# **Lecture 6: Discovering Domain Entities**

## 5. Discovering Domain Entities

- Lecture 2 briefly characterised, informally and also a bit more formally, what we mean by a domain.
- Lecture 3 informally and systematically characterised the four categories of entities: parts, actions, events and behaviours.
- Lectures 4–5 more-or-less "repeated" Sect. 3's material but by now giving more terse narratives (that is, informal descriptions) and, for the fist time, also formalisations.
- Lectures 4–5 did not hint at how one discovers domain parts (i.e., their types), actions, events and behaviours.
- In this section we try unravel a set of techniques and tools so-called 'discoverers' and 'analysers' using which the domain describer (scientist and/or engineer) can more-or-less systematically discover, analyse and describe a domain, informally and formally.

### 5.1. Preliminaries

• Before we present the discoverers and analysers we need establish some concepts. These are:

- Part Signatures	Slides 223–224
- Domain Indices	Slides 225–226
- Inherited Part Signatures	Slides 227–227
- Domain Signatures	Slides 228–231
- Simple and Compound Domain Signatures	Slides 233–234

## 5.1.1. Part Signatures

- Let us consider a part p:P.
- Let p: P, by definition, be the *principal part* of a domain.
  - Now we need to identify
    - \* (i) the type, P, of that part;
    - \* (ii) the types,  $S_1, \ldots, S_m$ , of its proper sub-parts (if p is composite);
    - \* (iii) the type, PI, of its unique identifier;
    - \* (iv) the possible types,  $MI_1, \ldots, MI_n$ , of its mereology; and
    - \* (v) the types,  $A_1$ , ...,  $A_o$ , of its attributes.
- We shall name that cluster of type identifications
  - the part signature.
- We refer to P as identifying the part signature.
- Each of the  $S_i$  (for  $i : \{1..m\}$ ) identifies sub-parts and hence sub-part, i.e., part signatures.

# **Example 41 (Net Domain and Sub-domain Part Signatures)**

The part signature of the hubs and the links are here chosen to be those of

- (i) **N**,
- (ii) Hs,
- (ii) Ls and
- Net\_name, Net\_owner, etc.

The part signatures Hs and Ls are

- $\bullet$  (ii) Hs = H-set, H
- (iii,iv) HI, LI-set
- (v) Hub\_Nm, Location, H $\Sigma$ , H $\Omega$ , ...
- (ii) Ls = L-set, L
- (iii,iv) LI, HI-set
- (v) Link\_Nm, L $\Sigma$ , L $\Omega$ , LEN, etc.

### 5.1.2. Domain Indices

- By a domain index we mean a list of part type names that identify a sequence of part signatures.
- More specifically
  - The domain  $\Delta$  has index  $\langle \Delta \rangle$ .
  - The sub-domains of  $\Delta$ , with part types A, B, ..., C, has indices  $\langle \Delta, A \rangle, \langle \Delta, B \rangle, ..., \langle \Delta, C \rangle$ .
  - The sub-domains of sub-domain with index  $\ell$  and with part types A, B, ..., C has indices  $\ell^{\hat{}}\langle A \rangle, \ell^{\hat{}}\langle B \rangle, ..., \ell^{\hat{}}\langle C \rangle$ .

**Example 42 (Indices of a Road Pricing Domain)** We refer to the the *Road-pricing Transport Domain*, cf. Example 36 on page 184. The sub-domain indices of the road-pricing transport domain,  $\Delta$ , are:

- $\bullet \langle \Delta \rangle$ ,
- $\bullet \langle \Delta, \mathsf{N} \rangle$ ,
- $\bullet \langle \Delta, \mathsf{F} \rangle$ ,
- $\bullet \langle \Delta, \mathsf{M} \rangle$ ,
- $\langle \Delta, N, Vs \rangle$ ,

- $\langle \Delta, \mathsf{N}, \mathsf{Ls} \rangle$ ,
- $\langle \Delta, N, H \rangle$ ,
- $\bullet \langle \Delta, N, L \rangle$  and
- $\bullet \langle \Delta, F, V \rangle$ .

## 5.1.3. Inherited Domain Signatures

- Let  $\langle \Delta, A, B, C, D \rangle$  be some domain index.
- Then

$$-\langle \Delta, A, B, C \rangle$$
  $-\langle \Delta, A, B \rangle$   $-\langle \Delta, A \rangle$   $-\langle \Delta \rangle$ 

are the inherited domain indices of  $\langle \Delta, A, B, C, D \rangle$ .

# 5.1.4. Domain and Sub-domain Categories

- ullet By the domain category of the domain indexed by  $\ell \widehat{\ } \langle \mathsf{D} \rangle$  we shall mean
  - the domain signature of D,

and the

- action,
- event and
- behaviour

definitions whose signatures involves

- just the types given in the domain signature of D or
- in inherited domain signatures.

# **Example 43 (The Road-pricing Domain Category)** The

road-pricing domain category consist of

- the types N, F and M,
- the create\_Net create\_Fleet and create\_M actions, and
- corresponding Net, Fleet and M behaviours

- By a sub-domain category, of index  $\ell$ , we shall mean
  - the sub-domain types of the sub-domain designated by index  $\ell$ , and the
  - actions,
  - events and
  - behaviours

whose signatures involves just the types

- of the  $\ell$  indexed sub-domain
- or of any prefix of  $\ell$  indexed sub-domain
- or of the root domain.

# **Example 44 (A Hub Category of a Road-pricing Transport Domain)**

The ancestor sub-domain types of the hub sub-domain are: HS, N and  $\Delta$ .

- The hub category thus includes
  - the part (etc.) types H, HI, ...,
  - the insert\_Hub and the delete\_Hub actions,
  - perhaps some saturated\_hub (and/or other) event(s),
  - but probably no hub behaviour as it would involve at least the type LI which is not in an ancestor sub-domain of the Hub sub-domain.

### **5.1.5.** Simple and Compound Indexes

- By a simple index we mean a domain or a sub-domain index.
- By a compound index we mean a set of two or more distinct indices of a domain  $\Delta$ .
- Compound indices,  $c_{idx}$ :  $\{\ell_i, \ell_j, \dots, \ell_k\}$ , designate parts, actions, events and behaviours each of whose types and signatures involve types defined by all of the simple indexes of  $c_{idx}$ .

