

Software Engineering I (02161)

Week 10

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Last Week

- ▶ Layered Architecture: Persistent Layer
- ▶ Software Development Processes
 - ▶ Waterfall
 - ▶ (Rational) Unified Process
 - ▶ Agile Processes: User story driven, travel light, Agile Manifesto

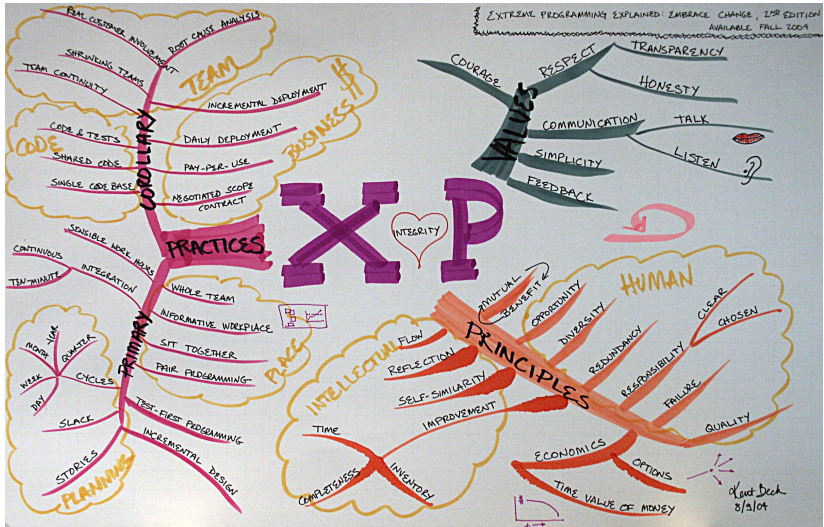
Contents

Software Development Process

Project planning

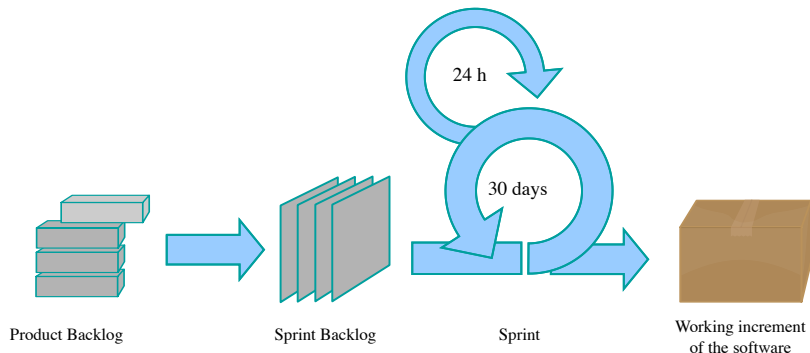
Design by Contract (DbC)

eXtreme Programming (XP)



Kent Beck, Extreme Programming 2nd ed.

Scrum



Wikipedia

- ▶ Robert Martin (Uncle Bob) about "The Land that Scrum Forgot"

<http://www.youtube.com/watch?v=hG4LH6P8Syk>

- History about agile methods, the agile manifesto, and Scrum and its relationship to XP

Lean Software Development

- ▶ Lean Production.

- ▶ Reduce the amount of *waste* in the production process
 - ▶ Generate *flow*

- ▶ Waste: resources used which do not produce value for the customer

- ▶ time needed to fix bugs
 - ▶ time to change the system because it does not fit the customers requirements
 - ▶ time waiting for approval
 - ▶ ...

