• defining infinite lists	7
• list indexing	5
• list operators	6
• examples of modeling using lists	8
• sets, lists and products	17

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.2		
Anne Haxthausen, IMM/DTU	Spring 2002	

What is a list?

A list is: an ordered collection of values of the same type.

Examples:

$\langle 1, 3, 3, 1, 5 \rangle$
 $\langle \text{true}, \text{false}, \text{true} \rangle$

List Type Expressions

- type_expr^*
denotes the type consisting of list values


$\langle v_1, \dots, v_n \rangle$

where $n \geq 0$, $v_i : \text{type_expr}$

- type_expr^ω
denotes the type consisting of list values

$\langle v_1, \dots, v_n \rangle,$
 $\langle v_1, \dots, v_n, \dots \rangle$

where $n \geq 0$, $v_i : \text{type_expr}$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.3		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Type Expressions

Examples

Bool*


denotes the type consisting of the lists:

$\langle \rangle$
 $\langle \text{true} \rangle, \langle \text{false} \rangle$
 $\langle \text{true}, \text{false} \rangle, \langle \text{false}, \text{true} \rangle, \langle \text{true}, \text{true} \rangle, \langle \text{false}, \text{false} \rangle,$
 $\langle \text{true}, \text{false}, \text{true} \rangle,$
 \vdots

Bool^ω

denotes the type consisting of the lists:

$\langle \rangle,$
 $\langle \text{true} \rangle, \langle \text{false} \rangle,$
 $\langle \text{true}, \text{false} \rangle, \langle \text{false}, \text{true} \rangle, \langle \text{true}, \text{true} \rangle, \langle \text{false}, \text{false} \rangle,$
 $\langle \text{true}, \text{false}, \text{true} \rangle$
 \vdots
 $\langle \text{false}, \text{true}, \text{true}, \text{true}, \text{false}, \dots \rangle$
 \vdots

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.4		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Value Expressions

Enumerated:

$\langle 1, 3, 3, 1, 5 \rangle$
 $\langle \text{true}, \text{false}, \text{true} \rangle$

$\langle \text{value_expr}_1, \dots, \text{value_expr}_n \rangle$

Ranged:


$\langle 3 .. 7 \rangle = \langle 3, 4, 5, 6, 7 \rangle$
 $\langle 3 .. 3 \rangle = \langle 3 \rangle$
 $\langle 3 .. 2 \rangle = \langle \rangle$

$\langle \text{value_expr}_1 .. \text{value_expr}_2 \rangle$

Comprehended:

$\langle 2 * n \mid n \text{ in } \langle 0 .. 3 \rangle \rangle$
 $\langle n \mid n \text{ in } \langle 0 .. 100 \rangle \cdot \text{is_even}(n) \rangle$

$\langle \text{value_expr}_1 \mid \text{binding in value_expr}_2 \cdot \text{logical_value_expr}_3 \rangle$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.5		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Indexing

Basic form:

$list_value_expr(value_expr_1)$

Example:


$\langle 2,5,3 \rangle(2) = 5$

Derived form:

$list_value_expr(value_expr_1) \dots (value_expr_n)$

Example:

$\langle \langle 2,5,3 \rangle, \langle 3 \rangle \rangle(1) = \langle 2,5,3 \rangle$
 $\langle \langle 2,5,3 \rangle, \langle 3 \rangle \rangle(1)(2) = \langle 2,5,3 \rangle(2) = 5$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.6		
Anne Haxthausen, IMM/DTU	Spring 2002	

Built-in List Operators

$\sim : T^* \times T^\omega \rightarrow T^\omega$

$hd : T^\omega \rightarrow T$

$tl : T^\omega \rightarrow T^\omega$

$len : T^\omega \rightarrow \mathbf{Nat}$

$elems : T^\omega \rightarrow \mathbf{T-infset}$

$inds : T^\omega \rightarrow \mathbf{Nat-infset}$

$\langle e_1, \dots, e_n \rangle \sim \langle e_{n+1}, \dots \rangle = \langle e_1, \dots, e_n, e_{n+1}, \dots \rangle$

$hd \langle e_1, e_2, \dots \rangle = e_1$

$tl \langle e_1, e_2, \dots \rangle = \langle e_2, \dots \rangle$


$len \langle e_1, \dots, e_n \rangle = n$

$len \text{ nil} \equiv \mathbf{chaos}$

$elems \langle e_1, e_2, \dots \rangle = \{e_1, e_2, \dots\}$

$inds \text{ fl} = \{1 .. len \text{ fl}\}$

$inds \text{ il} = \{\text{idx} \mid \text{idx} : \mathbf{Nat} \cdot \text{idx} \geq 1\}$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.7		
Anne Haxthausen, IMM/DTU	Spring 2002	

Defining Infinite Lists

value

$all_natural_numbers : \mathbf{Nat}^\omega$

axiom

$inds \text{ all_natural_numbers} = \{n \mid n : \mathbf{Nat} \cdot n > 0\}$,


$all_natural_numbers(1) = 0$,

$\forall \text{idx} : \mathbf{Nat} \cdot$

$\text{idx} \geq 2 \Rightarrow$

$all_natural_numbers(\text{idx}) =$

$all_natural_numbers(\text{idx} - 1) + 1$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.8		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Example 1

scheme QUEUE =

class

type

Element,

Queue = Element*

value

$empty : \text{Queue} = \langle \rangle$,

$put : \text{Element} \times \text{Queue} \rightarrow \text{Queue}$


$put(e,q) \equiv q \sim \langle e \rangle$,

$get : \text{Queue} \rightarrow \text{Queue} \times \text{Element}$

$get(q) \equiv (tl \ q, hd \ q)$

pre $q \neq empty$

end

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.9		
Anne Haxthausen, IMM/DTU	Spring 2002	