Example 45 (Compound Indices of the Road-pricing System) We show just one compound index:

•  $\{\langle \Delta, N, HS, H \rangle, \langle \Delta, N, LS, L \rangle \}$ .

# 5.1.6. Simple and Compound Domain Categories

- ullet By a simple domain category we shall mean any  $\ell$ -indexed [sub-]domain category.
- By the compound domain category of compound index  $c_{idx}: \{\ell_i, \ell_j, \dots, \ell_k\}$ , we shall mean
  - the set of types, actions, events and behaviours
  - as induced by compound index  $c_{idx}$ , that is,
  - parts, actions, events and behaviours
  - each of whose types and signatures involve types defined by all of the simple indexes of  $c_{idx}$ .

# **Example 46 (The Compound Domain Category of Hubs and Links)**

The compound domain category designated by  $\{\langle \Delta, N, HS, H \rangle, \langle \Delta, N, LS, L \rangle\}$  includes:

```
mereo_L: L \rightarrow HIs,
type
                                                        mereo H: H \rightarrow LIs
   HIs = HI-set
       axiom ∀ his:HIs·card his=2
                                                        attr H\Sigma: H \to H\Sigma
                                                        attr L\Sigma: L \rightarrow L\Sigma
   LIs = LI-set
                                                        attr H\Omega: H \to H\Omega
   H\Sigma = (LI \times LI)-set
   L\Sigma = (HI \times HI)-set
                                                        attr L\Omega: L\rightarrow L\Omega
   H\Omega = H\Sigma-set
                                                     axiom
   L\Omega = L\Sigma-set
                                                        \forall h: H \cdot attr\_H\Sigma(h) \subseteq attr\_H\Omega(h)
value
                                                        \forall l: L \cdot attr_L \Sigma(l) \subseteq attr_L \Omega(l)
```

## **5.1.7. Examples**

• We repeat some examples, but now "formalised".

**Example 47 (The Root Domain Category)** We start at the root,  $\Delta$ , of the Road Pricing Domain. See Fig. 13.





Figure 13: The  $\langle \Delta \rangle$  Root

- At the root we 'discover' the net, fleet and road pricing monitor.
- See Fig. 14 on the following page.

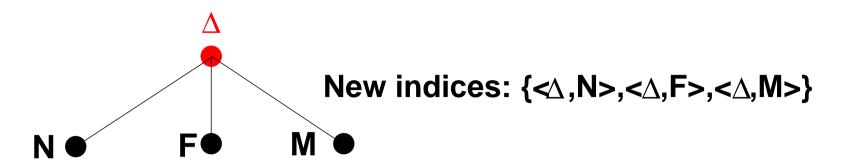


Figure 14: Exploring the root index  $\langle \Delta \rangle$  Index

When observing the very essence of the road pricing domain "at the  $\langle \Delta \rangle$  level" one observes:

```
type
N, F, M
value
obs_N: \Delta \rightarrow N
```

obs\_M:  $\Delta \to M$ 

obs\_F:  $\Delta \rightarrow F$ 

attr...:  $\Delta \rightarrow ...$ 

where ... stands for types of road pricing domain attributes.

J

**Example 48 (The Net Domain Category)** We then proceed to explore the domain at index  $\langle \Delta, N \rangle$ . See Fig. 15.

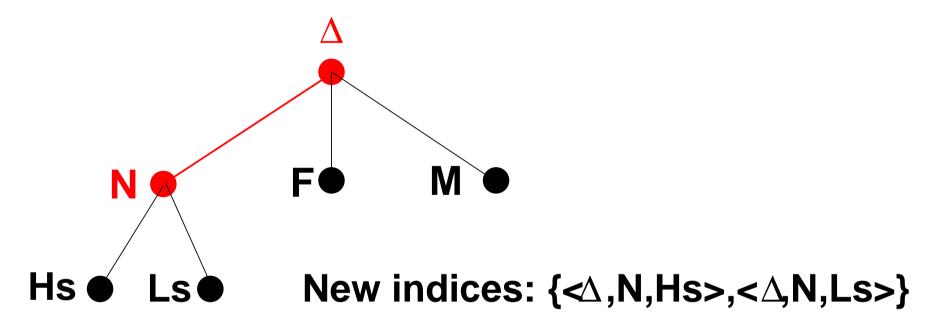


Figure 15: Exploring the  $\langle \Delta, N \rangle$  Index

When observing the very essence of the Net domain, "at the  $\langle \Delta, \mathbf{N} \rangle$  level" one observes:

```
type
Hs = H\text{-set}
Ls = L\text{-set}
H
L
...
value
obs\_Hs: N \to Hs
obs\_Ls: N \to Ls
attr\_Hs: Hs \to ...
attr\_Ls: Ls \to ...
```

where  $\dots$  stand for attributes of the Hs and the Ls parts of N.

**Example 49 (The Fleet Domain Category)** We then proceed to explore the domain at index  $\langle \Delta, F \rangle$ . See Fig. 16.

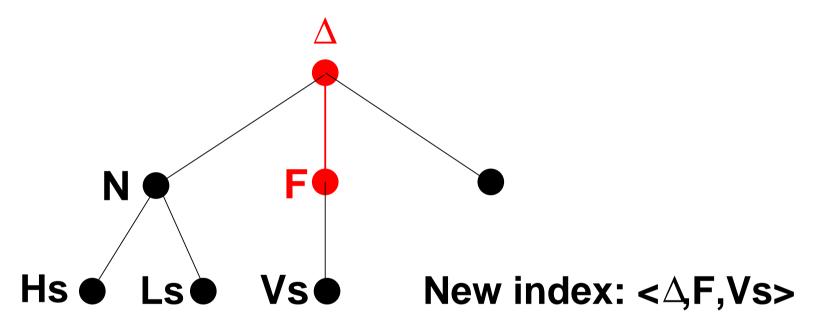


Figure 16: Exploring the  $\langle \Delta, F \rangle$  Index

When observing the very essence of the Fleet domain, "at the  $\langle \Delta, \mathsf{F} \rangle$  level" one observes:

```
type
Vs = V\text{-set}
V
...
value
obs_{-}Vs: F \rightarrow Vs
attr_{-}... Vs \rightarrow ...
```

where ... stand for attributes that we may wish to associate with Fleets of vehicles. **Example 50 (The Hub Domain Category)** We now switch "back" to explore the domain at index  $\langle \Delta, N, Hs \rangle$ . See Fig. 17.