Cycle time

Cycle time

Time it takes to go through the process one time

$$cycle_time = \frac{number_of_features \downarrow}{feature_implemantion_rate \uparrow}$$

► Example: Waterfall

- Batch size = number_of_features in an iteration
- Software: 250 features, feature_implementation_rate = 5 features/week
- $cycle_time = 250 \text{ f} / (5 \text{ f/w}) = \underline{50 \text{ weeks}}$
- Overall time: 50 weeks

→ 1 cycle

Goal: Reducing the cycle time

- ▶ Reduce batch size: 1 feature in an iteration
- ▶ Software: 250 features, $\text{feature_implementation_rate} = 5$ features/week

$$\text{cycle_time} = \frac{\text{number_of_features}}{\text{feature_implemantion_rate}}$$

- ▶ Agile: $\text{cycle_time} = 1 \text{ f} / (5 \text{ f/w}) = 1/5 \text{ week} = \underline{1 \text{ day}} = \underline{8 \text{ h}}$

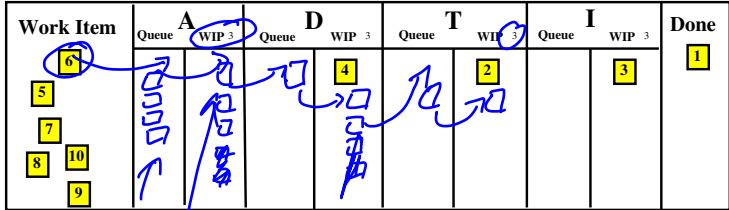
→ 250 cycles

- ▶ Advantages
 - ▶ Process adapts to changes in requirements
 - ▶ Process improvements and fine tuning

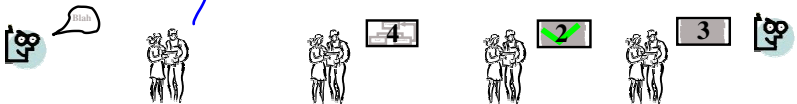
Generating flow using Pull and Kanban

WIP = Work in Progress Limit

highest
priority



bottleneck



Flow through Pull with Kanban



- ▶ Process controlling: local rules
- ▶ Load balancing: Kanban cards and *Work in Progress (WIP) limits*
- ▶ Integration in other processes

Online Kanban Tool: Trello

- ▶ `www.trello.com`: Electronic Kanban board useful for your project
- ▶ Example Kanban board `https://trello.com/b/4wddd1zf/kanban-workflow`

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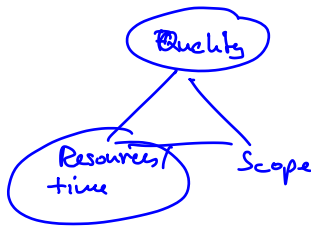
Software Development Process

Project planning

Design by Contract (DbC)

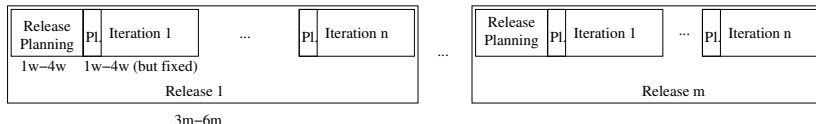
Project Planning

- ▶ Project plan
 - ▶ Defines:
 - ▶ How work is done
 - ▶ Estimate
 - ▶ Duration of work
 - ▶ Needed resources
 - Price
- ▶ Project planning
 - ▶ Proposal stage
 - Price
 - Time to finish
 - ▶ Project start-up
 - Staffing, ...
 - ▶ During the project
 - ▶ Progress (tracking)
 - ▶ Adapt to changes



Planning Agile Projects

- ▶ *fixed* general structure
- e.g. quarterly cycle / weekly cycle practices in XP / sprints in Scrum



- ▶ *time boxing*
 - ▶ fixed: release dates and iterations
 - ▶ adjustable: scope
- ▶ Planning: Which user story in which iteration / release

Planning game

- ▶ Goal of the game:
 - ▶ List of prioritized user stories
- ▶ Customer defines:
 - ▶ user stories
 - ▶ priorities
- ▶ Developer define:
 - ▶ costs, risks
 - ▶ suggest user stories
- ▶ Customer decides: is the user story worth its costs?
 - split a user story
 - change a user story

Scrum/XP: Project estimation and monitoring

- ▶ Estimation: two possibilities

- 1) Estimate *ideal time* (e.g. person days / week) * *load_factor*
- 2) Estimate *relative* to other user stories: *story points*

- ▶ **Monitoring**

- ad 1) New *load factor*: $\text{total_iteration_time} / \text{user_story_time finished}$

