

ACK Protocol (... Maybe Semantics...)

----- SENDER -----

Sender = wait for input $x \in M$;
Q(x)

Q($x \in M$) = send x to Receiver;
wait until receipt of msg ACK from Receiver;
Sender

----- RECEIVER -----

Receiver = wait until receipt of msg x from Sender;
send ACK to Sender;
output x;
Receiver

ACK/NACK Protocol

----- SENDER -----

Sender = wait for input $x \in M$;
 $Q(x)$

$Q(x \in M)$ = send x to Receiver;
 wait until receipt of msg y from Receiver;
 if $y == \text{ACK}$ then Sender
 else if $y == \text{NACK}$ then $Q(x)$

Receiver

Receiver = wait until receipt of msg x from Sender;
 if $x \in M$ then {
 send ACK to Sender;
 output x ;
 Receiver}
 else if $x \in M'$ then {
 send NACK to Sender;
 Receiver}

Polling protocol

----- RECEIVER -----

Receiver = send msg POLL to Sender;
Wait

Wait = wait until receipt of msg x from Sender;
if $x \in M$ then {
 output x;
 Receiver}
else if $x \in M'$ then {
 send REPT to Sender;
 Wait}

----- SENDER -----

Sender = Q(a) // a is a constant

Q($x \in M$) = wait until receipt of msg y from Receiver;
if $y == \text{POLL}$ then {
 wait for input $x \in M$ from user;
 send x to Receiver;
 Q(x)}
else if $y == \text{REPT}$ then {
 send x to Receiver;
 Q(x)}

ACK/NACK + TIMEOUT (... At-Least-Once Semantics...)

----- SENDER -----

```
Sender = wait for input  $x \in M$ ;  
        Q(x)  
  
Q( $x \in M$ ) = send x to Receiver;  
            start timer(s); // set timeout after s seconds  
            wait until (receipt of msg y from Receiver OR timeout);  
            if receipt of msg y from Receiver then {  
                if  $y == \text{ACK}$  then {  
                    reset timer;  
                    Sender}  
                else if  $y == \text{NACK}$  then {  
                    reset timer;  
                    Q(x)}  
            }  
            else { // timeout  
                reset timer;  
                Q(x)}
```

----- RECEIVER -----

```
Receiver = wait until receipt of msg x from Sender;  
          if  $x \in M$  then {  
              send ACK to Sender;  
              output x;  
              Receiver}  
          else if  $x \in M'$  then {  
              send NACK to Sender;  
              Receiver}
```

PAR Protocol (ACK + TIMEOUT + NUMBERING SCHEME)

----- SENDER -----

Sender = S(1)

$S(n \in N_0)$ = wait for input $x \in M$;
Q(n, x)

$Q(n \in N_0, x \in M)$ = send (n, x) to Receiver;
start timer(s);
wait until (receipt of msg y from Receiver OR timeout);
if receipt of msg y == ACK from Receiver then {
 reset timer;
 S(succ(n))}
else { // timeout
 reset timer;
 Q(n, x)}

----- RECEIVER -----

Receiver = R(0)

$R(n \in N_0)$ = wait until receipt of msg (i, x) from Sender;
if $x \in M$ then {
 if $i = \text{succ}(n)$ then { // expected msg
 send ACK to Sender;
 output x;
 R(succ(n))
 }
 elseif $i = n$ then { // same msg
 send ACK to Sender;
 R(n)}
 else R(n)}
else if $x \in M'$ then R(n)

ACK/NACK + TIMEOUT + NUMBERING SCHEME (... At-Most-Once Semantics...)

... Exercise...

PAR Protocol + TIMEOUT + NUMBERING SCHEME + NUMBERED ACK

----- SENDER -----

Sender = S(1)

$S(n \in N_0)$ = wait for input $x \in M$;
Q(n, x)

$Q(n \in N_0, x \in M)$ = send (n, x) to Receiver;
start timer(s);
wait until receipt of msg a from Receiver OR timeout;
if receipt of msg $a \in N_0$ from Receiver then {
 if (a == n) then { // expected ACK
 reset timer;
 S(succ(n))}
 else { // different ACK
 reset timer;
 Q(n, x)}}
else { // timeout or $a \notin N_0$
 reset timer;
 Q(n, x)}

----- RECEIVER -----

Receiver = R(0)

$R(n \in N_0)$ = wait until receipt of msg (i, x) from Sender;
if $x \in M$ then {
 if $i = \text{succ}(n)$ then { // expected msg
 send i to Sender;
 output x;
 R(i)}
 else {
 send n to Sender;
 R(n)}}
else if $x \in M'$ then R(n)

Polling protocol + TIMEOUT + NUMBERING SCHEME

... Exercise...