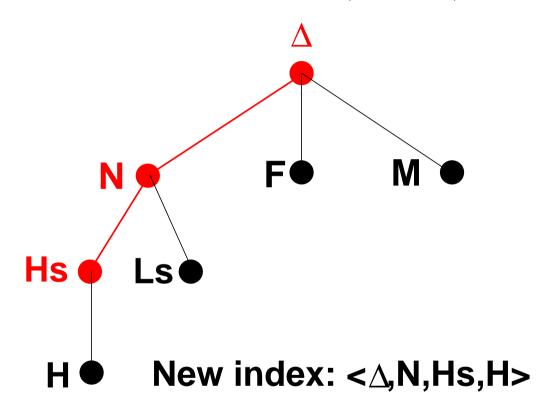


Figure 17: Exploring the  $\langle \Delta, N, Hs, H \rangle$  Index

When observing the very essence of the Fleet domain, "at the  $\langle \Delta, \mathsf{N}, \mathsf{Hs}, \mathsf{H} \rangle$  level" one observes:

```
type
HI
...

value
uid\_HI: H \rightarrow HI
attr\_...: H \rightarrow ...
```

where ... stand for LOCation, etc.

**Example 51 (The Link Domain Category)** Next we explore the link domain. See Fig. 18.

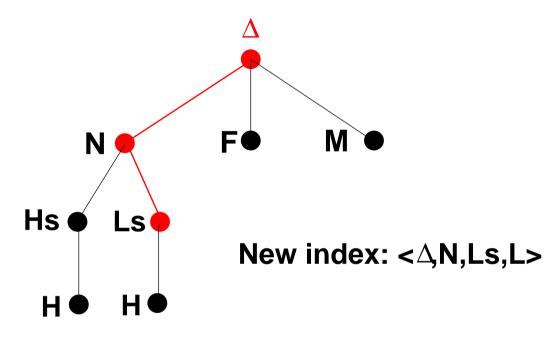


Figure 18: Exloring the  $\langle \Delta, N, Ls, L \rangle$  Index

When observing the very essence of the Fleet domain, "at the  $\langle \Delta, \mathsf{N}, \mathsf{Ls}, \mathsf{L} \rangle$  level" one observes:

```
type
LI
...

value
uid\_LI: L \rightarrow LI
attr\_...: L \rightarrow ...
```

where ... stand for LOCation, LENgth, etc.

# **Example 52 (The Compound Hub and Link Domain Category)**

We next explore a compound domain. See Fig. 19.

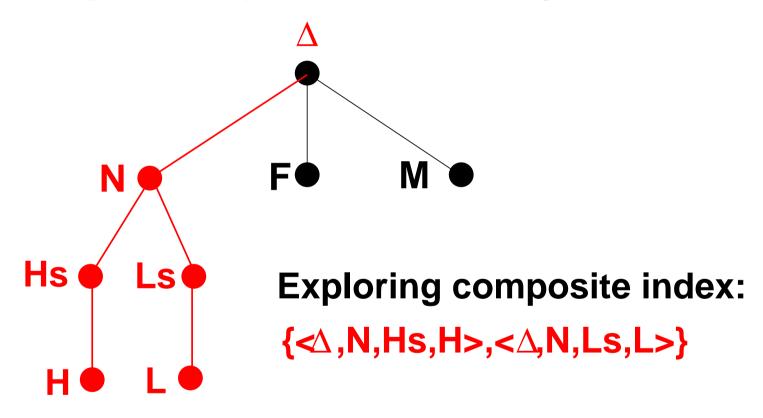


Figure 19: Exploring composite index  $\{\langle \Delta, N, Hs, H \rangle, \langle N, Ls, L \rangle\}$ 

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When observing the very essence of the Fleet domain, at the  $\{\langle \Delta, N, Hs, H \rangle, \langle \Delta, N, Ls, L \rangle\}$  level one observes:

```
type
      H\Sigma = (LI \times LI) - set, H\Omega = H\Sigma - set,
      L\Sigma = (HI \times HI)-set, L\Omega = L\Sigma-set
value
      attr_H\Sigma: H \to H\Sigma, attr_H\Omega: H \to H\Omega
      attr_L\Sigma: L \rightarrow L\Sigma, attr_L\Omega: L \rightarrow L\Omega
      mereo_L: L \rightarrow HI-set axiom \forall l:L:card mereo_L(l)=2
      mereo_H: H \rightarrow LI-set (= LI-set)
      remove_H: HI \rightarrow N \stackrel{\sim}{\rightarrow} N
      insert_L: L \rightarrow N \stackrel{\sim}{\rightarrow} N
      remove_L: LI \rightarrow N \stackrel{\sim}{\rightarrow} N
axiom
      \forall h\sigma: H\Sigma \dots \forall h\omega: H\Omega \dots
      \forall \ l\sigma: H\Sigma \dots, \forall \ l\omega: H\Omega \dots
```

#### 5.1.8. Discussion

- The previous examples (47-52), especially the last one (52),
  - illustrates the complexity of a domain category;
  - from just observing sub-part types and attributes (as in Examples 47–49),
  - Example 52 observations grow to intricate mereologies etcetera.
- The 'discoverers' that we shall propose
  - aim at structuring the discovery process by focusing, in turn, on

```
* part sorts,
* concrete part types,
* unique identifier types of parts,
* event signatures,
* part mereology,
* etc.
```

## 5.2. Proposed Type and Signature 'Discoverers'

- By a 'domain discoverer' we shall understand
  - a tool and a set of principles and techniques
  - for using this tool
  - in the discovery of the entities of a domain.
- In this section we shall put forward a set of type and signature discoverers.
  - Each discoverer is indexed by a simple or a compound domain index.
  - And each discoverer is dedicated to some aspect of some entities.
- Together the proposed discoverers should cover the most salient aspects of domains.
- Our presentation of type and signature discoverer does not claim to help analyse "all" of a domain.