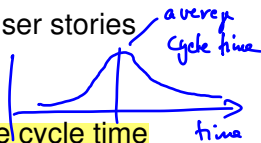
- ad 2) *velocity*: Number of points per iteration

→ What can be done in the next iteration

- ▶ *Yesterdays weather*: Calculate velocity/load_factor based on the *last* iteration only
- ▶ Important: If in trouble focus on *few stories* and finish them

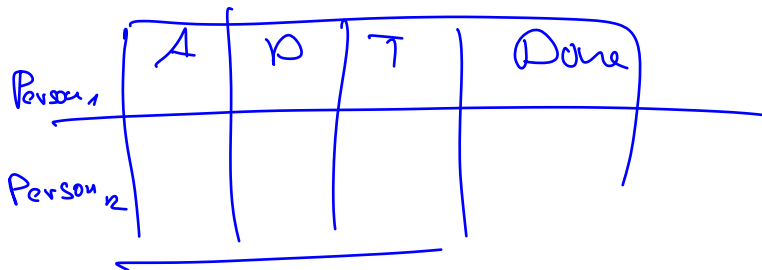
Lean / Kanban: User story estimation

- ▶ No "iterations": user stories come in and flow through the system
- Only a rough estimation of the size of the user stories
 - ▶ try to level the size of the user stories
 - ▶ Divide larger into smaller ones
- ▶ Measure process parameters, e.g., average cycle time
 - ▶ E.g. "After committing to a user story, it takes in average a week to have the user story finished"
- ▶ User average_cycle_time and WIP (Work In Progress) Limit to determine the capacity of the process and thus throughput



Example of a Kanban board for the exam project

- <https://trello.com/b/i029C07w/02161-example>



Contents

Software Development Process

Project planning

Design by Contract (DbC)

- Contracts

- Implementing DbC in Java

- Assertion vs Tests

- Inheritance

- Invariants

- Defensive Programming

What does this function do?

↪ qsort

sorting

Quicksort

```
public List<Integer> f(List<Integer> list) {  
    if (list.size() <= 1) return list;  
  
    int p = list.elementAt(0);  
  
    List<Integer> l1 = new ArrayList<Integer>();  
    List<Integer> l2 = new ArrayList<Integer>();  
    List<Integer> l3 = new ArrayList<Integer>();  
  
    g(p, list, l1, l2, l3);  
  
    List<Integer> r = f(l1);  
  
    r.addAll(l2);  
    r.addAll(f(l3));  
  
    return r;  
}  
  
public void g(int p, List<Integer> list,  
              List<Integer> l1, List<Integer> l2, List<Integer> l3) {  
    for (int i : list) {  
        if (i < p) l1.add(i);  
        if (i == p) l3.add(i);  
        if (i > p) l2.add(i);  
    }  
}
```

What does this function do?

```
public void testEmpty() {
    int[] a = {};
    List<Integer> r = f(Array.asList(a));
    assertTrue(r.isEmpty());
}

public void testOneElement() {
    int[] a = { 3 };
    List<Integer> r = f(Array.asList(a));
    assertEquals(Array.asList(3), r);
}

public void testTwoElements() {
    int[] a = {2, 1};
    List<Integer> r = f(Array.asList(a));
    assertEquals(Array.asList(1, 2), r);
}

public void testThreeElements() {
    int[] a = {2, 3, 1};
    List<Integer> r = f(Array.asList(a));
    assertEquals(Array.asList(1, 2, 3), r);
}

...
```

What does this function do?

