Solutions for CP Exercises, September 15

1. Solution for Andrews Ex. 3.2

```
var l : integer := 1;

process P[i : 1..n] =
  var s : integer;
  repeat
    non-critical section_i;
    DEC(l, s);
    while s > 0 do {
      INC(r, l);
      delay;
      DEC(l, s);
    }
    critical section_i;
    INC(l, s);
  forever;
```

Here, the lock l is used as in the test-and-set solution. However, if the lock is already "set" ($l < 1$), the effect of *DEC* must be undone by *INC*, before trying again. The correctness argument (or proof) follows the same line as for the test-and-set solution.

2. Solution for Andrews Ex. 3.13

A first attempt to be able to "use the barrier again" would be to let the last processes that arrives reset the counter thereby releasing everybody:

```
var count : integer := 0;

process Worker[i : 1..n] =
  repeat
    code to implement task i;
    ⟨ count := count + 1 ⟩;
    if count = n then count := 0;
    ⟨ await count = 0 ⟩;
  forever;
```

In this proposal, however, the processes first released by the reset may race to the next barrier round and increment *count* before all processes have seen $count = 0$. The result is a deadlock.

The solution is to use two counters ensuring that the first counter is not incremented until everybody have seen the reset and vice versa:

```
var count_1, count_2 : integer := 0;
```

```
process Worker[i : 1..n] =
  repeat
    code to implement task i;
    ⟨ count1 := count1 + 1 ⟩;
    if count1 = n then count1 := 0;
    ⟨ await count1 = 0 ⟩;
    ⟨ count2 := count2 + 1 ⟩;
    if count2 = n then count2 := 0;
    ⟨ await count2 = 0 ⟩;
  forever;
```

Using the *fetch-and-add* instruction, this can be readily implemented:

```
process Worker[i : 1..n] =
  repeat
    code to implement task i;
    if FA(count1, 1) = n − 1 then count1 := 0;
    while count1 ≠ 0 do skip;
    if FA(count2, 1) = n − 1 then count2 := 0;
    while count2 ≠ 0 do skip;
  forever;
```