• We need formally define what an index is.

### type

```
Index = Smpl_Idx \mid Cmpd_Idx
     Smpl_Idx = \{ | \langle \Delta \rangle \hat{i}dx | idx:Type_Name^* | \}
     Cmpd_Idx' = Smpl_Idx-set
     Cmpd_Idx = \{ | sis: Cmpd_Idx' \cdot wf_Cmpd_Idx(sis) | \}
value
     wf_Dmpd_Idx: Cmpd_Idx' \rightarrow Bool
     wf_Cmpd_Idx(sis) \equiv \forall si,si':Smpl_Idx \cdot \{si,si'\}\subseteq sis \land si \neq si'
     \mathcal{DISCOVERER\_KIND}: Index \rightarrow Text
     \mathcal{DISCOVER}_{\mathcal{K}I\mathcal{N}\mathcal{D}}(\ell^{\hat{}}\langle t \rangle) as text
          pre: \ell \hat{\ } \langle t \rangle is a valid index beginning with \Delta
          post: text is some, in our case, RSL text
```

• The idea of the  $\ell \hat{\phantom{a}} \langle t \rangle$  index is that it identifies a sub-domain, t, of  $\Delta$ 

— where  $\mathcal{DISCOVERER\_KIND}$  is any of the several different "kinds" of domain forms:

```
[90 (Slide 255)] PART_SORTS,

[91 (Slide 257)] HAS_A_CONCRETE_TYPE,

[92 (Slide 260)] PART_TYPES,

[93 (Slide 264)] UNIQUE_ID,

[96 (Slide 265)] MEREOLOGY,

[98 (Slide 267)] ATTRIBUTES,

[100 (Slide 275)] ACTION_SIGNATURES,

[101 (Slide 276)] EVENT_SIGNATURES and

[103 (Slide 280)] BEHAVIOUR_SIGNATURES.
```

- In a domain analysis (i.e., discovery) the domain description emerges "bit-by-bit".
  - Initially types are discovered and hence texts which define
    - \* unique identifier types and functions,
    - \* mereology types and functions, and
    - \* attribute types and functions.
  - Then the signatures of actions, events and behaviours.

## 5.2.1. Analysing Domain Parts

- The two most important aspects of an algebra are those of
  - its parts and
  - its operations.
- Rather than identifying, that is, discovering or analysing individual parts
  - we focus on discovering their types —
  - initially by defining these as sorts.
- And rather than focusing on defining what the operations achieve
  - we concentrate on the signature,
  - -i.e., the types of the operations.

- It (therefore) seems wise to start with the discovery of parts, and hence of their types.
  - Part types are present in the signatures of all actions, events and behaviours.
  - When observing part types we also observe a variety of part type analysers:
    - \* possible unique identities of parts,
    - \* the possible mereologies of composite parts, and
    - \* the types of the attributes of these parts.

### 5.2.1.1. Domain Part Sorts and Their Observers

- Initially we "discover" parts
  - by deciding upon their types,
  - in the form, first of sorts,
  - subsequently and possibly in the form of concrete types.

### 5.2.1.1.1. A Domain Sort Discoverer

- 90. A part type discoverer applies to a simply indexed domain, *index*, and yields
  - (a) a set of type names
  - (b) each paired with a part (sort) observer.

#### value

- 90.  $\mathbb{PART\_SORTS}$ : Index  $\overset{\sim}{\to}$  **Text**
- 90.  $PART\_SORTS(\ell^{\uparrow}\langle T \rangle)$ :
- 90(a).  $tns:\{T_1,T_2,...,T_m\}:TN-set \times$
- 90(b). { obs\_T<sub>j</sub>:  $T \rightarrow T_j \mid T_j$ :tns}

**Example 53 (Some Part Sort Discoveries)** We apply a concrete version of the above sort discoverer to the road-pricing transport domain  $\Delta$ :

```
PART_SORTS(\langle \Delta \rangle):
   type
       N, F, M
   value
       obs N: \Lambda \to N
       obs F: \Delta \to F
       obs M: \Lambda \to M
PART\_SORTS(\langle \Delta, F \rangle):
   type
       V_{S}
   value
       obs Cs: F \rightarrow Vs
```

```
PART_SORTS(\langle \Delta, N \rangle):

type

Hs, Ls

value

obs_Hs: N \to Hs

obs_Ls: N \to Ls
```

# 5.2.1.2. Domain Part Types and Their Observers 5.2.1.2.1. Do a Sort Have a Concrete Type?

- Sometimes we find it expedient
  - to endow a "discovered" sort with a concrete type expression, that is,
  - "turn" a sort definition into a concrete type definition.
- 91. Thus we introduce the "discoverer":
- 91  $\text{HAS\_A\_CONCRETE\_TYPE: Index} \rightarrow \mathbf{Bool}$
- 91  $\text{HAS\_A\_CONCRETE\_TYPE}(\ell^{\hat{}}\langle t \rangle): \mathbf{true}|\mathbf{false}|$

# **Example 54 (Some Type Definition Discoveries)** We exemplify

two true expressions:

$$\begin{split} &\texttt{HAS\_A\_CONCRETE\_TYPE}(\langle \Delta, N, Hs \rangle) \\ &\texttt{HAS\_A\_CONCRETE\_TYPE}(\langle \Delta, N, Ls \rangle) \\ &\sim &\texttt{HAS\_A\_CONCRETE\_TYPE}(\langle \Delta, N, Hs, H \rangle) \\ &\sim &\texttt{HAS\_A\_CONCRETE\_TYPE}(\langle \Delta, N, Ls, L \rangle) \end{split}$$

# 5.2.1.2.2. A Domain Part Type Observer

- The PART\_TYPES( $\ell^{\hat{}}(t)$ ) invocation yields one or more sort definitions of part types together with their observer functions.
  - The domain analyser can decide that some parts can be immediately analysed into concrete types.
  - Thus, together with yielding a type name, the PART\_TYPES can be expected to yield also a type definition, that is, a type expression (paired with the type name).
  - Not all type expressions make sense.
  - We suggest that only some make sense.