Two-Way Exchange (or Handshake) Protocol

Req: requests;
Accept: positive replies;
Refuse: negative replies;
Accept and Refuse are DISJOINT
ref, ERROR \in Refuse (internal message indicating refusal).

At (. . .), both parties are sufficiently finished to go on with the next part of their tasks.

----- SENDER -----

```
Sender = wait for input x  $\in$  Req;  
        send x to Receiver;  
        start timer(s);  
        SR  
  
SR      = wait until receipt of msg y from Receiver OR timeout;  
        if receipt of msg y from Receiver then {  
            if (y  $\in$  Accept) then { // y  $\in$  Accept  
                reset timer;  
                output y;  
                (...)}  
            else { // y  $\in$  Refuse  
                reset timer;  
                output y;  
                Sender}}  
        else { // timeout  
            reset timer;  
            output ERROR; // ERROR  $\in$  Refuse  
            Sender}
```

----- RECEIVER -----

```
Receiver = wait until receipt of msg  $y \in \text{Req}$  from Sender;
           output  $y$ ;
           start timer(s);
           RR

RR = wait for input  $x$  OR timeout;
    if receipt of input  $x$  then {
        if ( $x \in \text{Accept}$ ) then { //  $x \in \text{Accept}$ 
            reset timer;
            send  $x$  to Sender;
            (...)}
        else { //  $x \in \text{Refuse}$ 
            reset timer;
            send  $x$  to Sender;
            Receiver}}
    else { // timeout
        reset timer;
        output ERROR; // ERROR  $\in \text{Refuse}$ 
        send ERROR to Sender;
        Receiver}
```

Three-Way Handshake Protocol

----- SENDER -----

```
Sender      =   wait for input  $r \in \text{Req}$ ;  
              send  $(x, r)$  to Receiver;           //  $x \in \text{TypeOK}$   
              start timer(s);  
              SR(x)  
  
SR( $x \in \text{TypeOK}$ )      =   wait until receipt of msg  $y$  from Receiver OR timeout;  
                          if receipt of msg  $y = (p, q, c)$  from Receiver then {  
                              if  $((c \in \text{Accept}) \ \& \ (p, q \in \text{TypeOK}))$  then {  
                                  reset timer;  
                                  if  $(p = x)$  then {  
                                      send  $(p, q, \text{check})$  to Receiver;  
                                      output  $c$ ;  
                                      (...)}  
                                  else {  
                                      output ERROR;  
                                      Sender}}  
                              elseif  $((c \in \text{Refuse}) \ \& \ (p, q \in \text{TypeOK}))$  then {  
                                  if  $(p = x)$  then {  
                                      reset timer;  
                                      output  $c$ ;  
                                      Sender}  
                                  else  
                                      SR(x)}}  
                          else {  
                              reset timer;           // timeout  
                              output ERROR;         // ERROR  $\in$  Refuse  
                              Sender}
```

----- RECEIVER -----

Receiver = wait until receipt of msg $y = (x \in \text{TypeOK}, r \in \text{Req})$ from Sender;
output r ;
start timer(s);
RR(x)

RR ($x \in \text{TypeOK}$) = wait for input c or timeout;
if receipt of input c then {
 if ($c \in \text{Accept}$) then {
 send (x, y, c) to Sender; // $y \in \text{TypeOK}$
 RC(x, y)
 } elseif ($c \in \text{Refuse}$) then {
 send (x, y, c) to Sender;
 Receiver}}else { // timeout
 reset timer;
 send (x, y, ERROR) to Sender // $\text{ERROR} \in \text{Refuse}$
 output ERROR;
 Receiver}

RC ($x, y \in \text{TypeOK}$) = wait until receipt of msg $y = (p \in \text{TypeOK}, q \in \text{TypeOK}, c \in \text{Check})$ from Sender OR timeout;
if receipt of msg $y = (p, q, c)$ then {
 if ($p = x$) & ($q = y$) then {
 reset timer;
 (...)}
 } else RC(x, y)
else { // timeout
 reset timer;
 output ERROR;
 Receiver}