List Example 2

```

scheme LIST_PROPERTIES =
  class
  value
    is_permutation : Int* × Int* → Bool,
    is_permutation(l1,l2) ≡
      (∀ i : Int •
        card {idx | idx : Nat •
          idx ∈ inds l1 ∧ l1(idx) = i}
        =
        card {idx | idx : Nat •
          idx ∈ inds l2 ∧ l2(idx) = i}),

    is_sorted : Int* → Bool
    is_sorted(l) ≡
      (∀ idx1,idx2 : Nat •
        {idx1,idx2} ⊆ inds l ∧ idx1 < idx2 ⇒
          l(idx1) ≤ l(idx2))
  end

```


02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.10		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Example 2 – continued

```

scheme SORTING =
  extend LIST_PROPERTIES with
  class
  value
    sort : Int* → Int*
    sort(l) as l1 post is_permutation(l1,l) ∧ is_sorted(l1)
  end

```

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.11		
Anne Haxthausen, IMM/DTU	Spring 2002	

List Example 2 – continued


```

scheme SORTING =
  class
  value
    is_permutation : Int* × Int* → Bool,
    is_permutation(l1,l2) ≡
      (∀ i : Int •
        card {idx | idx : Nat •
          idx ∈ inds l1 ∧ l1(idx) = i}
        =
        card {idx | idx : Nat •
          idx ∈ inds l2 ∧ l2(idx) = i}),

    is_sorted : Int* → Bool
    is_sorted(l) ≡
      (∀ idx1,idx2 : Nat •
        {idx1,idx2} ⊆ inds l ∧ idx1 < idx2 ⇒
          l(idx1) ≤ l(idx2))

  value
    sort : Int* → Int*
    sort(l) as l1 post is_permutation(l1,l) ∧ is_sorted(l1)
  end

```


02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.12		
Anne Haxthausen, IMM/DTU	Spring 2002	

Example 3: Database Representations

- **type** Database = (Key × Data)**-set**

$$\{(k_1, d_1), (k_2, d_2)\}$$
- **type** Database = (Key × Data)*

$$\langle (k_1, d_1), (k_2, d_2) \rangle$$

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.13		
Anne Haxthausen, IMM/DTU	Spring 2002	


Example 3: Database Keys & Data

```

scheme KEY =
  class
    type
      Key
    value
      less_than : Key × Key → Bool
    axiom
      [anti_reflexive]
        ∀ k : Key • ~less_than(k,k),
      [transitive]
        ∀ k1,k2,k3 : Key •
          less_than(k1,k2) ∧ less_than(k2,k3) ⇒ less_than(k1,k3),
      [total_order]
        ∀ k1,k2,k3 : Key •
          less_than(k1,k2) ∨ less_than(k2,k1) ∨ k1 = k2
    end

scheme DATA =
  class
    type Data
    value not_found : Data
  end

```

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.14		
Anne Haxthausen, IMM/DTU	Spring 2002	

Example 3: Database Records


```

scheme RECORD =
  extend KEY with extend DATA with
  class
    type
      Record = Key × Data
    value
      new_record : Key × Data → Record
      new_record(k,d) ≡ (k,d),

      key_of : Record → Key
      key_of(k,d) ≡ k,

      data_of : Record → Data
      data_of(k,d) ≡ d
    end

```


02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.15		
Anne Haxthausen, IMM/DTU	Spring 2002	

Example 3: Database Modeled as Lists

```

scheme LIST_DATABASE =
  extend RECORD with
  class
    type
      Database = { | rl : Record* • is_wf_Database(rl) | }
    value
      is_wf_Database : Record* → Bool
      is_wf_Database(rl) ≡
        (∀ r1,r2 : Record, left,right : Record* •
          rl = left ~ ⟨r1,r2⟩ ~ right ⇒
            less_than(key_of(r1),key_of(r2))),
    end

```

02262: Formal Aspects of Software Engineering I		 Technical University of Denmark Informatics and Mathematical Modelling Computer Science and Technology
Lists; ch.9, pp.63-73		
Foil 6.16		
Anne Haxthausen, IMM/DTU	Spring 2002	

Example 3: LIST_DATABASE continued

```

empty : Database = ⟨ ⟩,

insert : Key × Data × Database → Database
insert(k,d,db) as db1
  post elems db1 =
    (elems remove(k,db)) ∪ {new_record(k,d)},

remove : Key × Database → Database
remove(k,db) ≡ ⟨ r | r in db • key_of(r) ≠ k ⟩,

defined : Key × Database → Bool
defined(k,db) ≡
  if db = ⟨ ⟩ ∨ less_than(k,key_of(hd db)) then false
  else key_of(hd db) = k ∨ defined(k,tl db)
  end,

lookup : Key × Database → Data
lookup(k,db) ≡
  if db = ⟨ ⟩ ∨ less_than(k,key_of(hd db))
  then not_found
  elseif key_of(hd db) = k then data_of(hd db)
  else lookup(k,tl db) end
end

```

Sets, Lists and Products

	number of elements	type of elements	duplicates	ordered
T-set	varying	same	no	no
T-infset	varying	same	no	no
T^*	varying	same	yes	yes
T^ω	varying	same	yes	yes
$T_1 \times \dots \times T_n$	fixed (n)	varying	yes	yes