- 92. The PART\_TYPES discoverer applies to a composite type, t, and yields
  - (a) a type definition, T = TE,
  - (b) together with the sort and/or type definitions of so far undefined type names of TE.
  - (c) The PART\_TYPES discoverer is not defined if the designated sort is judged to not warrant a concrete type definition.

```
92. \mathbb{PART}_{\mathbb{T}}\mathbb{TYPES}: Index \stackrel{\sim}{\to} \mathbf{Text}
```

92.  $PART_TYPES(\ell^{\uparrow}(t))$ :

```
92(a). \mathbf{type} \ t = te,
```

92(b). 
$$T_1 \text{ or } T_1 = TE_1$$

92(b). 
$$T_2 \text{ or } T_2 = TE_2$$

- 92(b). ...
- 92(b).  $T_n \text{ or } T_n = TE_n$
- 92(c). **pre**:  $\mathbb{HAS\_A\_CONCRETE\_TYPE}(\ell^{\hat{t}})$

#### **Example 55 (Some Part Type Discoveries)** We exemplify two discoveries:

```
PART_TYPES(\langle \Delta, N, Hs \rangle):
    type
         Н
         Hs = H-set
PART_TYPES(\langle \Delta, N, Ls \rangle):
    type
         Ls = L-set
PART_TYPES(\langle \Delta, F \rangle):
    type
         Vs = V-set
```

# 5.2.1.2.3. Concrete Part Types

- In Example 55 on the preceding page we illustrated one kind of concrete part type: sets.
- Practice shows that sorts often can be analysed into sets.
- Other analyses of part sorts are
  - Cartesians,
  - list, and
  - simple maps:

```
90(b). te: tn1 \times tn2 \times ... \times tnm
```

90(b). te: tn\*

90(b). te: Token  $\rightarrow$  tn

- where
  - tn's are part type usually sort names
  - some of which may have already been defined,
  - and where **Token** is some simple atomic (non-part) type.

# 5.2.1.3. Part Type Analysers

- There are three kinds of analysers:
  - unique identity analysers,
  - mereology analysers and
  - general attribute analysers. and

#### 5.2.1.3.1. Unique Identity Analysers

- We associate with every part type T, a unique identity type TI.
- 93. So, for every part type T we postulate a unique identity analyser function uid\_TI.

#### value

```
93. UNIQUE_ID: Index \rightarrow Text
```

93.  $UNIQUE_ID(\ell^{\uparrow}\langle T \rangle)$ :

93. **type** 

93. TI

93. value

93.  $uid_TI: T \rightarrow TI$ 

# 5.2.1.3.2. Mereology Analysers

- Given a part, p, of type T, the mereology, MEREOLOGY, of that part
  - is the set of all the unique identifiers of the other parts to which part p is partship-related
  - as "revealed" by the mereo\_ $\mathsf{TI}_i$  functions applied to p.
- 94. Let types  $T_1, T_2, \ldots, T_n$  be the types of all parts of a domain.
- 95. Let types  $TI_1, TI_2, \ldots, TI_n^{22}$ , be the types of the unique identifiers of all parts of that domain.
- 96. The mereology analyser MEREOLOGY is a generic function which applies to an index and yields the set of
  - (a) zero,
  - (b) one or
  - (c) more

mereology observers.

<sup>&</sup>lt;sup>22</sup>We here assume that all parts have unique identifications.

#### type

```
94. T = T_1 | T_2 | ... | T_n
95. T_{idr} = TI_1 | TI_2 | ... | TI_n
96. MEREOLOGY: Index \rightarrow Text
96. MEREOLOGY(\{\ell_i \hat{\ } \langle T_i \rangle, ..., \ell_k \hat{\ } \langle T_l \rangle \}):
96(a).
                 either: {}
96(b). or:
                             mereo_TI_x: T \to (TI_x|TI_x-set)
96(c). or: { mereo_TI<sub>x</sub>: T \to (TI_x | TI_x - \mathbf{set}),
96(c).
                              mereo_TI_{y}: T \to (TI_{y}|TI_{y}-\mathbf{set}),
96(c).
96(c).
                              mereo_TI_z: T \rightarrow (TI_z|TI_z-set) }
```

• where none of  $\mathsf{TI}_x$ ,  $\mathsf{TI}_y$ , ...,  $\mathsf{TI}_z$  are equal to  $\mathsf{TI}$  and each is some  $\mathsf{T}_{idx}$ .

# 5.2.1.3.3. General Attribute Analysers

- A general attribute analyser analyses parts beyond their unique identities and possible mereologies.
- 97. Part attributes have names. We consider these names to also abstractly name the corresponding attribute types, that is, the names function both as attribute names and sort names. Finally we allow attributes of two or more otherwise distinct part types to be the same.
- 98. ATTRIBUTES applies to parts of any part type t and yields
- 99. the set of attribute observer functions attr\_at, one for each attribute sort at of t.

#### type

97. 
$$AT = AT_1 | AT_2 | ... | AT_n$$

#### value

- 98. ATTRIBUTES: Index  $\rightarrow$  **Text**
- 98. ATTRIBUTES( $\ell^{\uparrow}(T)$ ):
- 99. **type**
- 99.  $AT_1, AT_2, ..., AT_m$
- 99. value
- 99.  $\operatorname{attr}_{A}T_{1}: T \to AT_{1}$
- 99.  $\operatorname{attr\_AT_2}: T \to \operatorname{AT_2}$
- 99. ...,
- 99.  $\operatorname{attr} AT_m: T \to AT_m, m \leq n$

**Example 56 (Example Part Attributes)** We exemplify attributes of composite and of atomic parts:

```
ATTRIBUTES(\langle \Delta \rangle):

type

Domain_Name, ...

value

attr_Name: \Delta \rightarrow Domain_Name

...
```

- where
  - Domain\_Name could include State Roads or Rail Net.
  - etcetera.

# $ATTRIBUTES(\langle \Delta, N \rangle)$ :

#### type

Sub\_Domain\_Location, Sub\_Domain\_Owner, Kms, ...

#### value

attr\_Location: N → Sub\_Domain\_Location

attr\_Owner: N → Sub-Domain\_Owner

attr\_Length:  $N \rightarrow Kms$ 

. . .