List<Integer> f(List<Integer> a)

Precondition: *a* is not **null**

Postcondition: For all *result*, $a \in \text{List<Integer>}$:

$\text{result} == f(a)$

if and only if

isSorted(result) and sameElements(a,result)

where

isSorted(a) if and only if

for all $0 \leq i, j < a.size()$:

$i \leq j$ implies $a.get(i) \leq a.get(j)$

and

sameElements(a,b) if and only if

for all $i \in \text{Integer}$: $\text{count}(a, i) = \text{count}(b, i)$

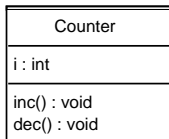
$$\left[\begin{array}{l} 3 \leq 10 \\ i \leq j \end{array} \right]$$

Example Counter

formal specification

```
{context Counter :: dec ()  
pre: i > 0  
post: i = i@pre - 1 }
```

contract for dec



```
{context Counter  
inv: i >= 0}
```

```
{context Counter :: inc ()  
post: i = i@pre + 1 }
```

pre: true

contract for inc.

client

```
public T n(Tl a1, ..., Tn an, Counter c)  
...  
// Here the precondition of c has to hold  
// to fulfil the contract of Counter::dec  
c.dec();  
// Before returning from dec, c has to ensure the  
// postcondition of dec  
...
```

Design by contract

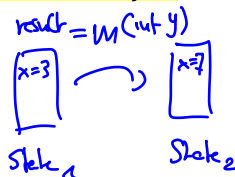
- ▶ Name invented by Bertrand Meyer (Eiffel programming language) for pre-/post-condition based formal methods applied to object-oriented designs/languages
- ▶ Pre-/post-conditions were invented by Tony Hoare and Robert W. Floyd

Contract for a method

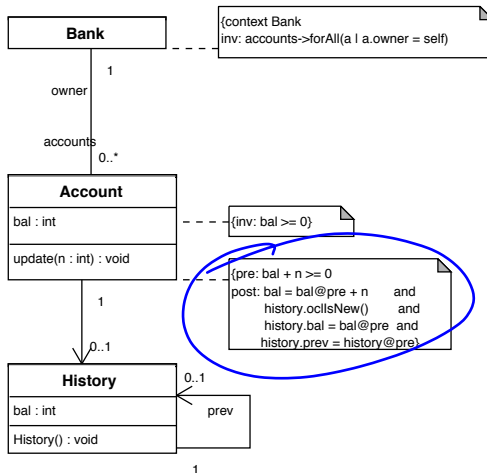
- ▶ precondition: a **boolean expression** over the **state** of the object and **arguments** *before* the execution of the method
- ▶ postcondition: a **boolean expression** over the state of the object and **arguments** *before the execution* of a method *and* the **result of the method** and the **state of the object** *after* the execution of the method

Contract between Caller and the Method

- ▶ Caller ensures precondition
- ▶ Method ensures postcondition



Bank example with constraints



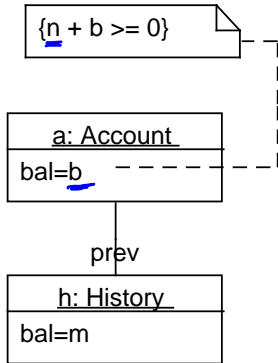
Update operation of Account

pre condition

```
{pre: bal + n >= 0  
post: bal = bal@pre + n    and  
      history.ocllsNew()    and  
      history.bal = bal@pre and  
      history.prev = history@pre}
```

State before executing

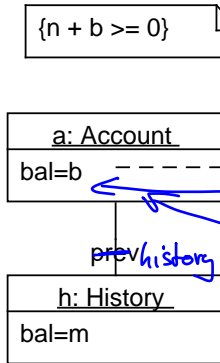
update (n)



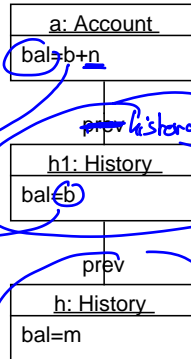
Update operation of Account

```
{pre: bal + n >= 0  
post: bal = bal@pre + n and  
      history.ocllsNew() and  
      history.bal = bal@pre and  
      history.prev = history@pre}
```

State **before** executing
update (n)



State **after** executing
update (n)



Example

Object Constraint Language

OCL expression for set union

formal language for UML

```
LibraryApp::addMedium(Medium m)
pre: adminLoggedIn
post: medium = medium@pre->including(m) and
      medium.library = this
```

```
LibraryApp::search(String string) : List<Medium>
post: result = medium->select(m |
    m.title.contains(string) or
    m.author.contains(string) or
    m.signature.contains(string))
medium = medium@pre
```

```
User::borrowMedium(Medium m)
pre: borrowedMedium->size < 10
      and m != null
      and not(borrowedMedium->exists(m' | m'.isOverdue))
post: m.borrowDate = libApp.getDate() and
      borrowedMedium = borrowedMedium@pre->including(m)
```