#### • where

- Sub\_Domain\_Location could include Denmark,
- Sub\_Domain\_Owner could include The Danish Road Directorate<sup>23</sup>, respectively BaneDanmark<sup>24</sup>,
- etcetera.

 $<sup>^{23}</sup> http://www.vejdirektoratet.dk/roaddirectorate.asp?page=dept\&objno=1024$ 

<sup>&</sup>lt;sup>24</sup>http://uk.bane.dk/default\_eng.asp?artikelID=931

# ATTRIBUTES( $\langle \Delta, N, Hs, L \rangle$ ): **type**LOC, LEN, ... **value**attr\_LOC: L $\rightarrow$ LOC attr\_LEN: L $\rightarrow$ LEN ...

```
ATTRIBUTES(\{\langle \Delta, N, Hs, L \rangle, \langle \Delta, N, Hs, H \rangle\}):

type
L\Sigma = \text{HI-set}, L\Omega - L\Sigma\text{-set}
H\Sigma = \text{LI-set}, H\Omega - H\Sigma\text{-set}
value
\text{attr} L\Sigma: L \to L\Sigma
\text{attr} L\Omega: L \to L\Omega
\text{attr} H\Sigma: H \to H\Omega
```

#### • where

- LOC might reveal some Bezier curve $^{25}$  representation of the possibly curved three dimensional location of the link in question,
- LEN might designate length in meters,
- $-L\Sigma$  [H $\Sigma$ ] designates the state of the link [hub],
- $-L\Omega$  [H $\Omega$ ] designates the space of all allowed states of the link [hub], etcetera.

 $<sup>^{25}</sup>$ http://en.wikipedia.org/wiki/Bézier\_curve

#### $\star$ Attribute Sort Exploration $\star$

- Once the attribute sorts of a part type have been determined
  - there remains to be "discovered" the concrete types of these sorts.
  - We omit treatment of this point in the present version of these these lectures.

# 5.2.2. Discovering Action Signatures 5.2.2.1. General

- We really should discover actions, but actually analyse function definitions.
- And we focus, in these lectures, on just "discovering" the function signatures of these actions.
- By a function signature, to repeat, we understand
  - a functions name, say fct, and
  - a function type expression (te), say  $dte \xrightarrow{\sim} rte$  where
    - \* dte defines the type of the function's definition set
    - \* and rte defines the type of the function's image, or range set.

# 5.2.2.2. Function Signatures Usually Depend on Compound Domains

- We use the term 'functions' to cover actions, events and behaviours.
- We shall in general find that the signatures of actions, events and behaviours depend on types of more than one domain.
  - Hence the schematic index set  $\{\ell_1 \hat{t}_1, \ell_2 \hat{t}_2, ..., \ell_n \hat{t}_n \}$
  - is used in all actions, events and behaviours discoverers.

#### 5.2.2.3. The ACTION SIGNATURES Discoverer

- 100. The ACTION\_SIGNATURES meta-function applies to an index set and yields
  - (a) a set of action signatures each consisting of an action name and a pair of definition set and range type expressions where
  - (b) the type names that occur in these type expressions are defined by in the domains indexed by the index set.

```
ACTION\_SIGNATURES: Index \xrightarrow{\sim} Text
100
             \mathbb{ACTION\_SIGNATURES}(\{\ell_1^{\wedge}\langle T_1\rangle, \ell_2^{\wedge}\langle T_2\rangle, ..., \ell_n^{\wedge}\langle T_n\rangle\}):
100
                              \operatorname{act\_fct}_i: \operatorname{te}_{i_d} \xrightarrow{\sim} \operatorname{te}_{i_r},
100(a)
                              \operatorname{act\_fct}_i: \operatorname{te}_{i_d} \xrightarrow{\sim} \operatorname{te}_{i_r},
100(a)
100(a)
                              \operatorname{act\_fct}_k: \operatorname{te}_{k_d} \xrightarrow{\sim} \operatorname{te}_{k_r}
100(a)
100(b)
                  where:
100(b)
                               type names in (te_{(i|j|...|k)_d}) and in (te_{(i|j|...|k)_r}) are
100(b)
                               type names defined by the indices which are prefixes of
100(b)
                               \ell_m \hat{\ } \langle T_m \rangle and where \mathsf{T}_m is in some signature \mathsf{act\_fct}_{i|j|...|k}.
```

# 5.2.3. Discovering Event Signature

- Events are from the point of view of signatures very much like actions.
- 101. The EVENT\_SIGNATURES meta-function applies to an index set and yields
  - (a) a set of action signatures each consisting of an action name and a pair of definition set and range type expressions where
  - (b) the type names that occur in these type expressions are defined either in the domains indexed by the index set or by the environment (i.e., "outside" the domain  $\Delta$ ).

```
EVENT_SIGNATURES: Index \stackrel{\sim}{\to} Text
101
           EVENT_SIGNATURES(\{\ell_1 \ (T_1), \ell_2 \ (T_2), ..., \ell_n \ (T_n)\}):
101
                         \operatorname{evt\_fct}_i: \operatorname{te}_{i_d} \xrightarrow{\sim} \operatorname{te}_{i_r},
101(a)
                         \operatorname{evt\_fct}_{j}: \operatorname{te}_{j_d} \xrightarrow{\sim} \operatorname{te}_{j_r},
101(a)
101(a)
                         \operatorname{evt\_fct}_k: \operatorname{te}_{k_d} \xrightarrow{\sim} \operatorname{te}_{k_r}
101(a)
101(b)
                     where:
                          type names of te_{(i|j|...|k)_d} and te_{(i|j|...|k)_r} are type names
101(b)
101(b)
                          defined by the indices which are prefixes of \ell_m \hat{t}_m
101(b)
                          and where t_m is in some signature act_i fct_{i|i|...|k} or may
                          refer to types definable only "outside" \Delta
101(b)
```

# 5.2.4. Discovering Behaviour Signatures

- We choose, in these lectures, to model behaviours in  $CSP^{26}$ .
- This means that we model (synchronisation and) communication between behaviours by means of messages m of type M, CSP channels (channel ch:M) and CSP

output: ch!e [offer to deliver value of expression e on channel ch], and input: ch? [offer to accept a value on channel ch].

ullet We allow for the declaration of single channels as well as of one, two, ..., n dimensional arrays of channels with indexes ranging over channel index types

```
type Idx, Cldx, Rldx ...:
  channel ch:M, { ch_v[vi]:M'|vi:Idx }, { ch_m[ci,ri]:M"|ci:Cldx,ri:Rldx }, ...
etcetera.
```

• We assume some familiarity with RSL/CSP.

<sup>&</sup>lt;sup>26</sup>Other behaviour modelling languages are Petri Nets, MSCs: Message Sequence Charts, Statechart etc.

• A behaviour usually involves two or more distinct sub-domains.

# **Example 57 (The Involved Subdomains of a Vehicle Behaviour)**

Let us illustrate that behaviours usually involve two or more distinct sub-domains.

- A vehicle behaviour, for example, involves
  - the vehicle subdomain,
  - the hub subdomain (as vehicles pass through hubs),
  - the link subdomain (as vehicles pass along links) and,
  - for the road pricing system, also the monitor subdomain.

- 102. The BEHAVIOUR\_SIGNATURES is a meta function.
- 103. It applies to a set of indices and results in a text,
- 104. The text contains
  - (a) a set of zero, one or more message types,
  - (b) a set of zero, one or more channel index types,
  - (c) a set of zero, one or more channel declarations,
  - (d) a set of one or more process signatures with each signature containing a behaviour name, an argument type expression, a result type expression, usually just **Unit**, and
  - (e) an input/output clause which refers to channels over which the signatured behaviour may interact with its environment.

```
BEHAVIOUR_SIGNATURES: Index \stackrel{\sim}{\to} Text
103
        BEHAVIOUR_SIGNATURES(\{\ell_1 \cap \langle T_1 \rangle, \ell_2 \cap \langle T_2 \rangle, ..., \ell_n \cap \langle T_n \rangle\}):
103.
                   type M = M_1 | M_2 | ... | M_m, m > 0
104(a).
104(b).
                               I = I_1 | I_2 | ... | I_n, n > 0
104(c). channel ch,vch[i],\{\text{vch}[i]:M|i:I_a\},\{\text{mch}[j,k]:M|j:I_b,k:I_c\},...
104(d).
             value
104(d).
                      bhv<sub>1</sub>: ate<sub>1</sub> \rightarrow inout<sub>1</sub> rte<sub>1</sub>,
104(d).
                       bhy: ate \rightarrow inout_2 rte.
104(d).
                      ••• ,
104(d).
                       bhv<sub>m</sub>: ate<sub>m</sub> \rightarrow inout<sub>m</sub> rte<sub>m</sub>,
104(d). where
                       type expressions atei_i and rte_i for all i involve at least two
                          types t'_i and t''_i of respective indexes \ell_i \hat{t}_i and \ell_j \hat{t}_j
104(d).
104(e). where inout_i: in k | out k | in,out k
104(e). where k: ch | ch[i] | \{ch[i]|i \in I_a\} | \{mch[j,k]:M|i:I_b,j:I_c\} | ...
```

**Example 58 (A Vehicle Behaviour Signature Discovery)** We refer, for example, to Examples 36 (Slides 184–198) and 40 (Slides 215–219).

```
let ih=\langle \Delta.N.LS,H \rangle,il=\langle \Delta.N.HS,L \rangle,iv=\langle \Delta,F,V \rangle,im=\langle \Delta,Monitor \rangle in
BEHAVIOUR_SIGNATURES({iv,ih,iv,im}) as text
   let n:N, hs=obs_HS(n), ls=obs_LS(n), vs=obs_F(PART_SORTS)(\langle \Delta \rangle) in
   where text:
        type
           VL_Msg, VH_Msg, VM_Msg
        channel
   62(a). \{\text{vh\_ch}[\text{attr\_VI(v),attr\_HI(h)}]|\text{v:V,h:H•v} \in \text{vd} \land \text{h} \in \text{hs}\}:\text{VH\_Msg}
   62(b). \{\text{vl\_ch}[\text{attr\_VI(v),attr\_LI(h)}]|\text{v:V,l:L•v} \in \text{vs} \land \text{h} \in \text{ls}\}:\text{VL\_Msg}
   62(c). m_ch:VM_Msg
        value
               vehicle: VI \rightarrow V \rightarrow VP \rightarrow
   64
                   out,in \{vl\_ch[vi,li]|li:LI·li \in xtr\_LIs(ls)\}
   83.
                              { vh_ch[vi,hi]|hi:HI \cdot hi \in xtr_HIs(hs) }  out m_ch,... Unit
   83.
end end
```

# 5.3. What Does Application Mean?

- Now what does it actually mean "to apply" a discover function?
- We repeat our list of discoverers.

```
[90 (Slide 255)] PART_SORTS,

[91 (Slide 257)] HAS_A_CONCRETE_TYPE,

[92 (Slide 260)] PART_TYPES,

[93 (Slide 264)] UNIQUE_ID,

[96 (Slide 265)] MEREOLOGY,

[98 (Slide 267)] ATTRIBUTES,

[100 (Slide 275)] ACTION_SIGNATURES,

[101 (Slide 276)] EVENT_SIGNATURES and

[103 (Slide 280)] BEHAVIOUR_SIGNATURES.
```

• It is the domain engineer cum scientist who "issues" the "commands".

- The first "formal" domain inquiry is that of PART\_SORTS( $\langle \Delta \rangle$ ).
  - We refer to Item 90 on page 255, for example as captured by the formulas, Items 90–90(b) (Slide 285).
- For the domain engineer to 'issue' one of the 'discovery commands' means that that person has
  - (i) prepared his mind to study the domain and is open to impressions,
  - (ii) decided which  $\mathcal{DISCOVERER\_KIND}$  to focus on, and
  - -(iii) studied the "rules of engagement" of that command, that is
    - \* which pre-requisite discoverers must first have been applied,
    - \* with which index, that is, in which context the command invocation should be placed,
    - \* and which results the invocation is generally expected to yield.