Implementing DbC with assertions

- ▶ Many languages have an assert construct. In Java:

`assert bexp;` or `assert bexp:string;`

- ▶ Contract for Counter::dec(i:int)

Pre: $i > 0$

Post: $i = i@pre - 1$

```
pre → void dec() {  
      → assert i > 0 : "Precondition violated"; // Precondition  
      → int iatpre = i; // Remember the value of the counter  
                          // to be used in the postcondition  
      i--;  
post → assert i == iatpre-1 : "Postcondition violated"; // Postcondition  
      }
```

- ▶ assert and assertTrue are not the same!

↑ ↳ JUnit used in a test
assertion used inside a method.

Implementing DbC in Java

Pre: *args* \neq *null* and *args.length* > 0

Post: $\forall n \in \text{args} : \text{min} \leq n \leq \text{max}$

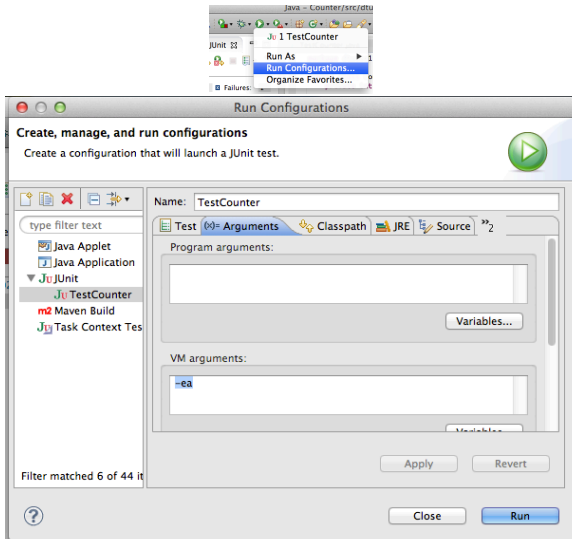
```
public class MinMax {
    int min, max;

    public void minmax(int[] args) throws Error {
        assert args != null && args.length != 0;
        min = max = args[0];
        for (int i = 1; i < args.length; i++) {
            int obs = args[i];
            if (obs > max)
                max = obs;
            else if (min < obs)
                min = obs;
        }
        assert isBetweenMinMax(args);
    }

    private boolean isBetweenMinMax(int[] array) {
        boolean result = true;
        for (int n : array) {
            result = result && (min <= n && n <= max);
        }
        return result;
    }
}
```

Important

- ▶ Assertion checking is switched off by default in Java
 - 1) Use `java -ea Main` to enable assertion checking
 - 2) In Eclipse



Assertions

- ▶ Advantage

- ▶ Postcondition is checked for each computation
- ▶ Precondition is checked for each computation

- ▶ Disadvantage

- ▶ Checking that a postcondition is satisfied can take as much time as computing the result

→ Performance problems

- ▶ Solution:

- ▶ Assertion checking is switched on during developing, debugging and testing and switched off in production systems

Assertion vs. Tests

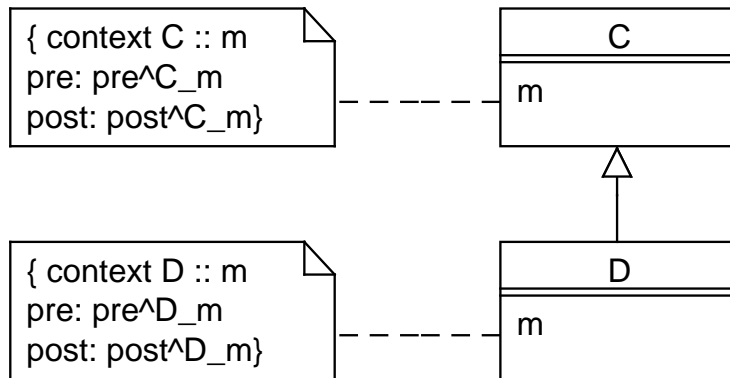
- ▶ Assertion

- ▶ Checks all computations (as long as assertion checking is switched on)
- ▶ Checks also for contract violations from the client (i.e. precondition violations)