#### 5.3.1. PART\_SORTS

• Let us review the PART\_SORTS discoverer:

#### value

```
90. PART_SORTS: Index \stackrel{\sim}{\to} Text
90. PART_SORTS(\ell \cap \langle T \rangle):
90(a). tns:\{T_1, T_2, ..., T_m\}:TN-set \times
90(b). \{\text{ obs-}T_j: T \to T_j \mid T_j: \text{tns}\}
```

- The domain analyser has decided to "position" the search at domain index  $\ell^{\hat{}}\langle \mathsf{T} \rangle$ 
  - where  $T = \Delta$  if  $\ell = \langle \rangle$  and
  - where T is some "previously discovered part type.

- From Item 90(a) the domain analyser is guided (i.e., advised) to analyse the domain "at position  $\ell^{\hat{}}\langle \mathsf{T} \rangle$ :
  - is the domain type T a type of one or more subpart types?
    - \* If so then decide which they are, that is:  $\mathsf{T}_1, \mathsf{T}_2, ..., \mathsf{T}_m$ , that is, the "generation" of the text **type**  $\mathsf{T}_1, \mathsf{T}_2, ..., \mathsf{T}_m$ ,
    - \* if not then tns={} and no text is "generated".
- Item 90(b), and given the domain analyser's resolution of Item 90(a), then directs
  - the "generation" of m observers  $\mathsf{obs}_{\mathsf{L}}\mathsf{T}_{j}$ :  $\mathsf{T} \to \mathsf{T}_{j}$  (for  $j:\{1..m\}$ ).

#### 5.3.2. HAS\_A\_CONCRETE\_TYPE

- Let us review the HAS\_A\_CONCRETE\_TYPE analyser:
- 91  $\text{HAS\_A\_CONCRETE\_TYPE: Index} \rightarrow \mathbf{Bool}$
- 91  $\mathbb{HAS\_A\_CONCRETE\_TYPE}(\ell^{\uparrow}\langle T \rangle)$ :**true**|**false** 
  - Item 91 directs the domain analyser to decide
    - whether the domain type T at "position"  $\ell \hat{t}$
    - should be given a concrete type definition.
  - It is a decision sôlely at the discretion of the domain analyser
    - whether domain type T should be given a concrete type definition,
    - and which concrete type it should then be "given",
    - that is, how it should be "concretely abstractly" modelled.

#### 5.3.3. PART\_TYPES

• Let us review the PART\_TYPES analyser:

```
92. \mathbb{PART}_{TYPES}: \operatorname{Index} \xrightarrow{\sim} \mathbf{Text}

92. \mathbb{PART}_{TYPES}(\ell^{\wedge}\langle t \rangle):

92(a). \mathbf{type} \ T = TE,

92(b). T_1 \text{ or } T_1 = TE_1

92(b). T_2 \text{ or } T_2 = TE_2

92(b). ...

92(b). T_n \text{ or } T_n = TE_n

92(c). \mathbf{pre}: \mathbb{HAS}_{A}_{ONCRETE}
```

- The domain analyser has decided to "position" the search at domain index  $\ell^{\hat{}}\langle \mathsf{T} \rangle$ 
  - where  $T = \Delta$  if  $\ell = \langle \rangle$  and
  - where T is some "previously discovered part type.
- From Item 92 the domain analyser is guided (i.e., advised) to analyse the domain "at position  $\ell^{\hat{}}\langle \mathsf{T} \rangle$ :
  - -can an abstract, yet concrete type definion be given for T?
  - If so then decide which it should be, that is, should it be an atomic type
    - \* a number type, Intg, Rat, Real,
    - \* a Boolean type, **Bool**, or
    - \* a token type, f.ex. TOKEN;

- or should it be a composite type

```
* either a set type: TE: T_s-set of te: T_s-infset,

* or a Cartesian type: TE: T_1 \times T_2 \times ... T_m,

* or a list type: TE: T_\ell^* or te: T_\ell^\omega

* or a map type: TE: T_d \xrightarrow{m} T_t?
```

• In either case the text

$$\begin{array}{lll} - \mbox{ type } T = TE, \\ - & T_1 \mbox{ or } T_1 \! = \! TE_1, \\ - & T_2 \mbox{ or } T_1 \! = \! TE_2, \\ - & ..., \\ - & T_n \mbox{ or } T_n \! = \! TE_n \end{array}$$

is generated

- where  $\mathsf{TE}\ (\mathsf{TE}_x)$  is a type expression whose
- so far undefined type names  $T_1$ ,  $T_2$ , ...,  $T_n$  must be defined, either as sorts, or a concrete types.

# 5.3.4. UNIQUE\_ID

• Let us review the UNIQUE\_ID analyser:

#### value

```
93. UNIQUE_ID: Index \rightarrow Text
93.a UNIQUE_ID(\ell^{\hat{}}\langle T \rangle):
93.b type
93.c TI
93.d value
93.e uid_TI: T \rightarrow TI
```

- Item 93.a inquires as to the
- Line 93.b type name
- Line 93.c of the inquired part type's unique identifiers
- Line 93.d and the function signature value
- Line 93.e of
  - the observer, uid\_TI, name,
    Thus, the only real "new"
    the definition set type (T, of course) and
    "discovery" here is the name,
    TI, of the unique identifier type.

• Etcetera, etcetera.

#### 5.4. Discussion

• We have presented a set of discoverers:

```
[90 (Slide 255)] PART_SORTS,
```

[91 (Slide 257)] HAS\_A\_CONCRETE\_TYPE,

[92 (Slide 260)] PART\_TYPES,

[93 (Slide 264)] UNIQUE\_ID,

[96 (Slide 265)] MEREOLOGY,

[98 (Slide 267)] ATTRIBUTES,

[100 (Slide 275)] ACTION\_SIGNATURES,

[101 (Slide 276)] EVENT\_SIGNATURES and

[103 (Slide 280)] BEHAVIOUR\_SIGNATURES.

- There is much more to be said:
  - About a meta-state component in which is kept the "text" so far generated.
  - A component from which one can see which indices and hence which type names have so far been "generated",
  - and on the basis of which one can perform tests of well-formedness of generated text,
  - etcetera, etcetera,

# **End Lecture 6: Discovering Domain Entities**