- ▶ Tests

- ▶ Only checks test cases (concrete values)
- ▶ Cannot check that the clients establish the precondition

Contracts and inheritance



Contracts and Inheritance

Liskov / Wing Substitution principle:

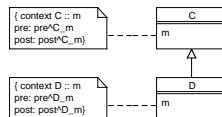
At every place, where one can use objects of the superclass C, one can use objects of the subclass D

```
public T n(C c)
...
// n has to ensure  $\text{Pre}^C_m$ 
c.m();
// n can rely on  $\text{Post}^C_m$ 
...
```

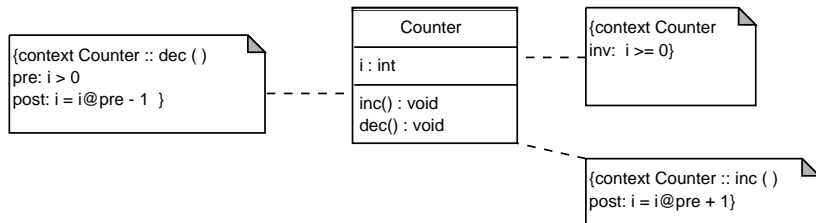
$t.n(\text{new } C())$ vs. $t.n(\text{new } D())$.

→ $\text{Pre}^C_m \implies \text{Pre}^D_m$ weaken precondition

→ $\text{Post}^D_m \implies \text{Post}^C_m$ strengthen postcondition



Invariants: Counter



- ▶ **Methods**
 - ▶ assume that invariant holds
 - ▶ ensure invariants
- ▶ **When does an invariant hold?**
 - ▶ After construction
 - ▶ After each *public* method

Invariants

- ▶ Constructor has to ensure invariant

```
public Counter() {  
    i = 0;  
    assert i >= 0; // Invariant  
}
```

- ▶ Operations ensure and assume invariant

```
void dec() {  
    assert i >= 0; // Invariant  
    assert i > 0; // Precondition  
    int iatpre = i; // Remember the value of the counter  
                  // to be used in the postcondition  
    i--;  
    assert i == iatpre-1; // Postcondition  
    assert i >= 0; // Invariant  
}
```

Defensive Programming

- ▶ Can one trust the client to ensure the precondition?

```
void dec() { i--; }
```

- ▶ Depends if the programmer controls the client or not
 - ▶ e.g. if dec is private, only the programmer of the method can call dec
 - ▶ if dec is public, potentially others can call the method

Defensive Programming

- ▶ If one does not trust the client
- ▶ Check explicitly that the precondition of a method is satisfied

- ▶ Either

```
void dec() { if (i > 0) { i--; } }
```

- ▶ Or

```
void dec() {  
    if (i <= 0) {  
        throw new Exception("Dec not allowed ...");  
    }  
    i--;  
}
```

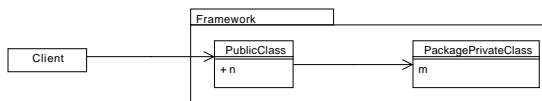
- ▶ Don't rely on the `assert` statement.

- ▶ Why?

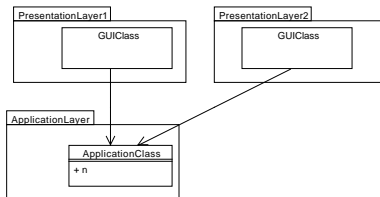
```
void dec() {  
    assert i <= 0;  
    i--;  
}
```

Defensive Programming

- ▶ Use defensive programming with public methods
- ▶ Use asserts with private or package private methods
- ▶ For example public method of a library



- ▶ Public method of a class in the application/domain layer



Next week

- ▶ Principles of Good Design
- ▶ Design Patterns